New fairness concepts for allocating indivisible items

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This talk ...

- An overview of well-known results in fair division with indivisible items
- New technical results (C., Garg, Rathi, Sharma, & Varriccio, 2023)

Fair division: some indicative problems

- An inheritance, consisting of a jewellery collection, pieces of antique furniture, and estate property, is to be divided among heirs
- Food donated to a food bank has to be given to charities
- Access to rainwater reservoirs has to be granted to farmers
- A territorial dispute has to be resolved between neighbouring countries
- A partnership is dissolved, and the ex-partners have to split assets and liabilities
- Responsibility for the protection of refugees has to be shared among EU countries

The research agenda: conceptual and computational challenges in fair division

- Computational questions: How should fair division procedures for these scenarios work?
- Before that: need to define fairness as a concept

Allocating indivisible items

The basic setting

• Indivisible items











Agents with valuations for the items (additivity)

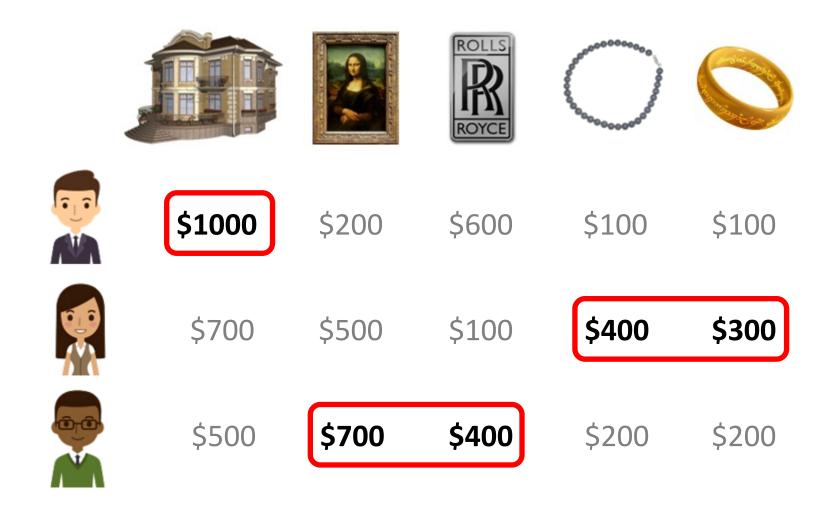


• Goal: divide the items among the agents in a fair manner

An example



An example





- Two interpretations of fairness:
 - Comparative: to evaluate an allocation as fair, each agent compares the bundle of items allocated to her to the bundles allocated to other agents
 - In absolute terms: each agent defines a threshold value based on her view of the items to be allocated and evaluates as fair those allocations which give her value higher than the threshold



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- Fairness notions
 - Envy freeness: every agent prefers her bundle to that given to any other agent $\forall i, j \colon v_i(A_i) \geq v_i(A_j)$





- Two interpretations of fairness:
 - Comparative: to evaluate an allocation as fair, each agent compares the bundle of items allocated to her to value of agent i for the to other agents
 - For every pair of agents i and j bundle A_i allocated to her bundle A_j allocated to agent j and j allocated to agent j
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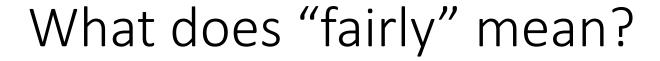


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Fairness notions

- Envy freeness: every agent prefers her bundle to that given to any other agent $\forall i, j \colon v_i(A_i) \geq v_i(A_j)$
- Proportionality: every agent feels that she gets at least 1/n-th of all items

$$\forall i: v_i(A_i) \ge \frac{1}{n} v_i(G)$$





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Threshold value: 1/nth of the total value of

- Fair For every agent i
 - gent preference $\forall i,j \colon \iota \quad A_i) \geq v_i(A_j)$ agent i for all items
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Fairness notions

Unfortunately, envy free and proportional allocations may not exist

When is a fairness concept important/useful?

- Must be fair ©
- Should always exist
- Must be efficiently computable

Relaxing envy-freeness



 Envy freeness up to some item (EF1): every agent prefers her own bundle to the bundle of any other agent after eliminating some item from the latter

$$\forall i, j: \exists g \in A_j \text{ s. t. } v_i(A_i) \geq v_i(A_j \setminus \{g\})$$

Proposed by Budish (2011)





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$$\forall i, j: \exists g \in A_j \text{ s.t. } v_i(A_i) \geq v_i(A_j \setminus \{g\})$$

- Proposed by Budish (2011)
- EF1 always exist and can be computed in polynomial time
- Via the draft mechanism (folklore), envy-cycle elimination (Lipton, Markakis, Mossel, & Saberi, 2004), the maximum Nash welfare allocation (C., Kurokawa, Moulin, Procaccia, Shah, & Wang, 2019)

• Drafting order:































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• Drafting order:































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• Drafting order:





























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• Drafting order:































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• Drafting order:































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Envy-cycle elimination

- Allocate items one by one
- In each step *j*:
 - Allocate item j to an agent that nobody envies
 - If this creates a "cycle of envy", redistribute the bundles along the cycle
- Crucial property:
 - Envy can be eliminated by removing just a single good
 - Implies EF1
- Lipton, Markakis, Mossel, & Saberi (2004)

