



Building a Shared Sorting Function for a Group of Decision Makers

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Outline

- 1 Context and motivation
- 2 Why not the obvious solution?
- 3 Our framework
- 4 Improving the procedure
- 5 Conclusions & Future work

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Our goal

- Obtain a way of evaluating objects (*alternatives*),
- by **sorting** them into preference-ordered *categories*, e.g. {Good, Medium, Bad},
- on the basis of several (objective) performance measures (*criteria*).
- Resulting sorting function must be **consensual** among multiple Decision Makers (DMs).

General framework

- Alternatives \mathcal{A}
- Criteria \mathcal{J}
- Performances $g_j : \mathcal{A} \rightarrow X_j, \forall j \in \mathcal{J}$
- Preference orders $\succeq_j, \forall j \in J$ (total orders)
- Preference ordered set of categories \mathcal{C}
 - C_1 worst category
- Set of decision makers \mathcal{T}

	g_1	g_2	g_3	g_4		
a_1	5	6	1	+++	→ ?	C_3
a_2	4	7	8	+	→ ?	
a_3	2	8	4	-	→ ?	C_2
a_4	4	6	9	+	→ ?	
...						C_1

Example: student evaluation

- Determine a way of evaluating students at the end of the year
 - Obtain a sorting function
- \mathcal{A} , space of all possible evaluations (may be infinite or large)
- We may reason on specific cases (possibly fictitious)
- Involves subjective appreciations
 - DMs may have different opinions

	math.	lang.	phys.	partic.		
St 1	B	D	A	++	→ ?	C_4 : highest d.
St 2	A	A	B	+	→ ?	C_3 : distinction
St 3	A	C	C	-	→ ?	C_2 : succeeds
St 4	C	A	A	+	→ ?	C_1 : fails
...						

Example: research projects

- Researchers submit projects to the board
- The board wants some way of evaluating these research projects
- Criteria: Redaction quality, scientific quality, experience of the team, publication score, ...

	redac.	sci.	exp.	publ.		
Pr 1	2	5	3	75.1	→ ?	C ₂ : fund project
Pr 2	5	5	4	32.2	→ ?	
Pr 3	5	3	3	63.4	→ ?	
Pr 4	3	5	4	61.7	→ ?	
...						C ₁ : reject

Example: green labelling

- Evaluate ecological quality of consumer products
- Criteria: amount of pollutants of different sorts, road distance, ...
- Use a representative set of products
- Obtain a transparent decision procedure

pol. 1 pol. 2 rd dist.

Pr 1	2	5	3	→ ?
Pr 2	5	5	4	→ ?
Pr 3	5	3	3	→ ?
Pr 4	3	5	4	→ ?
...				

A+
A
B
C
D

Important features of the problem setting

Hypothesis

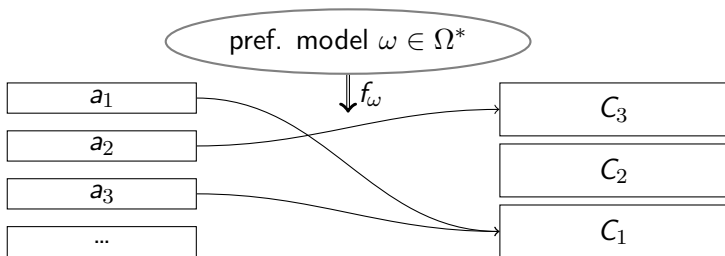
The situation is *not* about bargaining.

Sorting function depends on objective and **subjective** data.

- Objective, or consensual, data: most importantly, performances;
- subjective data: how the performances relate.

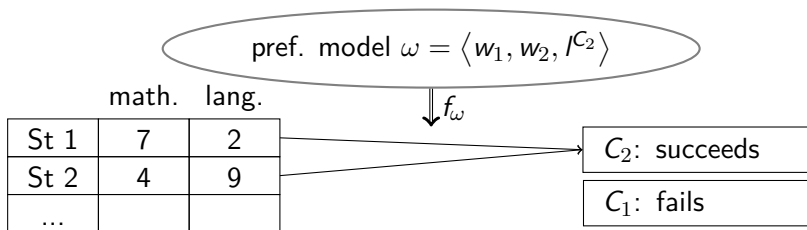
Modelling subjectivity

- Sorting function is parameterized
- Vector of parameters ω captures the subjective aspect of the sorting
- Set of possible parameter values is Ω^*
- Choosing Ω^* defines the set of possible sorting functions
- Can be done for one individual or for the group



