



Filippos Vogiatzian

Quantum Key Distribution

Quantum Cryptography

Motivation

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- Backwards unreliable secrecy

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- As soon as we have quantum computers:
 - Shor's Algorithm: Integer factorisation in polynomial time!

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One-Time Pad !!!

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When Classical Cryptography fails...

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Quantum Cryptography

Post-quantum Cryptography

When Classical Cryptography fails...

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- Perfect Security
- Relies only on the laws of nature

Post-quantum Cryptography

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Post-quantum Cryptography

- Quantum Computational Security
- Relies on primitives that are equally hard for classical and quantum computers to solve
 - Lattice-Based Cryptography

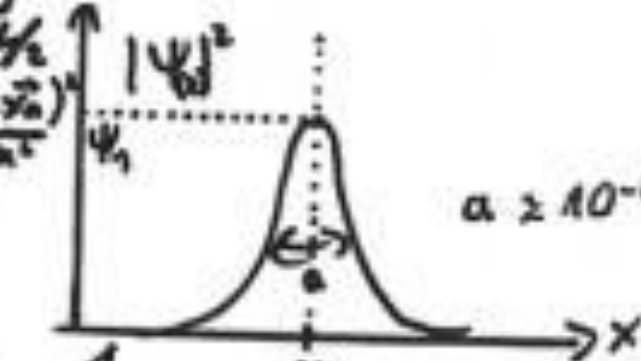
QKD 2PC and more

- Quantum Key Distribution (QKD): Two parties (Alice and Bob) communicate with perfect secrecy in the presence of an eavesdropper (Eve) [1,3,4]
- Two-Party Cooperation (2PC): Two parties that don't trust each other cooperate in a secure way
- How to encrypt or authenticate a quantum state
- Implementations

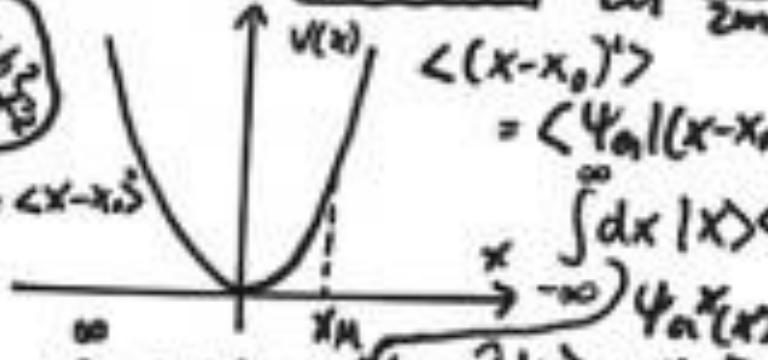
$$\langle \phi_n | \phi_{n'} \rangle = \langle \phi_n | \int_{-l/2}^{+l/2} dx |x\rangle \langle x| \phi_{n'} \rangle \Rightarrow \left(\frac{2\pi}{L}\right)^{1/2} \cos\left[\frac{\pi}{L}(2n-1)x\right] \quad \text{for } l = 2L \quad \text{for } l = 2L \quad \text{for } l = 2L$$

$$\langle \phi_n | \phi_{n'} \rangle = \int_{-l/2}^{+l/2} dx \phi_n^*(x) \phi_{n'}(x) \quad \Psi_n(x) = \sqrt{\frac{2}{L}} \cos\left[\frac{\pi}{L}(2n-1)x\right]; \quad \Psi_n(x) = \sqrt{\frac{2}{L}} \sin\left[\frac{2\pi}{L}nx\right]$$

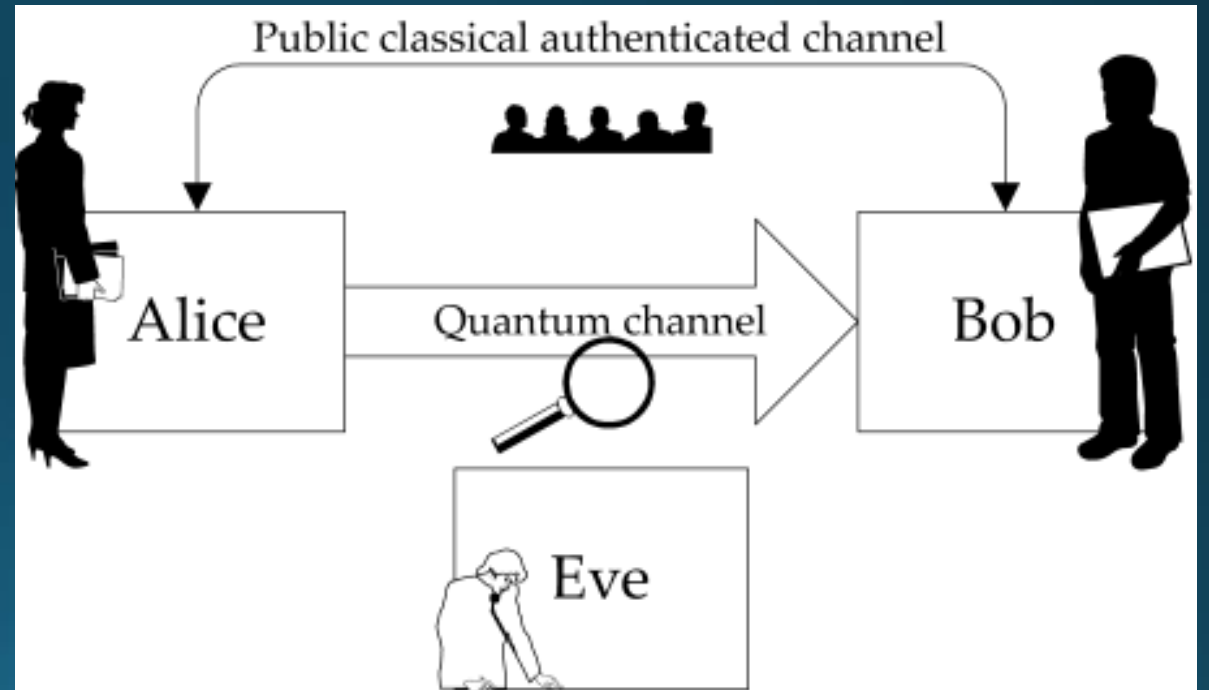
$$\langle \phi_n | \phi_{n'} \rangle = \frac{1}{L} \int_{-l/2}^{+l/2} dx e^{-ikx} e^{ik'x} \stackrel{!}{=} 0; \quad \hat{H} \Psi_{ns}(x) = -\frac{\hbar^2}{2m} \partial_x^2 \Psi_{ns}(x) = \frac{\hbar^2}{2m} \left(\frac{\pi}{L}(2n-1)\right)^2 \Psi_{ns}(x)$$