Relaxing envy-freeness



- Envy freeness up to any item (EFX): every agent prefers her own bundle to the bundle of any other agent after eliminating any item from the latter $\forall i, j, \forall g \in A_i : v_i(A_i) \geq v_i(A_j \setminus \{g\})$
- Proposed by C., Kurokawa, Moulin, Procaccia, Shah, & Wang (2019), Gourves, Monnot, & Tilane (2014)

Relaxing envy-freeness



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- Proposed by C., Kurokawa, Moulin, Procaccia, Shah, & Wang (2019), Gourves, Monnot, & Tilane (2014)
- Not known whether it always exists for general instances
- Known results for agents with identical valuations, ordered valuations, three agents, and a few more
 - Plaut & Roughgarden (2020), Chaudhuri, Garg, & Mehlhorn (2020)
- Known results for relaxations of EFX (approximations, EFX with charity, etc.)
 - Amanatidis, Markakis, & Ntokos (2020), C., Gravin, & Huang (2019), Chaudhuri, Kayitha Mehlhorn & Sgouritsa (2021), Chaudhuri, Garg, Mehlhorn, Ruta, & Misra,





• Maximin share fairness (MMS): each agent's threshold is equal to the best guarantee when dividing the items into n bundles and getting the least valuable bundle

$$\forall i, v_i(A_i) \ge \theta_i = \max_B \min_j v_i(B_j)$$

Proposed by Budish (2011)

For every agent *i*

Agent *i*'s value is above the MMS threshold

MMS threshold = the maximum over all allocations B of the minimum value agent i has from B's bundles

• Maximin share fairness (MMS): The agent's threshold is equal to the best guarantee when dividing the item into n bundles and getting the least valuable bundles.

$$\forall i, v_i(A_i) \ge \theta_i = \max_B \min_j v_i(B_j)$$

Proposed by Budish (2011)

Relaxing proportionality



• Maximin share fairness (MMS): each agent's threshold is equal to the best guarantee when dividing the items into n bundles and getting the least valuable bundle

$$\forall i, v_i(A_i) \ge \theta_i = \max_B \min_j v_i(B_j)$$

- Proposed by Budish (2011)
- Unfortunately, MMS allocations may not exist
 - Procaccia & Wang (2014), Kurokawa, Procaccia, & Wang (2018)
- Research has focused on achieving MMS-approximations in poly time
 - Amanatidis, Markakis, Nikzad, & Saberi (2017), Ghodsi, Hajiaghayi, Seddighin, Seddighin, & Yami (2018), Barman & Krishnamurthy (2020), Garg & Taki (2020)

Summarizing so far

- EF1: always exists, easy to achieve, not fair
- EFX: not known whether it can be always satisfied, fair
- MMS: may not exist, fair (if exists)



• See Bouveret & Lemaitre (2016), Aziz, Bouveret, C., Giagkousi, & Lang (2018) for taxonomies including more fairness concepts

Summarizing so far

• EF1: always exists, easy to achieve, nq

EFX: not known whether it can be alw

• MMS: may not exist, fair (if existed

EF EFX EF1

Prop MMS

Still, EFX seems to be the most promising fairness property we have for indivisible items

• See Bouveret & Lemaitre (2016), Aziz, Bouveret, C., Giagkousi, & Lang (2018) for taxonomies including more fairness concepts

New fairness concepts

Fairness and knowledge

- What kind of knowledge do the agents need to have?
- Knowledge about the items and the number of agents only:
 - Proportionality, MMS
- Knowledge about the whole allocation:
 - EF, EFX, EF1

Epistemic envy-freeness (EEF)

- Informally: a relaxation of EF with a definition that uses only knowledge about items and number of agents
- Formal definition: the allocation $(A_1, A_2, ..., A_n)$ is EEF if, for every agent i, there is a **reallocation** $(B_1, ..., B_{i-1}, A_i, B_{i+1}, ..., B_n)$ of the items in which agent i is not envious, i.e., $v_i(A_i) \geq v_i(B_j)$ for every other agent j
- Aziz, C., Bouveret, Giagkousi, & Lang (2018)
- Unfortunately, EEF allocations may not exist

Epistemic envy-freeness up to any item (EEFX)

- Informally: a relaxation of EFX with a definition that uses only knowledge about items and number of agents
- Formal definition: the allocation $(A_1, A_2, ..., A_n)$ is EEFX if, for every agent i, there is a **reallocation** $(B_1, ..., B_{i-1}, A_i, B_{i+1}, ..., B_n)$ of the items in which the EFX conditions for agent i are satisfied

$$\forall i, j \neq i, \forall g \in B_j: v_i(A_i) \geq v_i(A_j \setminus \{g\})$$

• C., Garg, Rathi, Sharma, & Varricchio (2023)

Minimum EFX value fairness (MXS)