Example: weighted sum with thresholds

- Criteria \mathcal{J} , $|\mathcal{J}| = n$, with functions $g_j : A \rightarrow \mathbb{R}$
- Preference model $\omega = \langle W, L \rangle$:
 - a vector of weights $W \in \mathbb{R}_+^n$,
 - a set of thresholds $L \in \mathbb{R}_+^{K-1}$ (K categories), l^C is the low threshold for category C .
- Class of models is $\Omega^* = \mathbb{R}_+^n \times \mathbb{R}_+^{K-1}$
- Sorting function $f_\omega : A \rightarrow \mathcal{C}$ compares the score of a to the thresholds l^{C_2}, l^{C_3}, \dots



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An obvious way

Two difficulties

- 1 Multiple criteria to aggregate
- 2 Multiple DMs whose points of view must be aggregated

A simple way to get around the second difficulty and come back to the mono-DM case:

- Come up with individual preference models $\omega_t, \forall t \in \mathcal{T}$ (typically using utility functions)
- The group preference model is some aggregation of the individual preference models;
- or the group sorting function is some aggregation of the individual sorting functions $\{f_{\omega_t}, t \in \mathcal{T}\}$.

Why not?

What we want

- Avoid unnecessary sacrifices
 - use of preference lability
- Achieve a better understanding of the points of consensus and disagreements
- Explore non utility-based classes of preference models
- Ask questions in terms of the problem
 - effects on the sorting results
- Ask easy questions
 - assignment examples

(Not all of these points are specific to the approach presented here.)

Preference lability

- Preferences are labile: not determined precisely in one's head
- DMs may not know their own preferences
- Maybe several equally good ways of aggregating the criteria, for a given DM
- Documented (with a different perspective) in numerous experiments
[Kahneman and Tversky, 2000, Lichtenstein and Slovic, 2006]
 - “Preferences are constructed in context”

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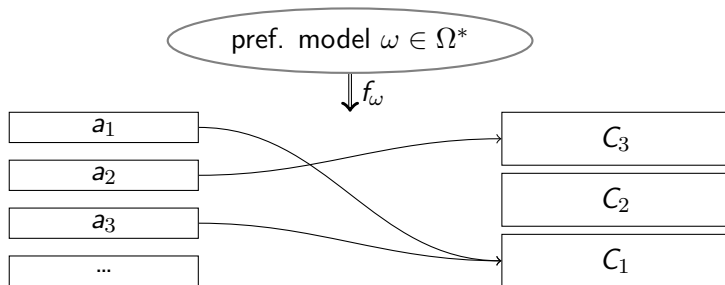
Sketch of the procedure

- We work with (a variant of) ELECTRE TRI as the class of sorting functions
- Parameters to be elicited: $\omega = \langle L, W, \lambda \rangle$
- We ask for assignment examples (e.g. $a_1 \rightarrow C_2$)
- These constrain the set of candidate preference models
 - We must have $f_\omega(a_1) = C_2$
- If no preference model satisfy all examples:
 - Search which constraints should be removed to restore consistency

Defining the group sorting function

Explanation proceeds in two steps.

- Show how (the variant of) ELECTRE TRI works: how f_ω is defined, *assuming* the preference model ω is defined
- Then, explain how we find a suitable preference model for the group of DMs.



Sorting method: a variant of ELECTRE TRI

Preference parameters

- *Category limits* $L = \langle I^C, C \in \mathcal{C} \setminus C_1 \rangle$: determine when the alternative is *good enough* on a criterion
- Weights $W = \langle w_j, j \in \mathcal{J} \rangle$, and a majority threshold λ : determine when the alternative is *globally good enough*

Alternatives \mathcal{A}

	g_1	g_2	g_3
a_1	3	1	1
a_2	5	3	3
a_3	0	5	1
a_4	2	0	2

Cat. limits L

	g_1	g_2	g_3
	C ₃ : Good		
I^C_3	4	4	3
	C ₂ : Average		
I^C_2	3	3	2
	C ₁ : Bad		

Weights W, λ

	g_1	g_2	g_3
W	0.2	0.6	0.2
$\lambda = 0.8$			

Sorting method: a variant of ELECTRE TRI

- a may reach at least C iff $\sum_{j \text{ in favor}} w_j \geq \lambda$ ($C \neq C_1$).
- j in favor of a reaching C iff $g_j(a) \geq I_j^C$.
- Thus, a sorted into the best category s.t.

