

# Proportionality in Multiwinner Voting with Weighted Seats

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- ★ **Multiwinner Voting:** A job panel must produce a shortlist of  $k$  candidates to continue to the next interview stage.
- ★ **Participatory Budgeting:** Citizens must decide on the public projects, each coming with a cost, that are to be implemented by the local municipality, subject to a budget.

We look *another* such complex domain.

# Talk Outline

- ★ Standard Multwinner Voting (MWV) Model
- ★ Proportionality in MWV.
- ★ MWV with Weighted Seats.

## (Approval-based) MWV Model

- ★ Candidates  $C = \{a, b, c, \dots\}$ .
- ★ Agents  $N = \{1, \dots, n\}$ .
- ★ Each agent submits an *approval ballot*  $A_i \subseteq C$ .
- ★ Outcome is a committee  $W \subseteq C$  of size  $k$ .

## Definition ( $\ell$ -cohesiveness)

For an integer  $\ell \in \{1, \dots, k\}$ , a group of agents  $N' \subseteq N$  is  $\ell$ -cohesive if  $|N'| \geq n \cdot \frac{\ell}{k}$  and  $|\bigcap_{i \in N'} A_i| \geq \ell$ .

## Example

- Candidates  $C = \{a, b, c, d\}$  with  $k = 3$ .
- Agents  $N = \{1, 2, 3\}$ .
- Approval ballots are  $A_1 = \{a, b\}$ ,  $A_2 = \{a, b, c\}$  and  $A_3 = \{c, d\}$ .
- $\{1, 2\}$  is 2-cohesive.
- $\{2, 3\}$  and  $\{3\}$  are 1-cohesive.

# Proportionality in MWV

**Natural axiom:** if a group is  $\ell$ -cohesive then  $\ell$  of their common candidates should be elected to the committee.

## Definition (Strong Justified Representation (SJR))

A committee  $W$  provides SJR if for every  $\ell$ -cohesive group  $N'$ , it holds that  $|W \cap \bigcap_{i \in N'} A_i| \geq \ell$ .

However, this requirement is too strong, even when  $\ell = 1$ .

## Example

- Candidates  $C = \{a, b, c, d\}$  with  $k = 3$ .
- Agents  $N = \{1, \dots, 9\}$ .
- Suppose 2 agents approve  $\{a\}$ , another 2 agents approve  $\{d\}$ , and 1 agent each approves of  $\{b\}$ ,  $\{c\}$ ,  $\{a, b\}$ ,  $\{b, c\}$ ,  $\{c, d\}$ .
- Each candidate  $c \in \{a, b, c, d\}$  must be elected to provide SJR.

**A weaker axiom:** if a group is  $\ell$ -cohesive then at least one group member should be represented by  $\ell$  committee members.

## Definition (Extended Justified Representation (EJR))

A committee  $W$  provides EJR if for every  $\ell$ -cohesive group  $N'$ , there exists an agent  $i \in N'$  such that  $|W \cap A_i| \geq \ell$ .

## Multiwinner Voting with Weighted Seats

Joint work with Ulle Endriss, Ronald de Haan, Adrian Haret and Jan Maly.



# MWV with Weighted Seats

## Example

Each seat represents a role and some roles are more valuable than others.

- The committee has 5 seats with the following roles:  
(chair, treasurer, secretary, member, member).

## Example

Each seat has an associated budget that is available for the seat's elected candidate to spend.

- The committee has 5 seats with the following budgets:  
(\$3278, \$1400, \$560, \$100, \$4).

# Model

- ★ Candidates  $C = \{a, b, c, \dots\}$ .
- ★ Agents  $N = \{1, \dots, n\}$ .
- ★ Each agent submits an approval ballot  $A_i \subseteq C$ .
- ★ A weight vector  $\mathbf{w} = (w_1, \dots, w_k)$  with a weight for each of the  $k$  seats.
- ★  $W$  is the sum of all the weights.
- ★ Outcome is a committee  $\mathbf{c} = (c_1, \dots, c_k)$ .
- ★ For any set of candidates  $A \subseteq C$ , the satisfaction from a committee  $\mathbf{c}$  is  $\text{sat}(\mathbf{A}, \mathbf{c}) = \sum_{j=1}^k \mathbb{1}_{c_j \in A} \cdot w_j$ .

# Proportionality

For weight vector  $\mathbf{w}$ , the set of all *possible* satisfaction values is  $\text{SAT}(\mathbf{w})$ .

## Example

If  $\mathbf{w} = (5, 3, 1)$ , then  $\text{SAT}(\mathbf{w}) = \{1, 3, 4, 5, 6, 8, 9\}$ .

## Definition ( $\ell$ -WS-cohesiveness)

For an integer  $\ell \in \text{SAT}(\mathbf{w})$ , a group of agents  $N'$  is  $\ell$ -WS-cohesive if  $|N'| \geq n \cdot \frac{\ell}{W}$  and there exists a  $C' \subseteq \bigcap_{i \in N'} A_i$  with  $|C'| = t$  such that there exists a committee  $\mathbf{c}$  where  $\text{sat}(C', \mathbf{c}) \geq \ell$ , and  $|N'| \geq n \cdot \frac{t}{k}$ .