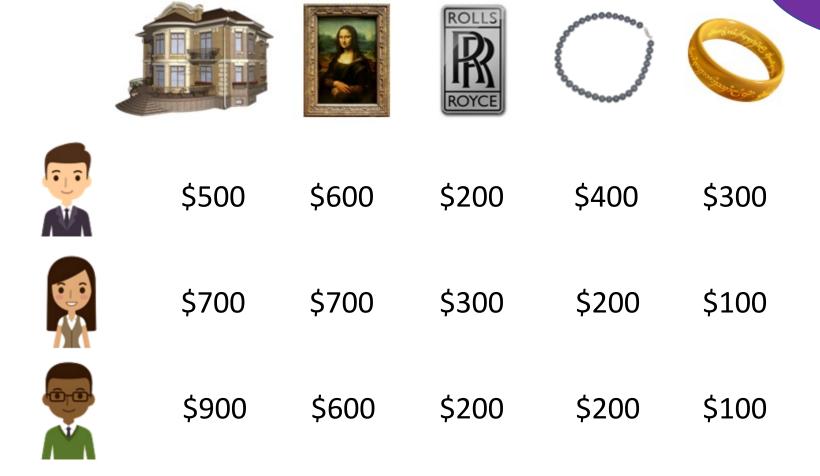
- Informally: Each agent i gets a value that is at least as high as the minimum value agent i gets among all allocations where the EFX conditions for her are satisfied
- Formal definition: the allocation $(A_1, A_2, ..., A_n)$ is MXS if $\forall i : v_i(A_i) \geq \theta_i = \min_{B \in EFX_i} v_i(B_i)$

where the set EFX_i consists of those allocations $B=(B_1,B_2,\ldots,B_n)$ such that

$$\forall j \neq i, g \in B_j: v_i(B_i) \geq v_i(B_j \setminus \{g\})$$

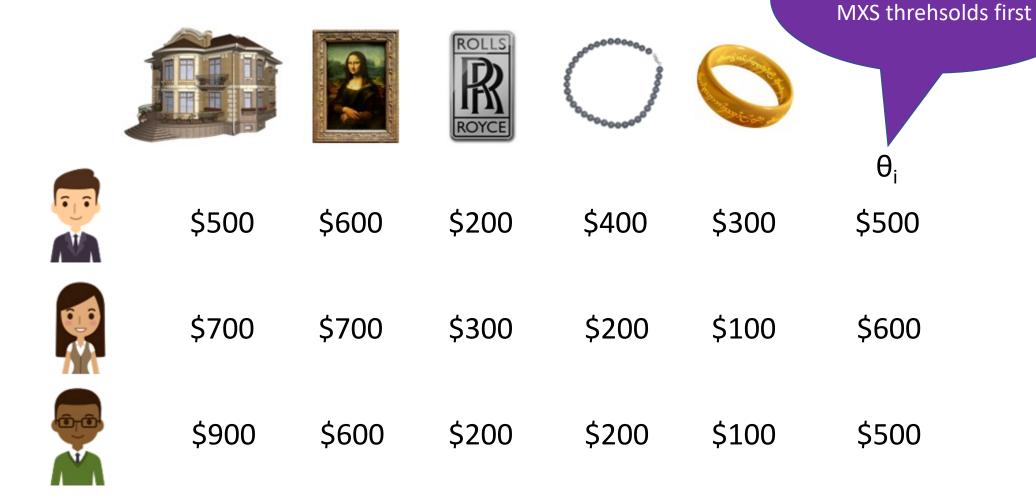
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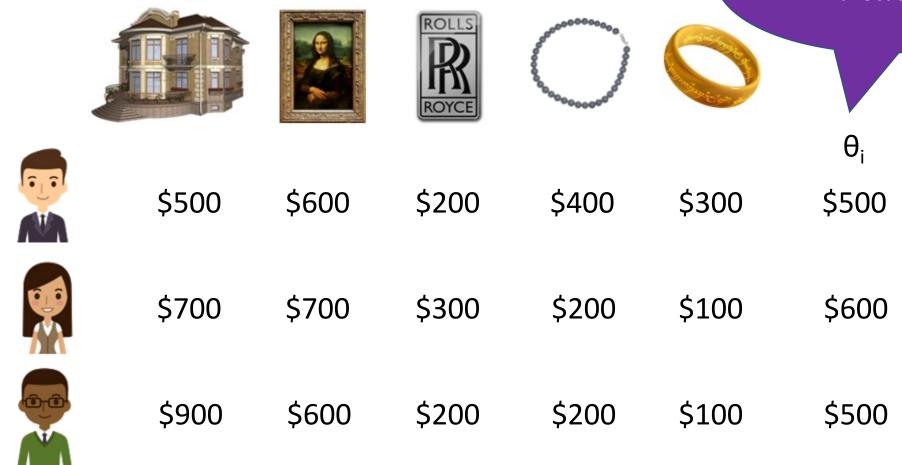


