

# Computational Complexity

## Lecture 8: Some Sort of Recap

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## Recap

*What we saw last time..*

- The classes  $\Sigma_i^P$  and  $\Pi_i^P$
- The Polynomial Hierarchy
- $\Sigma_i^P$ -complete and  $\Pi_i^P$ -complete QBF problems
- Characterizations using oracles and ATMs

## What will we do today?

- Reflecting on what we've seen before
- Mostly using examples (of games)

## One-/two-liner overview of complexity classes

- L: deterministic algorithm, logarithmic space (and polynomial time)
- NL: nondeterministic algorithm, logarithmic space (and polynomial time)
- P: solvable in (deterministic) polynomial time
- NP: solutions (*for yes-answers*) can be guessed/checked in polynomial time
- coNP: solutions (*for no-answers*) can be guessed/checked in polynomial time
- $\Sigma_2^P$ : solutions (*for yes-answers*) have “ $\exists\forall$  structure”
- $\Pi_2^P$ : solutions (*for yes-answers*) have “ $\forall\exists$  structure”
- PSPACE: (non)deterministic algorithm, polynomial space (and exponential time)  
OR: unbounded “ $\exists\forall\exists\forall\exists\cdots$  structure”
- EXP: solvable in (deterministic) exponential time

## Some oracle questions..

- Is it the case that  $P^P = P$ ? Yes
- Is it the case that  $NP^{NP} = NP$ ? We don't know..
- Is it the case that  $PSPACE^{PSPACE} = PSPACE$ ? Yes
- Is it the case that  $EXP^{EXP} = EXP$ ? No
- Is it the case that  $DTIME(n^2)^{DTIME(n^2)} = DTIME(n^2)$ ? No
- Is it the case that  $NTIME(n)^{NTIME(n)} = NTIME(n)$ ? We don't know..

- Polls on  $P \stackrel{?}{=} NP$  have been held among computational complexity researchers:
  - In 2002, see: <https://tiny.cc/pnp-poll1>
  - In 2012, see: <https://tiny.cc/pnp-poll2>
  - In 2019, see: <https://tiny.cc/pnp-poll3>
- In these papers, there are some very interesting opinions on the question (and some nerdy jokes)
- Short answer: we have no clue (really), why  $P = NP$  or  $P \neq NP$  would be true, but most think that  $P \neq NP$ .

## Quiz example #1: checking if a given solution is unique

- What is the complexity of this problem?
- *Input:* A propositional formula  $\varphi$ , and a satisfying truth assignment  $\alpha$  for  $\varphi$ .  
*Question:* Is  $\alpha$  the only satisfying assignment for  $\varphi$ ?

## Quiz example #2: finding a minimal equivalent DNF formula

- What is the complexity of this problem?
- *Input:* A propositional formula  $\varphi$ , and  $1^k$  for some  $k \in \mathbb{N}$ .  
*Question:* Is there a DNF formula  $\psi$  of size  $\leq k$  such that  $\varphi \equiv \psi$ ?

## Quiz example #3: equivalence of propositional logic formulas

- What is the complexity of this problem?
- *Input:* Two propositional formulas  $\varphi_1, \varphi_2$ .  
*Question:*  $\varphi_1 \equiv \varphi_2$ ?

## Quiz example #4: 2SAT

- What is the complexity of this problem?
- *Input:* A propositional 2CNF formula  $\varphi$ .  
*Question:* Is  $\varphi$  satisfiable?

## Quiz example #5: satisfiability of modal logic K

- What is the complexity of this problem?
- *Input:* A basic modal logic formula  $\varphi$ .  
*Question:* Is  $\varphi$  satisfiable?

## Quiz example #6: satisfiability of modal logic S5

- What is the complexity of this problem?
- *Input:* A modal logic formula  $\varphi$ .  
*Question:* Is there an S5 Kripke model where  $\varphi$  is true?
- **Theorem:** if there is an S5 Kripke model where  $\varphi$  is true, then there exists an S5 Kripke model with at most  $|\varphi|$  states where  $\varphi$  is true.

## Quiz example #7: Tiling I

- What is the complexity of this problem?

- *Input:* A set of 4-sided tile types, and  $n, m \in \mathbb{N}$ .

*Question:* Can we use these tile types to fill an  $n \times m$  grid, so that

- (1) the outsides of the grid all have side  $s_0$ , and
- (2) neighboring tiles have matching sides?

- What is the complexity of this problem?

- *Input:* A set of 4-sided tile types, and  $n \in \mathbb{N}$ .

*Question:* Can we use these tile types to fill an  $n \times m$  grid, for some  $m \in \mathbb{N}$ , so that

- (1) the outsides of the grid all have side  $s_0$ , and
- (2) neighboring tiles have matching sides?