$$\sum_{j|g_j(a) \geq I_j^C} w_j \geq \lambda.$$

Alternatives \mathcal{A}

	g_1	g_2	g_3
a_1	3	1	1
a_2	5	3	3
a_3	0	5	1
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Weights W, λ

	g_1	g_2	g_3
W	0.2	0.6	0.2
$\lambda = 0.8$			

Strengths of ELECTRE TRI

- A variant has been axiomatized [Bouyssou and Marchant, 2007a, Bouyssou and Marchant, 2007b].
- Justification of assignment easy to grasp (no complex computation needed).
- Might ease discussion among DMs.

Determining a group preference model

We want to determine a suitable $\omega = \langle L, W, \lambda \rangle$ for the group of DMs.

- We ask for assignment examples
- $A^* \subseteq \mathcal{A}$ the set of alternatives used as examples
- $\forall a \in A^*$, we know the category a should go into, according to one DM at least
- Assignment examples should be non contradictory

Mathematical program

Having assignment examples $E \subseteq A^* \times C$,

- We want to find $\omega = \langle L, W, \lambda \rangle$ such that
 - $\forall (a, C) \in E : f_\omega(a) = C$.
- Idea (finding ω satisfying examples) existed already [Mousseau and Słowiński, 1998] but no efficient tools to solve it.
- We solve a Mixed Integer Program (MIP) [Cailloux et al., 2012].
- The MIP must represent the assignment examples as constraints on the decision variables L, W, λ .

Defining (some of) the constraints: idea

Consider example $a \rightarrow C_k$, with $C_k \neq C_1, C_{|C|}$.

- Criterion j thinks that a deserves to reach C_k iff
 - $g_j(a) \geq f_j^{C_k}$.
- Introduce binary variable $b_j^{a, C_k} = 1$ iff a may reach C_k according to j .

Examples E

	g_1	g_2	g_3	C
a	3	1	5	C_k
...				

Cat. limits L

	g_1	g_2	g_3
	C_3 : Good		
f^{C_3}	$f_1^{C_3}$	$f_2^{C_3}$	$f_3^{C_3}$
	C_2 : Average		
f^{C_2}	$f_1^{C_2}$	$f_2^{C_2}$	$f_3^{C_2}$
	C_1 : Bad		

Weights W, λ

g_1	g_2	g_3
w_1	w_2	w_3

Defining (some of) the constraints: idea

Consider example $a \rightarrow C_k$, with $C_k \neq C_1, C_{|C|}$.

- a deserves to reach at least C_k iff

$$\sum_{j|b_j^{a,C_k}} w_j \geq \lambda.$$

Examples E

	g_1	g_2	g_3	C
a	3	1	5	C_k
...				

Cat. limits L

	g_1	g_2	g_3
	C_3 : Good		
I_{C_3}	$I_1^{C_3}$	$I_2^{C_3}$	$I_3^{C_3}$
	C_2 : Average		
I_{C_2}	$I_1^{C_2}$	$I_2^{C_2}$	$I_3^{C_2}$
	C_1 : Bad		

Weights W, λ

g_1	g_2	g_3
w_1	w_2	w_3

Defining (some of) the constraints

Consider example $a \rightarrow C_k$, with $C_k \neq C_1, C_{|C|}$.

We want that

- a reaches at least C_k ;
- a does not reach C_{k+1} :

$$\sum_{j|b_j^{a,C_k}} w_j \geq \lambda \wedge \sum_{j|b_j^{a,C_{k+1}}} w_j < \lambda.$$

Define continuous variable [Meyer et al., 2008]:

$$v_j^{a,C_k} = \begin{cases} w_j & \text{if } b_j^{a,C_k} = 1, \\ 0 & \text{otherwise.} \end{cases}$$

Therefore:

$$\sum_j v_j^{a,C_k} \geq \lambda \wedge \sum_j v_j^{a,C_{k+1}} < \lambda.$$

Recap of the procedure

- Obtain assignment examples
- Run the MIP
- Find a preference model $\omega = \langle L, W, \lambda \rangle$ such that f_ω satisfies the examples
- Present the results to the DMs by applying the function to a larger set of alternatives, or by explaining how it “reasons”
- They might want to correct or add examples

Also possible:

- No satisfying model exist
- Then some examples must be changed
- Existing procedures can find minimal sets of constraints to remove [Mousseau et al., 2006]

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Assignments stability

- Typically, DMs provide few examples compared to number of examples required to define a sorting function.
- Thus, great variability in possible assignments of other alternatives (from $\mathcal{A} \setminus A^*$).
- When one example changes (thus ω is changed to ω'), the other assignments may completely change (f_ω may be completely different than $f_{\omega'}$).
- This is called instability.
- As the procedure is used interactively, instability can occur at some point.
- Convergence may be slow and hard to see.