$|\Psi(x)|^2 = |\Psi_0|^2 e^{-\frac{(x-x_0)^2}{2a^2}}$

 $\int_{-\infty}^{+\infty} dx e^{-Ax^2} = \sqrt{\frac{\pi}{A}}$
 $A = \frac{1}{2a^2} \Rightarrow |\Psi_0| = \frac{1}{(2\pi a^2)^{1/4}}$
 $a \geq 10^{-10} \text{ m}$
 $\hat{H} \Psi_a = -\frac{\hbar^2}{2m} \partial_x^2 \Psi_a(x) = \frac{\hbar^2}{2m} \frac{1}{2a^2} \Psi_a(x) - \frac{\hbar^2}{2m} \frac{1}{4a^4} (x-x_0)^2 \Psi_a(x)$
 $= -\frac{\hbar^2}{2m} \left(-\frac{1}{2a^2} + \frac{(x-x_0)^2}{4a^4}\right) \Psi_a(x) \quad ; \quad V(x) = \frac{\hbar^2}{2m} \frac{1}{4a^4} (x-x_0)^2$
 $\hat{H} \rightarrow \hat{H} = -\frac{\hbar^2}{2m} \partial_x^2 + V(x); \quad \hat{H} \Psi_a = \frac{\hbar^2}{2m} \frac{1}{2a^2} \Psi_a = E_a \Psi_a$
 $V(x) = \frac{1}{2} m \omega^2 (x-x_0)^2 \rightarrow m \omega^2 = \frac{\hbar^2}{m 4a^4} \Rightarrow \omega = \frac{\hbar}{2ma^2}$
 $E_0 = \frac{\hbar^2}{2m} \frac{1}{2a^2}$

$[\hat{p}, \hat{x}] = \frac{\hbar}{i}; \quad \hat{p} = \frac{\hbar}{i} \partial_x \quad ; \quad \hat{H} = \frac{\hat{p}^2}{2m} + \frac{1}{2} m \omega^2 \hat{x}^2$
 $1. \quad a^2 + b^2 = (a+ib)(a-ib); \quad a, b \in \mathbb{R}; \quad 2. \quad (a\hat{p} + ib\hat{x})(a\hat{p} - ib\hat{x}), \quad a, b \in \mathbb{R}$
 $= a^2 \hat{p}^2 + iba\hat{x}\hat{p} - iab\hat{p}\hat{x} + b^2 \hat{x}^2 = a^2 \hat{p}^2 + b^2 \hat{x}^2 - \hbar ab$
 $\hat{H} = (a\hat{p} + ib\hat{x})(a\hat{p} - ib\hat{x}) = \hbar ab; \quad a^2 = \frac{1}{2m}; \quad b^2 = \frac{1}{2} m \omega^2$
 $\text{Def: } C^+ = \frac{1}{\sqrt{\hbar \omega}} (a\hat{p} + ib\hat{x}); \quad C^- = \frac{1}{\sqrt{\hbar \omega}} (a\hat{p} - ib\hat{x}) \Rightarrow \hat{H} = \hbar \omega C^+ C^- + \frac{1}{2} \hbar \omega$
 $\left(\frac{\omega}{\omega_0} \in \mathbb{C}\right) \quad \left\{ \begin{matrix} +1 \\ -1 \end{matrix} \right\} \text{SU}(2) \cong S^3 \quad A \rightarrow \omega \bar{A} \omega^{-1}$