Let's compute the MXS threhsolds first

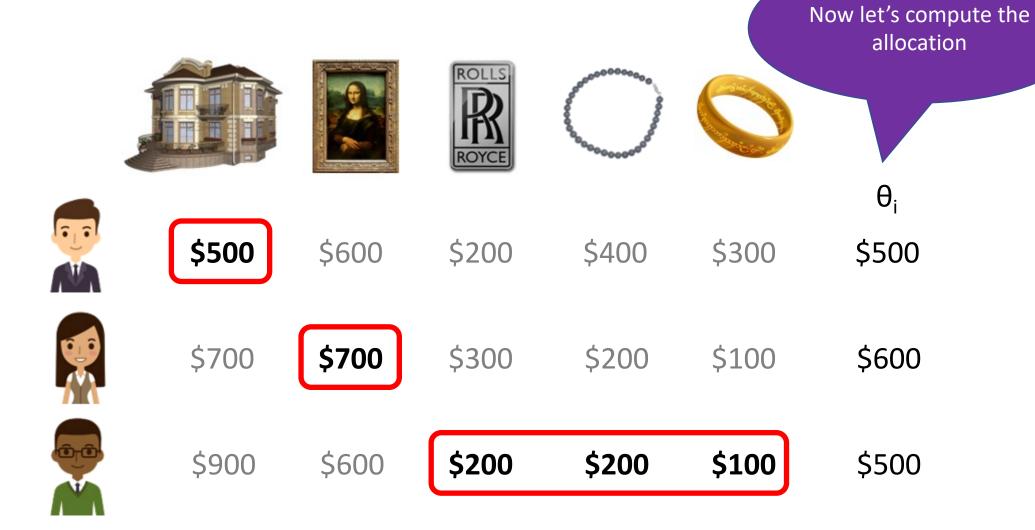
 θ_{i}



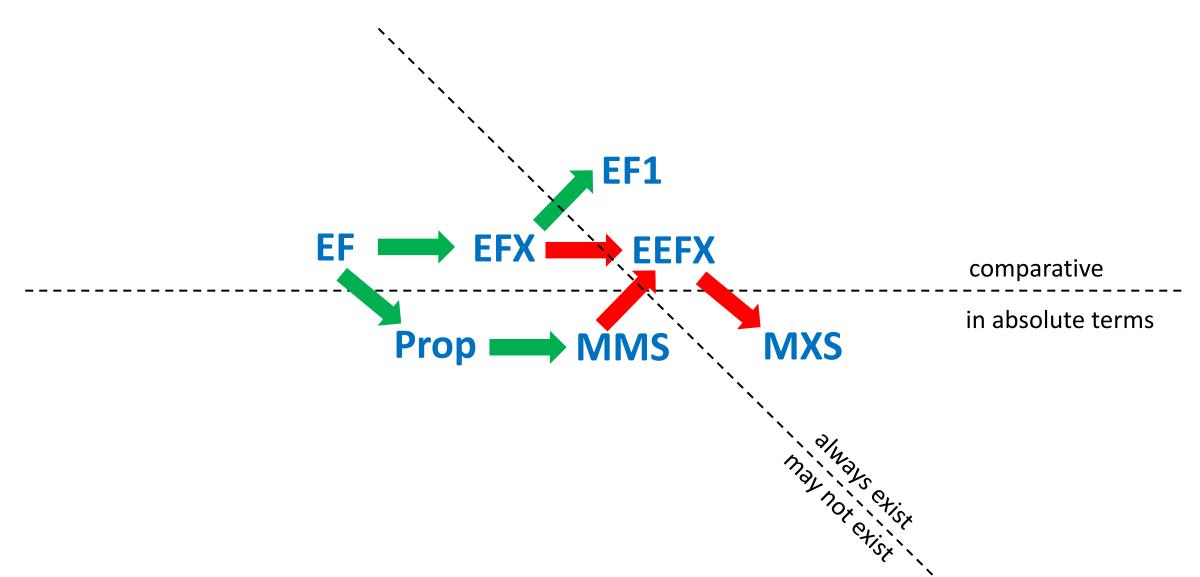
Let's compute the



Now let's compute the allocation



A geometry of fairness properties



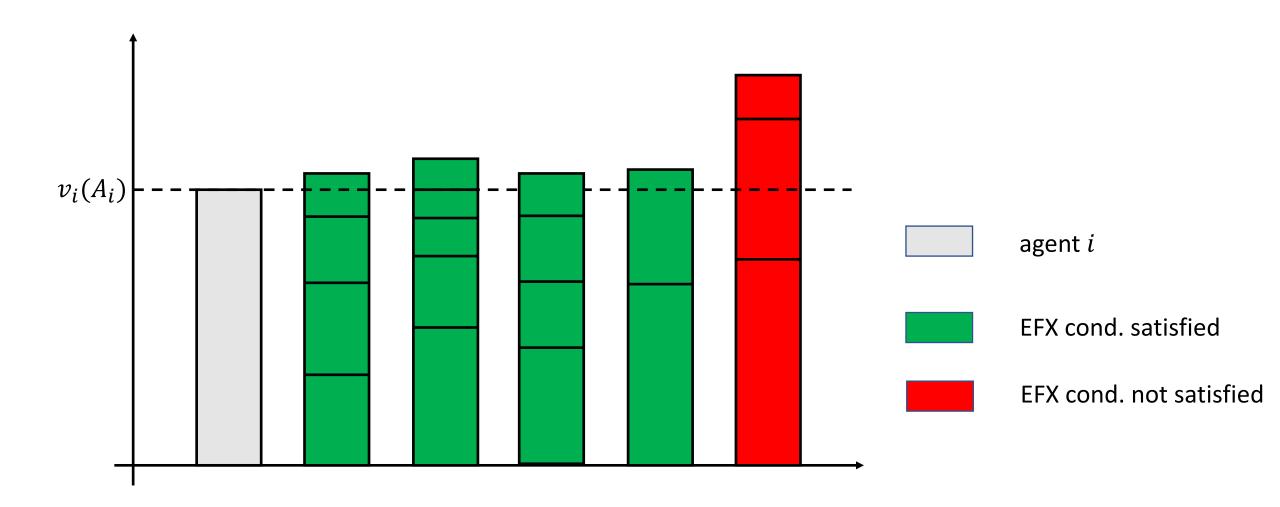
EEFX — MXS