## Definition ( $\ell$ -WSJR)

A committee  $\mathbf{c}$  provides  $\ell$ -WSJR if for every  $\ell$ -WS-cohesive group  $N'$ , there exists an agent  $i \in N'$  such that  $\text{sat}(A_i, \mathbf{c}) \geq \ell$ .

Unfortunately,  $\ell$ -WSJR is not always satisfiable.

## Example

- Candidates  $C = \{a, b, c\}$ .
- Agents  $N = \{1, 2, 3\}$ .
- Weight vector  $\mathbf{w} = (3, 2, 1)$ .
- Approval ballots are  $A_1 = \{a\}$ ,  $A_2 = \{b\}$  and  $A_3 = \{c\}$ .

Also, even if such a committee exists, it is computationally hard to compute it.

**What now?** Weaken the axiom.

# Weakening $\ell$ -WSJR: Part 1

**Intuition:** some cohesive group member is just one ‘swap’ away from the deserved satisfaction?

$I_{\mathbf{c}}(A)$  is the vector of positions within the committee  $\mathbf{c}$  of candidates in  $A$ .

## Definition ( $\ell$ -WSJR-1)

A committee  $\mathbf{c}$  provides  $\ell$ -WSJR-1 if for every  $\ell$ -WS-cohesive group  $N'$ , there exists an agent  $i \in N'$  and some  $j \in I_{\mathbf{c}}(C \setminus A_i)$  such that either (i), we have  $w_j + \text{sat}(A_i, \mathbf{c}) \geq \ell$  if there exists some candidate  $c \in A_i$  with  $c \notin \mathbf{c}$ , or (ii), for some  $h \in I_{\mathbf{c}}(A_i)$ , it holds that  $w_j + \text{sat}(A_i, \mathbf{c}) - w_h \geq \ell$ .

Can  $\ell$ -WSJR-1 always be satisfied?

Inspired by the Method of Equal Shares (MES) rule in standard MWV.

The rule works in  $k$  rounds where agents pay to assign candidates to weights from  $\mathbf{w} = (w_1, \dots, w_k)$ :

- ★ In round  $r \in \{1, \dots, k\}$ , agents consider assignments to weight  $w_r$ .
- ★  $b_i(r)$  is agent  $i$ 's budget to start round  $r$ , and in round 1, we set  $b_i(1) = \frac{W}{n}$ .
- ★ In round  $r$ , we say a pair  $(c, w_r)$  is  $q$ -affordable for some  $q \in \mathbb{R}_{\geq 0}$ , with  $c$  currently unelected, if:

$$\sum_{i \in N: c \in A_i} \min(q, b_i(r)) \geq w_r.$$

- ★ If no pair is  $q$ -affordable then go to the next round, otherwise, for a  $q$ -affordable pair  $(c, w_r)$  for a minimum  $q$ , assign  $c$  to  $w_r$  and continue to the next round.

**Good news** in the following *restricted setting*.

**Party-list elections:** An election where for every pair of agents  $i, j \in N$ , it holds that either  $A_i = A_j$ , or  $A_i \cap A_j = \emptyset$ , and for every agent  $i$ , we have  $|A_i| \geq k$ .

## Theorem

*w*-MES satisfies  $\ell$ -WSJR-1 on party-list elections.

## Weakening $l$ -WSJR: Part 2

Use **LOWSAT**( $\mathbf{w}$ ) =  $(l_1, l_2, \dots, l_k)$  where  $l_t = \sum_{j=1}^t w_j$ .

### Example

If  $\mathbf{w} = (5, 3, 3, 1)$ , then **LOWSAT**( $\mathbf{w}$ ) =  $(1, 4, 7, 12)$ .

### Definition (Lower $l$ -WS-cohesiveness)

For an integer  $l \in \mathbf{LOWSAT}(\mathbf{w})$ , a group of agents  $N'$  is *lower  $l$ -WS-cohesive* if  $|N'| \geq n \cdot \frac{l}{W}$  and there exists a  $C' \subseteq \bigcap_{i \in N'} A_i$  with  $|C'| = t$  such that there exists a committee  $\mathbf{c}$  where  $\text{sat}(C', \mathbf{c}) \geq l$ , and  $|N'| \geq n \cdot \frac{t}{k}$ .

### Definition (Lower $l$ -WSJR)

A committee  $\mathbf{c}$  provides *lower  $l$ -WSJR* if for every *lower  $l$ -WS-cohesive* group  $N'$ , there exists an agent  $i \in N'$  such that  $\text{sat}(A_i, \mathbf{c}) \geq l$ .



**Bad news!**  $w$ -MES does not satisfy *lower*  $\ell$ -WSJR.

Is *lower*  $\ell$ -WSJR is always satisfiable? Yes, use MES as in standard MWV.

- ★ Treat all seats as having weight 1.
- ★ Run MES where each agent  $i$  has initial budget  $b_i(1) = \frac{k}{n}$  instead of  $\frac{W}{n}$ .
- ★ When a seat is bought for a candidate  $c$ , assign  $c$  to some weight.
- ★ MES ensures that cohesive groups get the seats that they deserve.

# Future Work

- ★ Test more rules.
- ★ Define other fairness notions.
- ★ More axioms for the setting.