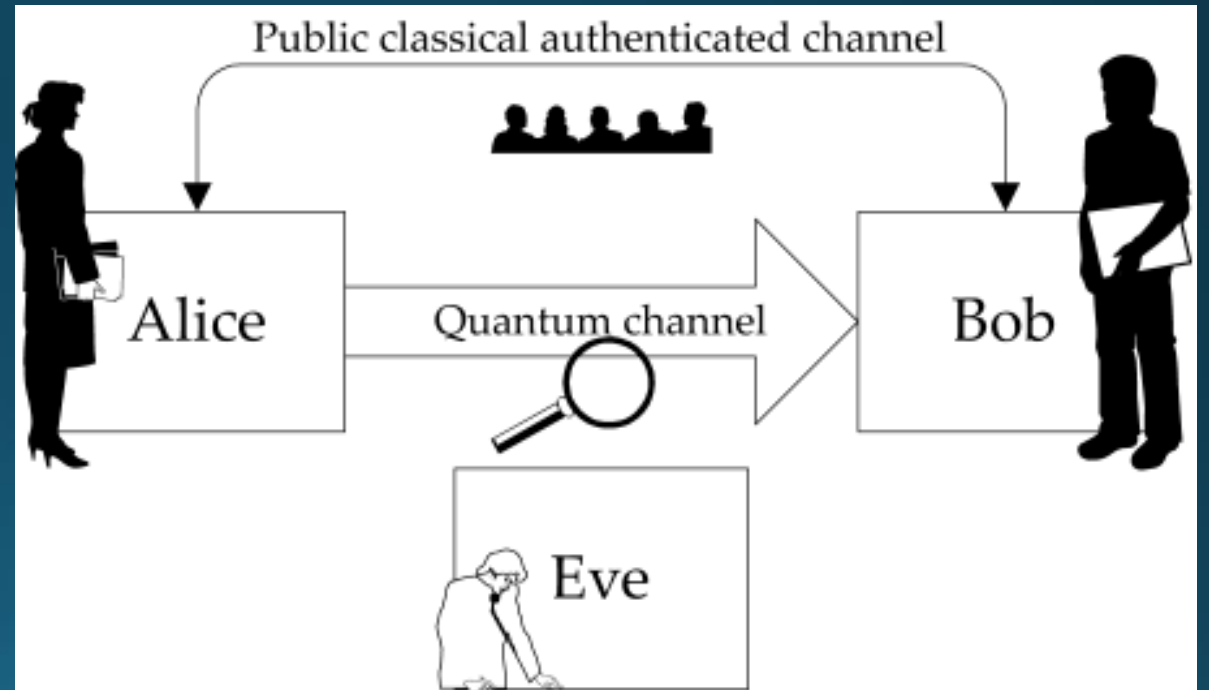

 $\langle (x-x_0)^n \rangle = \langle \Psi_0 | (x-x_0)^n | \Psi_0 \rangle = \int_{-\infty}^{+\infty} dx \Psi_0^*(x) (x-x_0)^n \Psi_0(x)$