- Proof: Let $A = (A_1, ..., A_n)$ be **EEFX**. Then, for every agent i, there exists a reallocation $B = (B_1, ..., B_{i-1}, A_i, B_{i+1}, ..., B_n)$ so that the EFX conditions are satisfied for agent $i, B \in EFX_i$
- Hence,

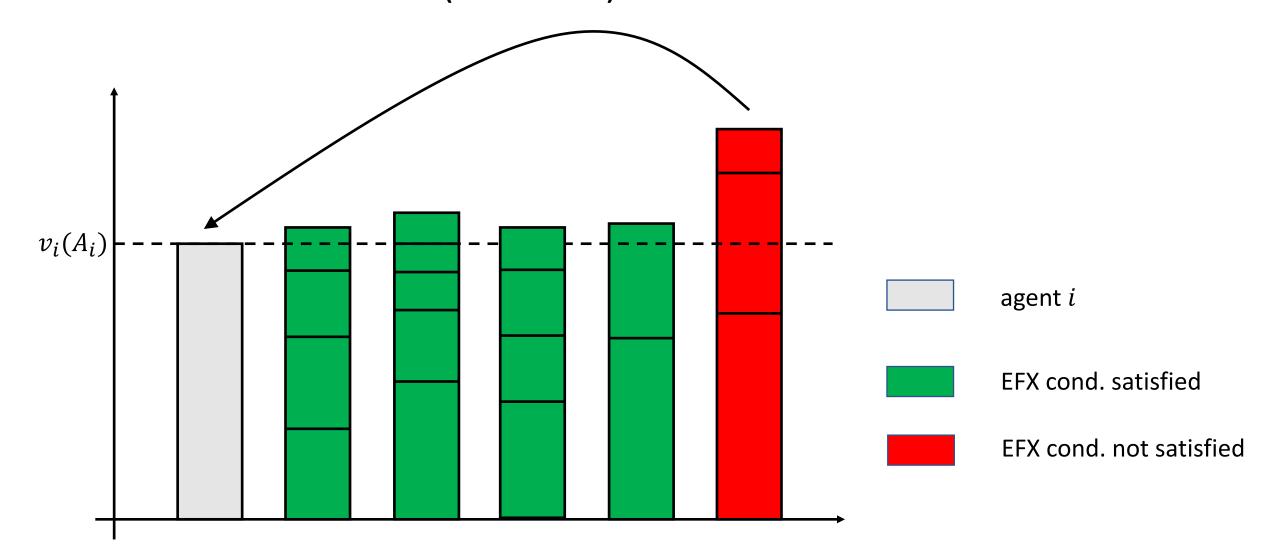
$$v_i(A_i) \ge \min_{B' \in EFX_i} v_i(B'_i) = MXS_i$$