## Quiz example #9: Greedy Spiders

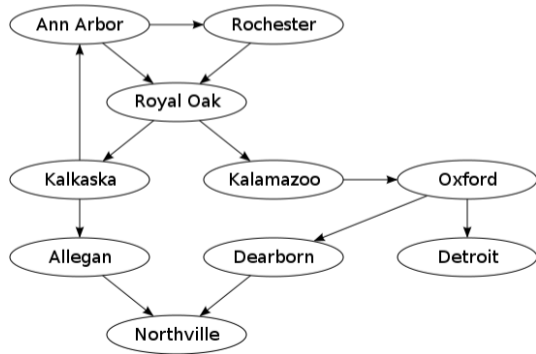
- What is the complexity of this problem?
- *Input:* An instance  $I$  of the *greedy spiders* puzzle.  
*Question:* Is there a solution for  $I$ ?

## Quiz example #10: Generalized Geography

- What is the complexity of this problem?  
(See: [https://en.wikipedia.org/wiki/Generalized\\_geography](https://en.wikipedia.org/wiki/Generalized_geography))

- *Input:* An instance  $I$  of *generalized geography*.

*Question:* Does Player 1 have a winning strategy?

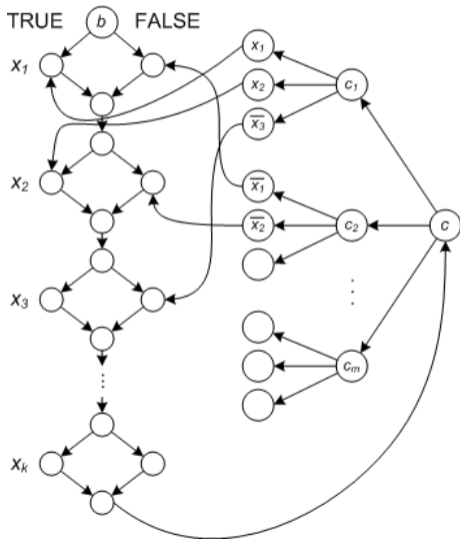


## Quiz example #10: Generalized Geography

- What is the complexity of this problem?  
(See: [https://en.wikipedia.org/wiki/Generalized\\_geography](https://en.wikipedia.org/wiki/Generalized_geography))

- Input:* An instance  $I$  of *generalized geography*.

*Question:* Does Player 1 have a winning strategy?



## Quiz example #11: Game of the Amazons (2 players)

- What is the complexity of this problem?  
(See: [https://en.wikipedia.org/wiki/Game\\_of\\_the\\_Amazons](https://en.wikipedia.org/wiki/Game_of_the_Amazons))
- *Input:* A (2-player) Game of the Amazons position (on an  $n \times n$  board).  
*Question:* Does Player 1 have a winning strategy?

## Quiz example #12: Game of the Amazons (1 player)

- What is the complexity of this problem?

(See: [https://en.wikipedia.org/wiki/Game\\_of\\_the\\_Amazons](https://en.wikipedia.org/wiki/Game_of_the_Amazons))

- *Input:* A (1-player) Game of the Amazons position (on an  $n \times n$  board), and  $1^k$  for some  $k \in \mathbb{N}$ .

*Question:* Can Player 1 make at least  $k$  consecutive moves?

## Quiz example #13: reachability in succinctly represented graphs

- What is the complexity of this problem?
- *Input:* A propositional logic formula  $\varphi(x_1, \dots, x_n, x'_1, \dots, x'_n)$ , and two binary vectors  $s, t \in \{0, 1\}^n$ .

*Question:* Consider the directed graph  $G = (V, E)$ , where:  
 $V = \{0, 1\}^n$ , and for each  $\bar{v}, \bar{w} \in V$ ,  
 $(\bar{v}, \bar{w}) \in E$  if and only if  $\varphi[\bar{v}, \bar{w}]$  is true.

Is  $t$  reachable from  $s$  in  $G$ ?

## Quiz example #14: 3-colorability for succinctly represented graphs

- What is the complexity of this problem?
- *Input:* A propositional logic formula  $\varphi(x_1, \dots, x_n, x'_1, \dots, x'_n)$ , and two binary vectors  $s, t \in \{0, 1\}^n$ .

*Question:* Consider the undirected graph  $G = (V, E)$ , where:  
 $V = \{0, 1\}^n$ , and for each  $\bar{v}, \bar{w} \in V$ ,  
 $\{\bar{v}, \bar{w}\} \in E$  if and only if  $\varphi[\bar{v}, \bar{w}]$  is true.

Is  $G$  3-colorable?

- Non-uniform complexity
- Circuit complexity
- TMs that take advice
- The Karp-Lipton Theorem