Quantum Channel



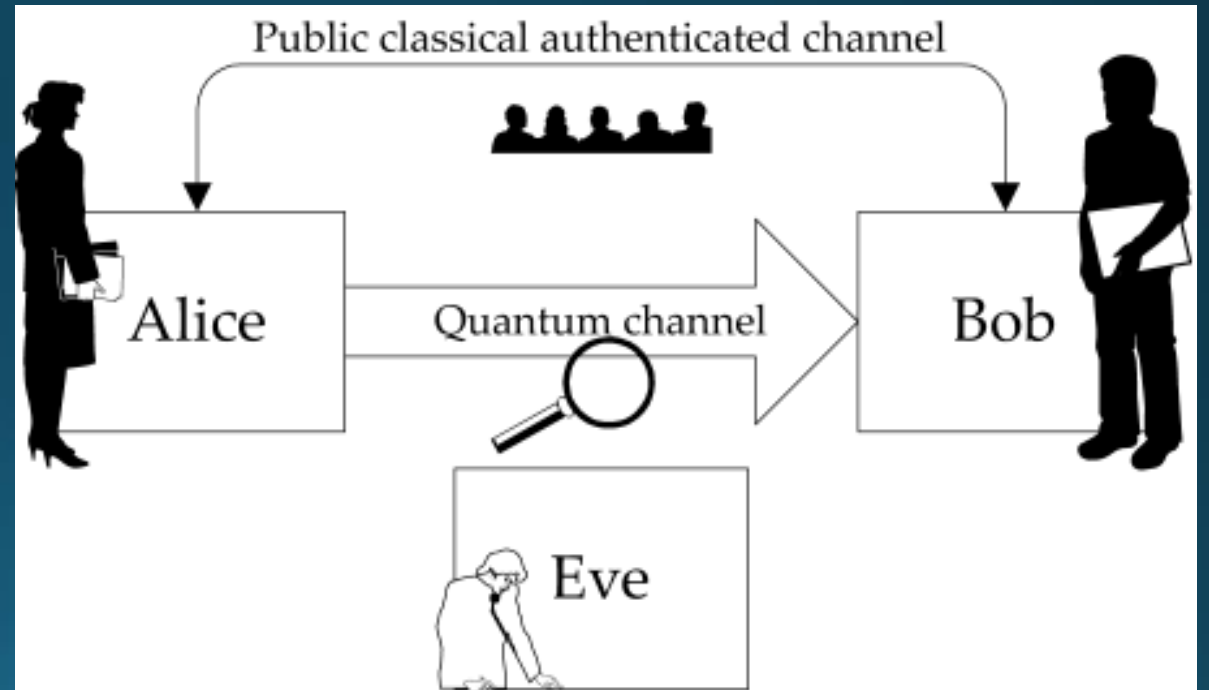
Quantum Channel

- Eve has complete control over the channel



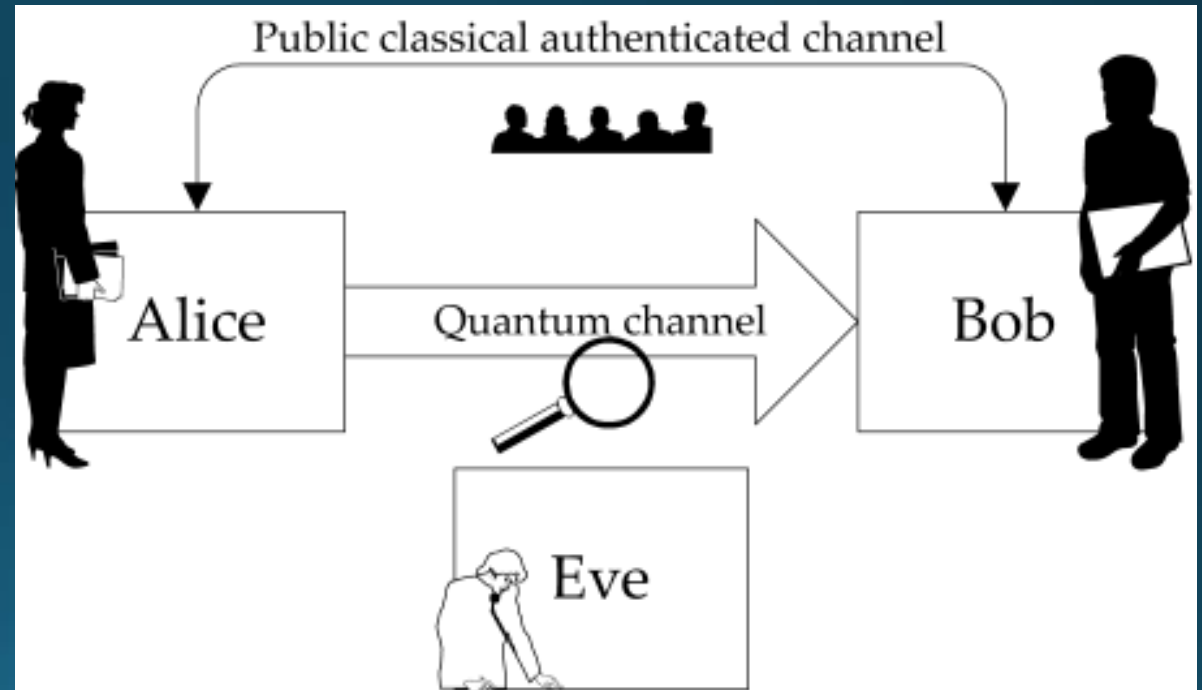
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- She can intercept or measure the sent qubits



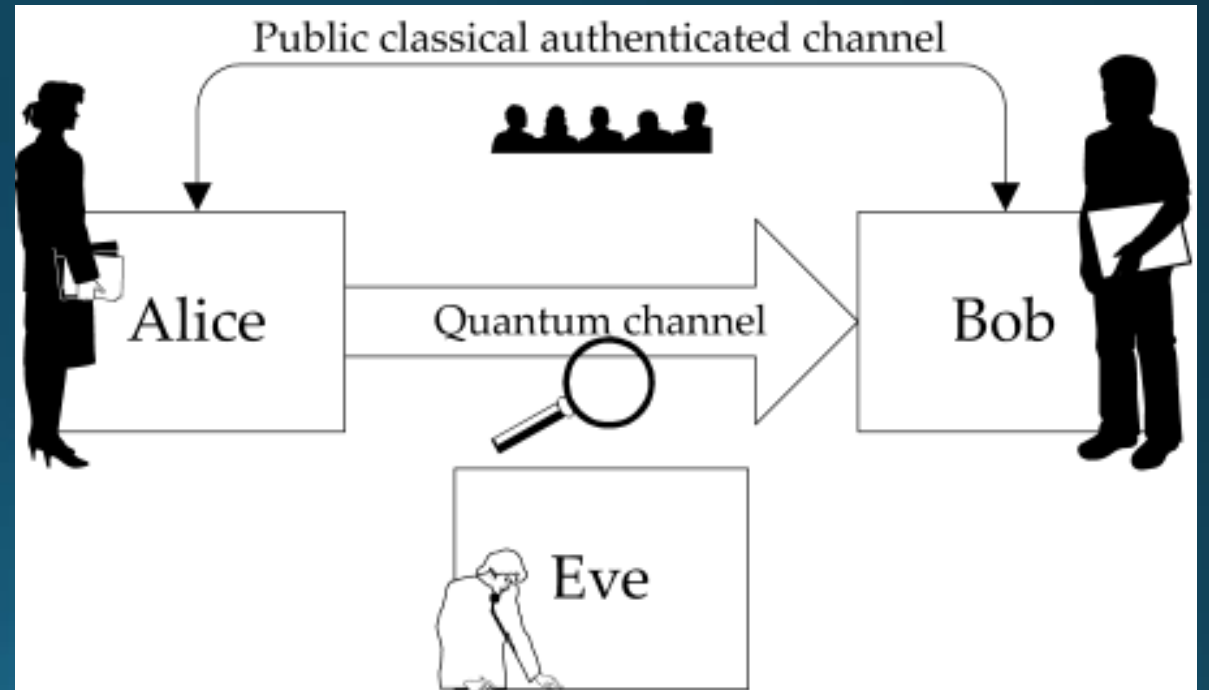
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- She can intercept or measure the sent qubits
- No cloning theorem: Forbidden to clone of an unknown quantum state (Wooters and Zurek and Dieks 1982)
- Eve can block the channel by sending random qubits and prevent communication over the channel



Quantum Key Distribution

- Alice and Bob use the public quantum channel to agree on a private secure key
- Eve has no information about the key
- Having a private key they can use any other classical encryption scheme to communicate through the public classical channel
- If they use OTP they can communicate with perfect secrecy!