• I.e., *A* is also **MXS**

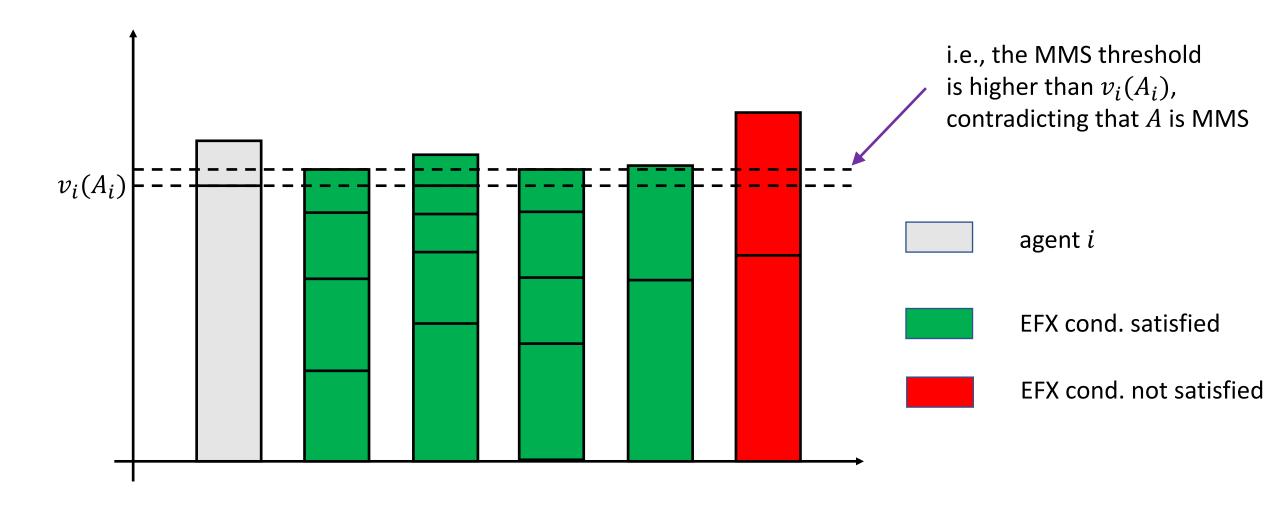
MMS — EEFX



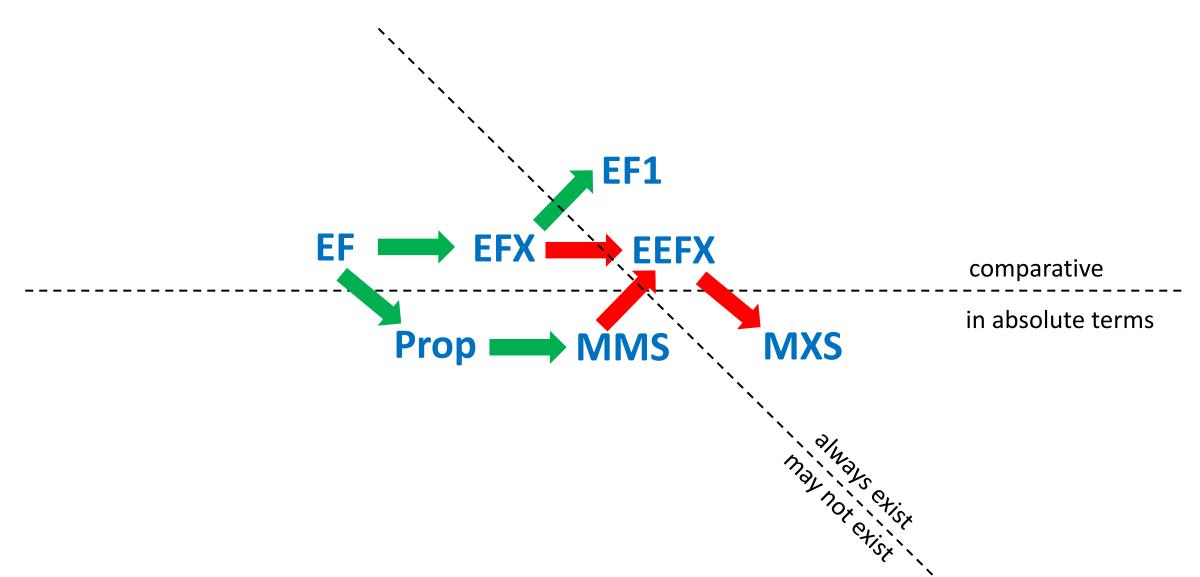
MMS — EEFX (contd.)



MMS — EEFX (contd.)



A geometry of fairness properties



Main result: EEFX and MXS are awesome!

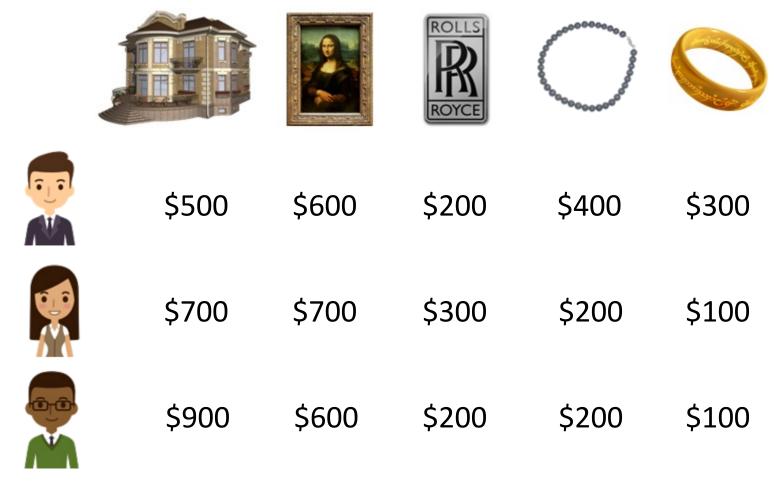
 Theorem: EEFX and MXS allocations always exist and can be computed in polynomial time

- Step 1: Enumerate the items as g_1, g_2, \dots, g_m and redistribute the values so that each agent has her j-th highest value for item g_j
- Step 2: Run envy-cycle elimination on this ordered instance
- Step 3: Redistribute the items to the bundles. For $j=1,\ldots,m$, agent who currently has item g_i is asked to pick her best available item

Envy-cycle elimination (implementation of step 2)