Improving the procedure

- Possible instability
- Possibly no consensus from the start
- What if starting examples are contradictory?
- We want the procedure to be more incremental
- First, agree on the category limits (L), then on the weights (W, λ).

Partially shared parameters

- We ask for examples E^t to each DM $t \in \mathcal{T}$.
- These can be contradictory, e.g. $(a_1 \xrightarrow{t_1} C_3), (a_1 \xrightarrow{t_2} C_2)$.
- For each DM $t \in \mathcal{T}$, we search for individual preference models $\omega^t = \langle L, W^t, \lambda^t \rangle$ satisfying examples E^t .
- Thus, category limits are shared but weights are chosen individually.
- This may exist even though there is no shared model satisfying all examples.
- This decomposes the problem into two simpler problems.
- Once shared category limits are found, better stability.
- Connects with existing procedures to find shared weights [Damart et al., 2007].

Implementation

- Search procedure may be implemented with a MIP.
- Has been implemented in a Java free (libre) package:
www.decision-deck.org/j-mcda/ [Cailloux, 2012].
- Available as a web service in the Decision Deck framework.
- Can be used through a client program (diviz):
<http://www.decision-deck.org/diviz/>.

Performance test

Objective

Examine whether the MIP is able to find shared profiles within a reasonable time frame.

- Random generation : [3–10] criteria, [1–4] DMs, [2–5] categories, [1–700] examples per DM.
- Random performances $g_j(a)$.
- Shared profiles used to sort examples, then forgotten.

Performances

- Nb binaries = $|\mathcal{J}| \times |A^*| \times (|\mathcal{C}| - 1)$.
- 6 criteria, 3 categories, 3 DMs giving each 30 different examples: ≈ 1000 binaries

Results

- Solved within 90 minutes using less than 3 GB disk space:

Binary variables	Sample size	Problems solved
[0, 399]	477	100%
[400, 799]	441	87%
[800, 1199]	362	80%
[1200, 1599]	290	78%
[1600, 1999]	268	75%
[2000, 2199]	121	69%

- Mainly depend on the number of criteria.

Computing restrictions on weights

- We can now propose shared L to the DMs.
- For each $t \in \mathcal{T}$, accepting L imposes a restriction on possibility of choosing weights W^t satisfying the examples E^t .
- We define an LP to compute these restrictions,
 $\forall t \in \mathcal{T}, j_1, j_2 \in \mathcal{J}$:

$$j_1 \triangleright^t j_2 \Leftrightarrow w_{j_1}^t > w_{j_2}^t, \forall \langle w_j^t, j \in \mathcal{J}, \lambda^t \rangle \text{ satisfying } E^t.$$

- This may help the DMs in choosing shared category limits.

Variants and extensions

- Possible to set direct constraints on the model parameters.
- Possible to search for models including ELECTRE TRI vetoes.
- Possible to specify constraints on the (weighted) category size [Zheng et al., 2011].
 - Select research projects that fit the budget.
 - Mainly for consensual constraints.
- DMs may give imprecise assignments ($a \xrightarrow{t} [C_1, C_2]$).
- Also implemented in software.

Possible use

This approach can be considered as a supplementary tool in the analyst's toolbox. Here is only one possible use.

- Ask for examples E^t .
- Resolve possible individual inconsistencies.
- Search for consensual model ω .
- If no such model, search for shared category limits.
- If still no model, allow for vetoes.
- Present resulting category limits and individual weights with restrictions on the weights.
- If not acceptable, DMs may provide supplementary examples or other constraints.

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Conclusion

- A path out of the “average of individual opinions” strategy.
- We search for consensus instead of compromise.
- We use a “divide and conquer” approach.
- Using a model possibly more intuitive than utility functions.
- Asks easy questions.
- Results are easily interpretable.
- Computation time: could be improved.