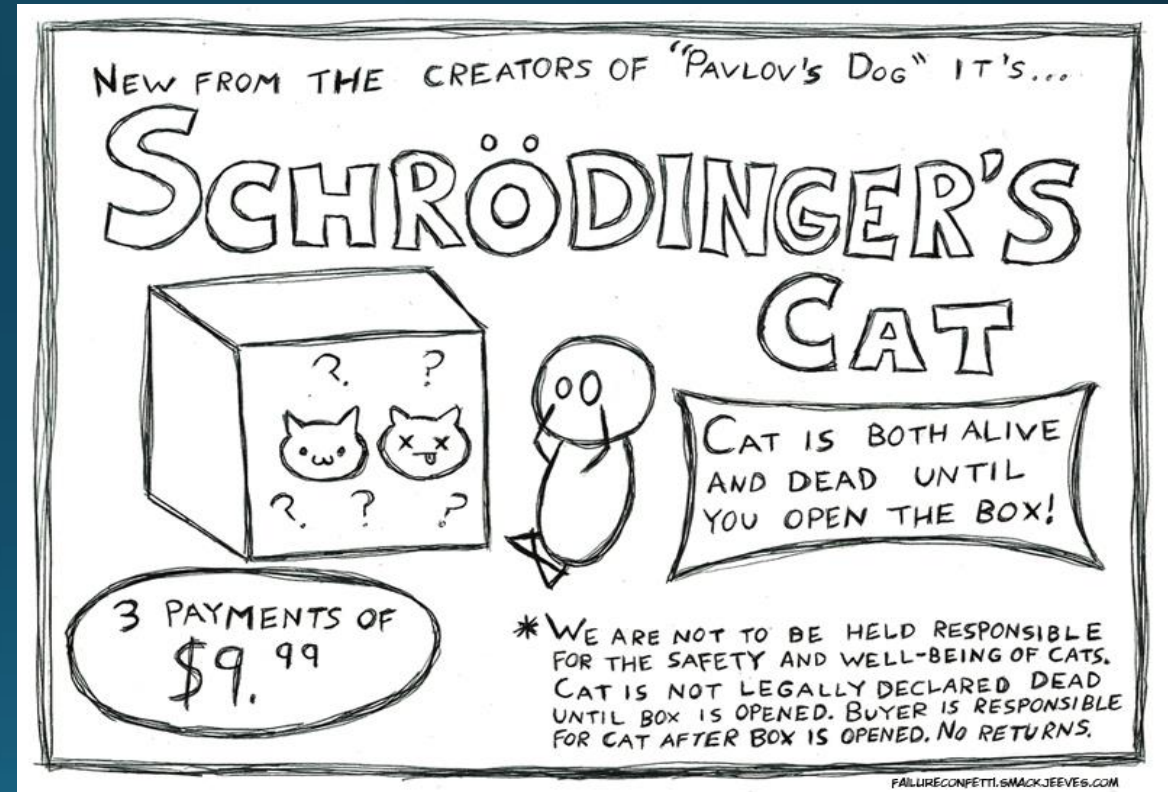
Quantum Bits (qubits)

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- Two-state quantum-mechanical system