- Lipton, Markakis, Mossel, & Saberi (2004)
- Allocate items one by one (ordered from the most to the least valued one)
- In each step *j*:
 - Allocate item j to an agent that nobody envies
 - If this creates a "cycle of envy", redistribute the bundles along the cycle
- Crucial property:
 - Envy can be eliminated by removing a single item (the last one inserted in a bundle)
 - Implies EF1 (actually, EFX)
- Barman & Krishnamourthy (2020)

An example



Step 1: redistributing the values

	1	2	3	4	5
الم الم	\$600	\$500	\$400	\$300	\$200
	\$700	\$700	\$300	\$200	\$100
	\$900	\$600	\$200	\$200	\$100

	1	2	3	4	5
١	\$600	\$500	\$400	\$300	\$200
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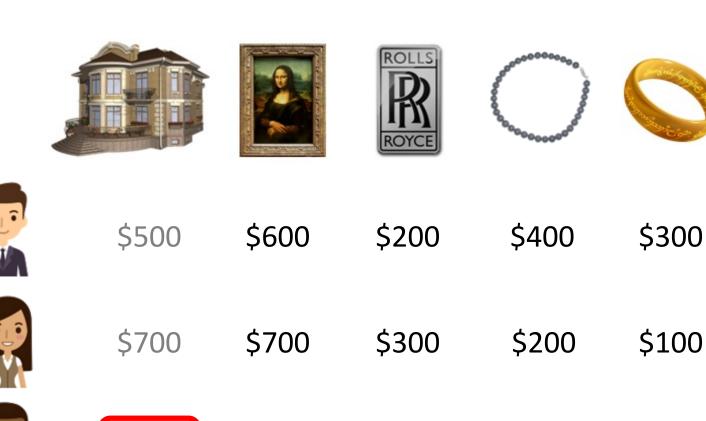
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\$600