Future work

- Generalise the idea and separate the specifics to ELECTRE TRI.
- May also apply to other classes of models (such as utility functions)
- May apply to other problem types, e.g. ranking instead of sorting.
- Separate parameters in different manners?
- Formal description of the relation between “what we want” and “what we do”.
- Validation of the model class.

Thank you for your attention!



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max s s.t.

$$\left\{ \begin{array}{l} \sum_{j \in J} w_j = 1 \\ f_j^{C_h} \leq f_j^{C_{h+1}} \end{array} \right.$$

Variables

$$\left\{ \begin{array}{l} f_j^{C_h}, \forall j \in J, C_h \in \mathcal{C} \\ w_j, \forall j \in J \\ \lambda \\ b_j^{a, C_h} \text{ (binaries)} \\ v_j^{a, C_h} \\ n(a, C_h) \text{ (binaries)} \end{array} \right.$$

$$\left\{ \begin{array}{l} \frac{(g_j(a) - f_j^{C_h}) + \varepsilon}{M} \leq b_j^{a, h} \leq \frac{g_j(a) - f_j^{C_h}}{M} + 1 \\ v_j^{a, C_h} \leq w_j; b_j^{a, C_h} + w_j - 1 \leq v_j^{a, C_h} \leq b_j^{a, C_h} \\ \sum_{j \in J} v_j^{a, C_h} \geq \lambda + s \quad \forall a \rightarrow h, h \geq 2 \\ \sum_{j \in J} v_j^{a, C_{h+1}} + s \leq \lambda - \varepsilon \quad \forall a \rightarrow h, h < k \\ n(a, C_h) \leq 1 + \sum_{j \in J} v_j^{a, C_h} - \lambda \\ n(a, C_h) \leq 1 + \lambda - \sum_{j \in J} v_j^{a, C_{h+1}} - \varepsilon \\ \sum_{1 \leq h \leq k} n(a, C_h) = 1 \\ \underline{n}_h \leq \sum_{a \in A} n(a, C_h) P(a) \leq \bar{n}_h \quad \forall \langle C_h, P, \underline{n}_h, \bar{n}_h \rangle \end{array} \right.$$

All constraints (with qualifiers)

max s s.t.

$$\left\{ \begin{array}{l} \frac{(g_j(a) - f_j^{C_h}) + \varepsilon}{M} \leq b_j^{a,C_h} \leq \frac{g_j(a) - f_j^{C_h}}{M} + 1 \quad \forall j \in J, a \in A, h \geq 2 \\ v_j^{a,C_h} \leq w_j; b_j^{a,C_h} + w_j - 1 \leq v_j^{a,C_h} \leq b_j^{a,C_h} \quad \forall j \in J, a \in A, h \geq 2 \end{array} \right.$$

$$\left\{ \begin{array}{l} \sum_{j \in J} w_j = 1 \\ f_j^{C_h} \leq f_j^{C_{h+1}} \quad \forall j \in J, C_h \in \mathcal{C} \end{array} \right. \quad \left\{ \begin{array}{l} \sum_{j \in J} v_j^{a,C_h} \geq \lambda + s \quad \forall a \rightarrow h, h \geq 2 \\ \sum_{j \in J} v_j^{a,C_{h+1}} + s \leq \lambda - \varepsilon \quad \forall a \rightarrow h, h < k \end{array} \right.$$

Variables

$$\left\{ \begin{array}{l} f_j^{C_h}, \forall j \in J, C_h \in \mathcal{C} \\ w_j, \forall j \in J; \lambda \\ b_j^{a,C_h} \text{ (binaries)}; v_j^{a,C_h}; \\ n(a, C_h) \text{ (binaries)}, \\ \forall j \in J, a \in A, C_h \in \mathcal{C} \end{array} \right.$$

$$\left\{ \begin{array}{l} n(a, C_h) \leq 1 + \sum_{j \in J} v_j^{a,C_h} - \lambda \quad \forall a \in A, h \geq 2 \\ n(a, C_h) \leq 1 + \lambda - \sum_{j \in J} v_j^{a,C_{h+1}} - \varepsilon \quad \forall a \in A, h \leq k-1 \\ \sum_{1 \leq h \leq k} n(a, C_h) = 1 \quad \forall a \in A \\ \underline{n}_h \leq \sum_{a \in A} n(a, C_h) P(a) \leq \bar{n}_h \quad \forall \langle C_h, P, \underline{n}_h, \bar{n}_h \rangle \end{array} \right.$$