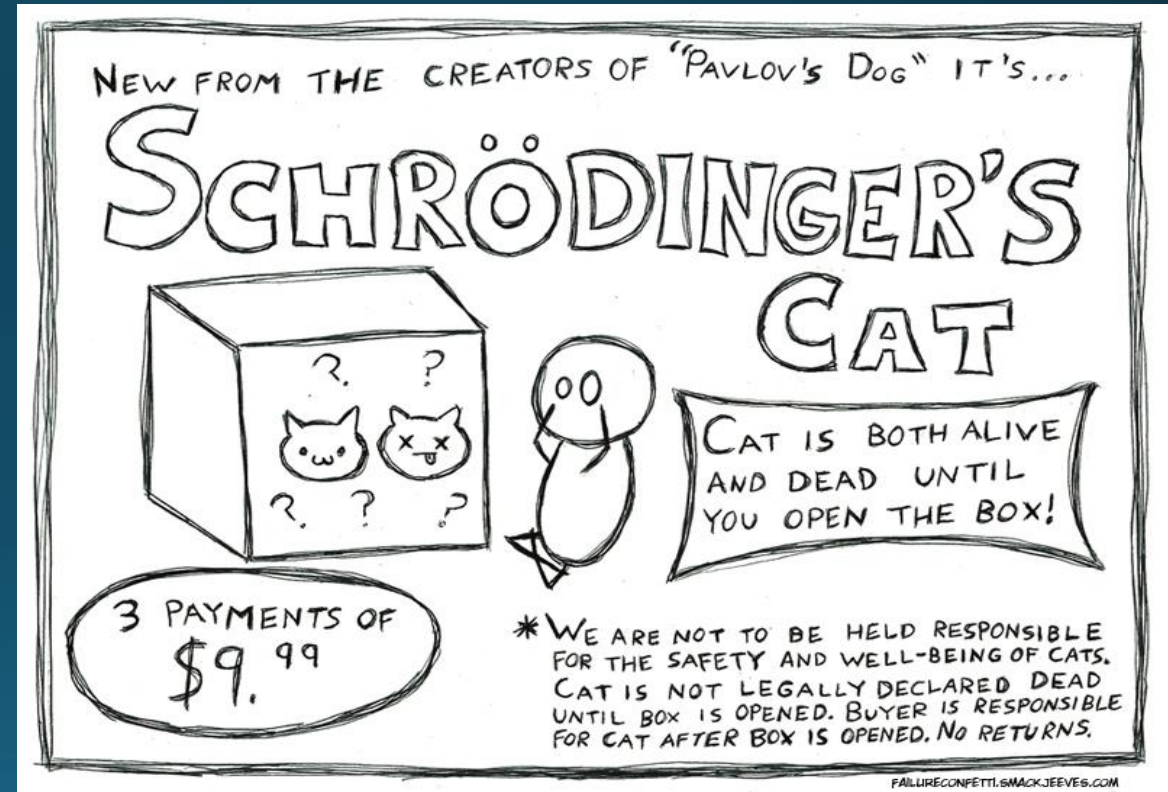
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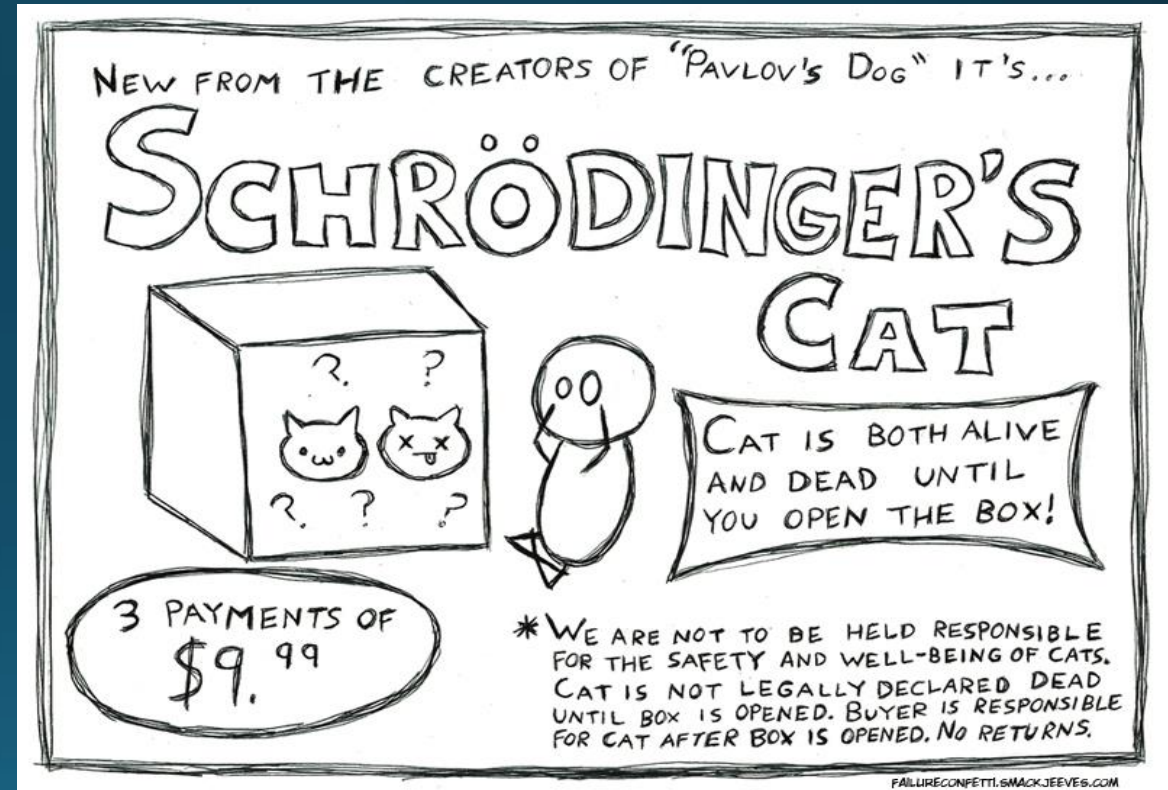
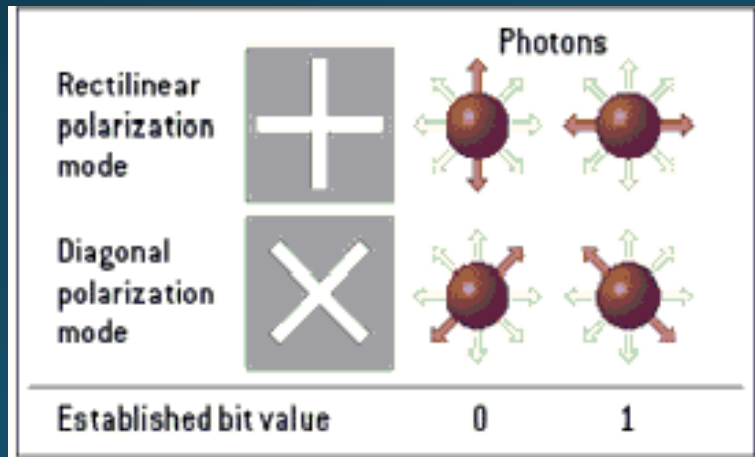
Quantum Bits (qubits)

- Two-state quantum-mechanical system
- Polarisation of photon
 - rectilinear / diagonal polarization