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picking sequence

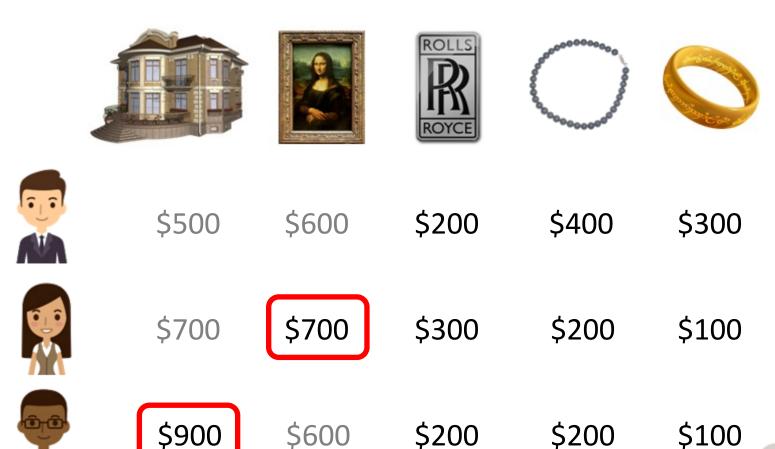


\$100







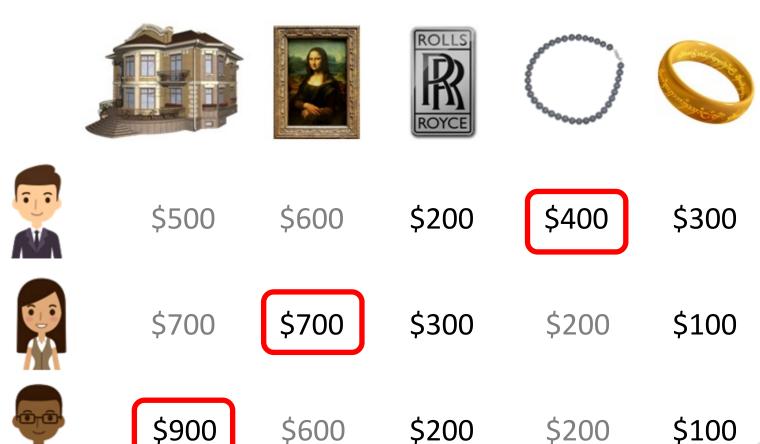










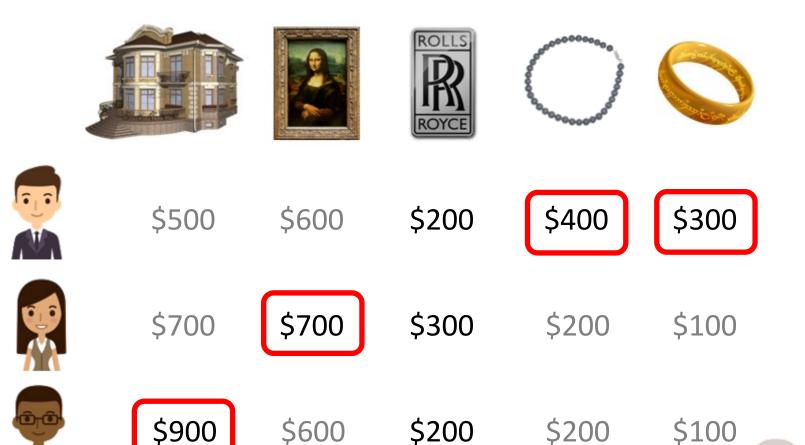










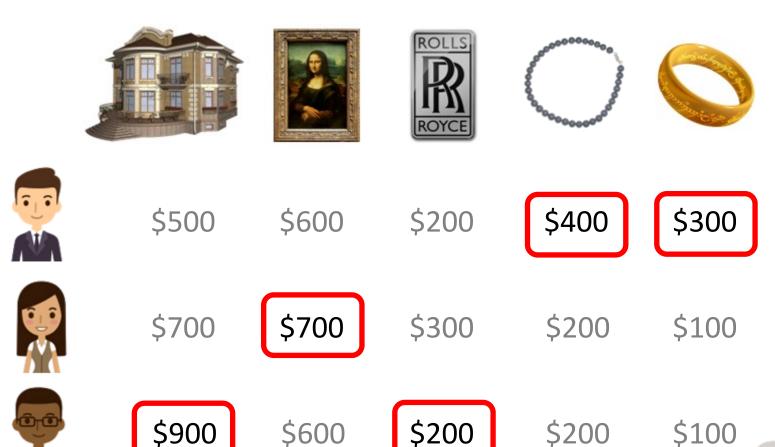


















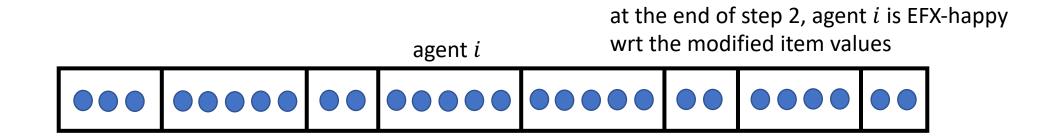


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 - Bouveret & Lemaitre (2016)

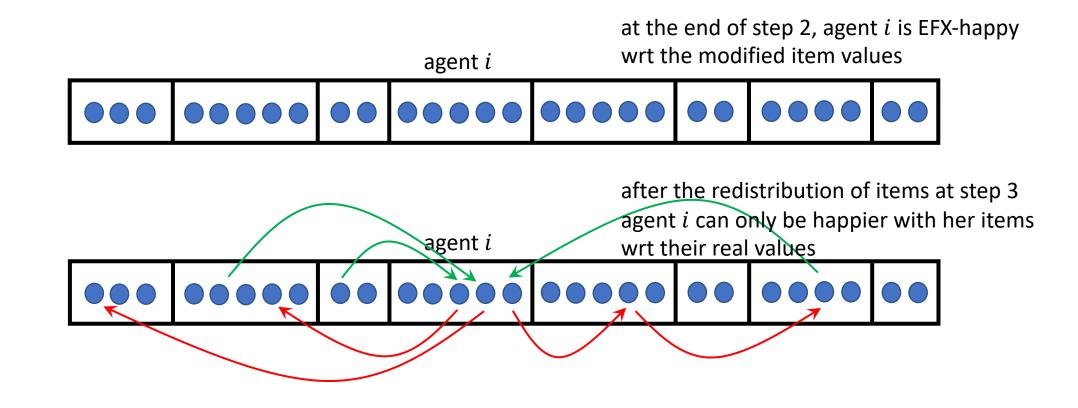
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 - Yields an EFX allocation for the ordered instance (Barman & Krishnamourthy, 2020, Plaut & Roughgarden, 2020)

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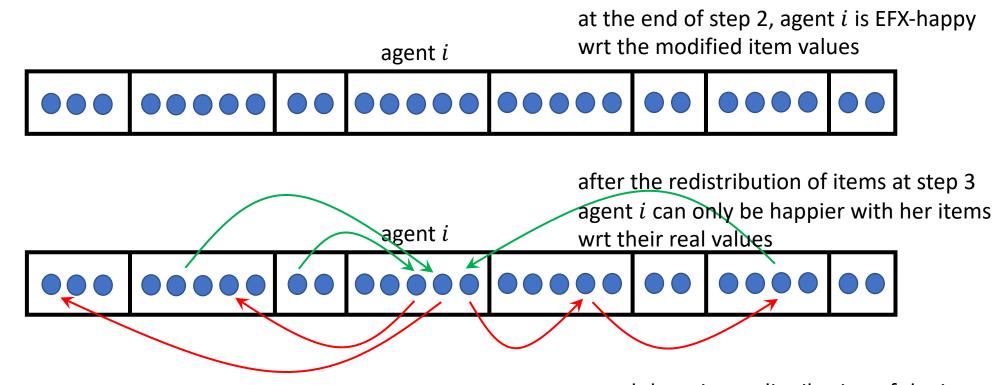
What happens at step 3?



What happens at step 3?



What happens at step 3?



... and there is a redistribution of the items (i.e., their assignment at the end of step 2) which makes agent *i* EFX-happy

Takeaway message

- EFX is still an important property and we should further explore it
- But why not focusing on alternative fairness concepts in parallel?
- In particular, on concepts that are related to it, like EEFX and MXS
- Reconsider existing algorithms (they may do more than we think)

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