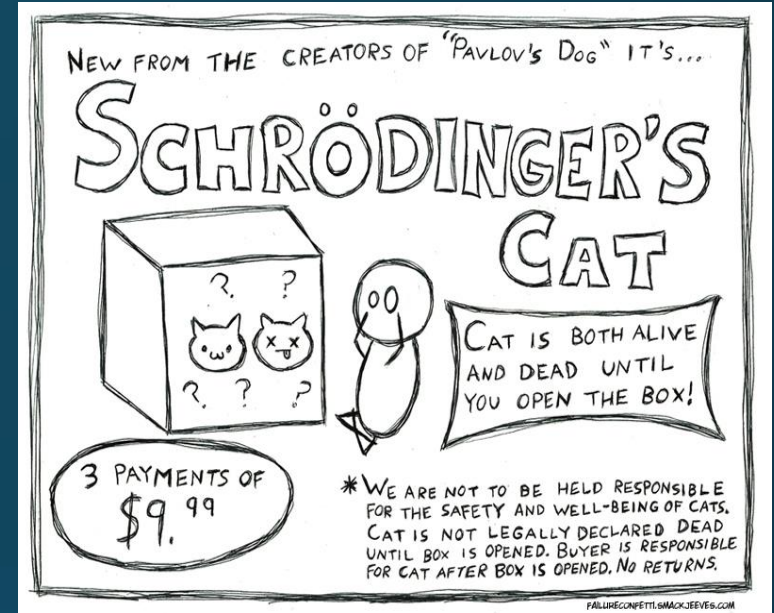


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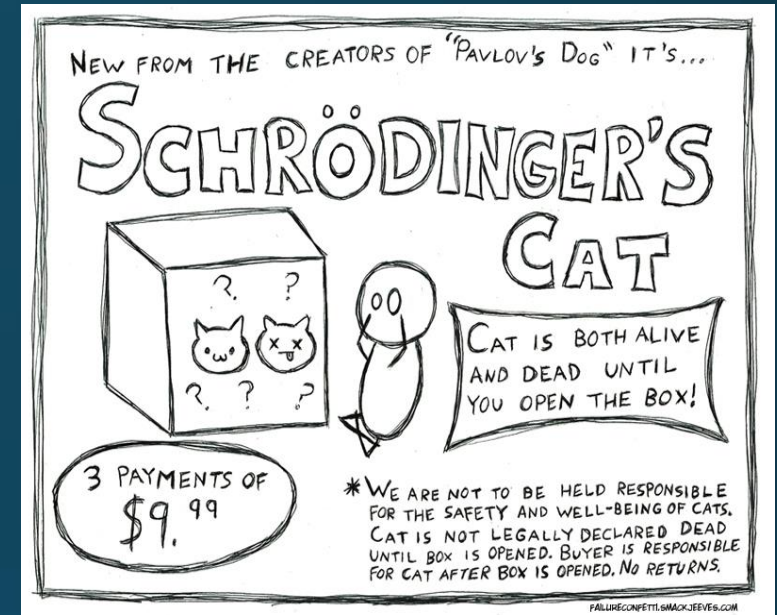


Measuring Qubits



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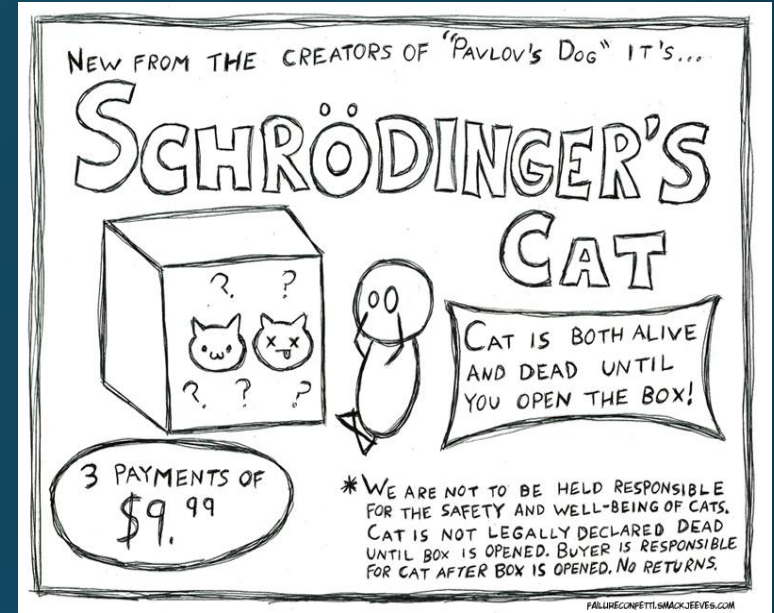
- Measuring a qubit:
 - opening the box



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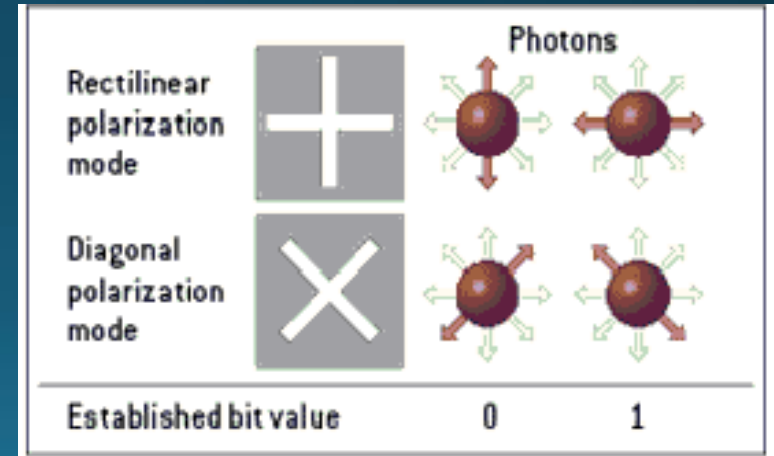
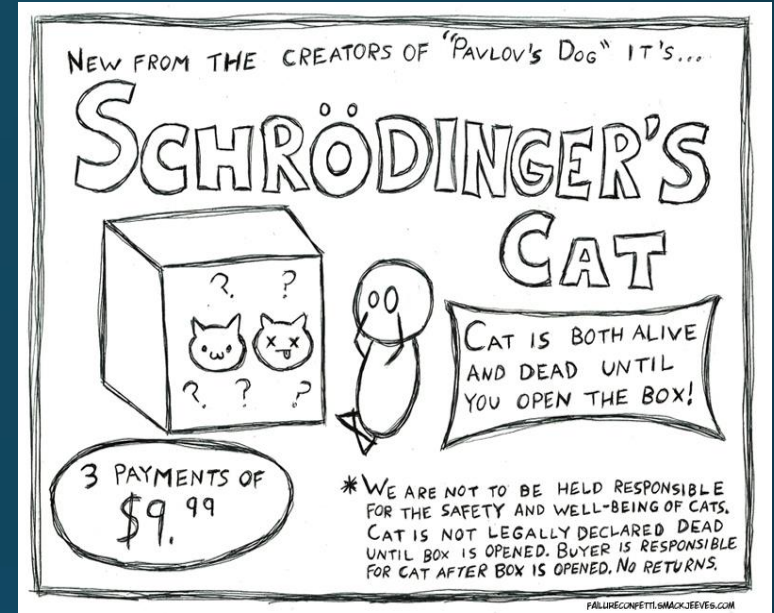
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- Collapses the wavefunction to one of the two states:
 - The cat is DEAD or ALIVE



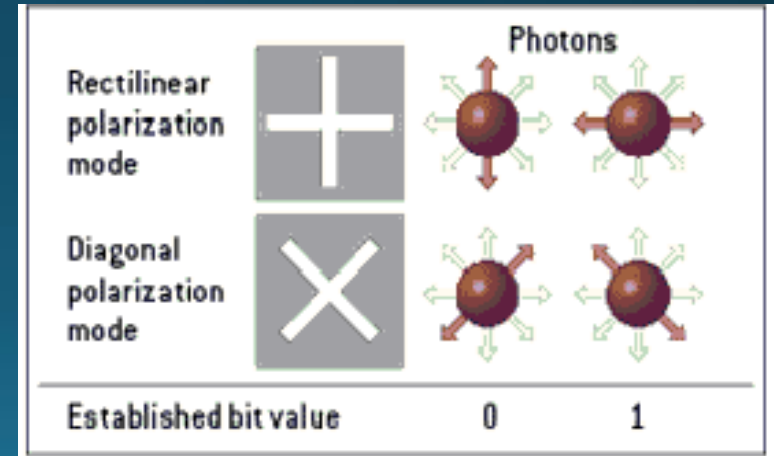
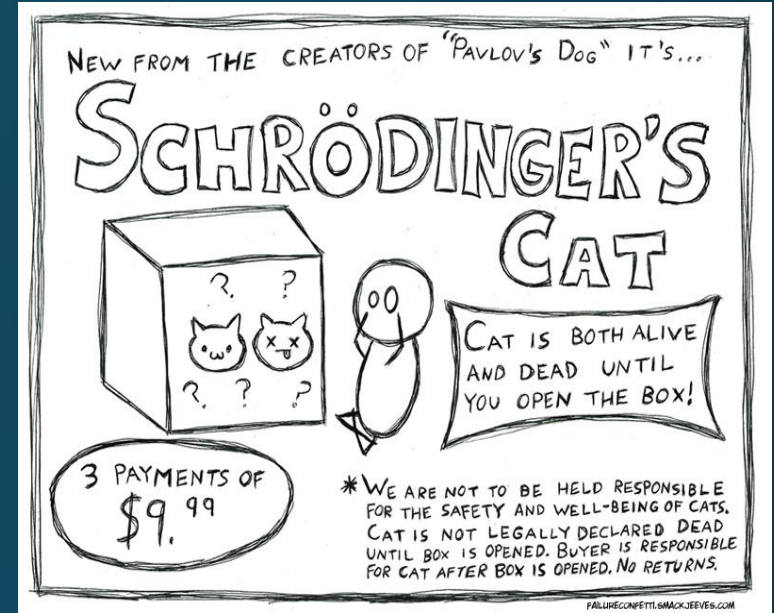
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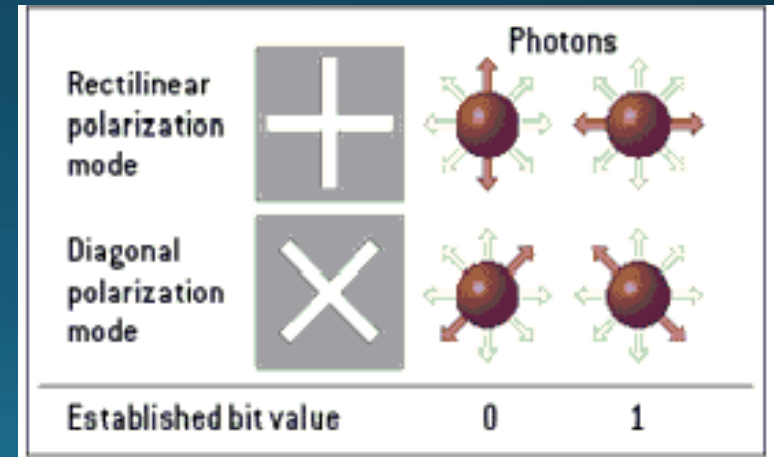
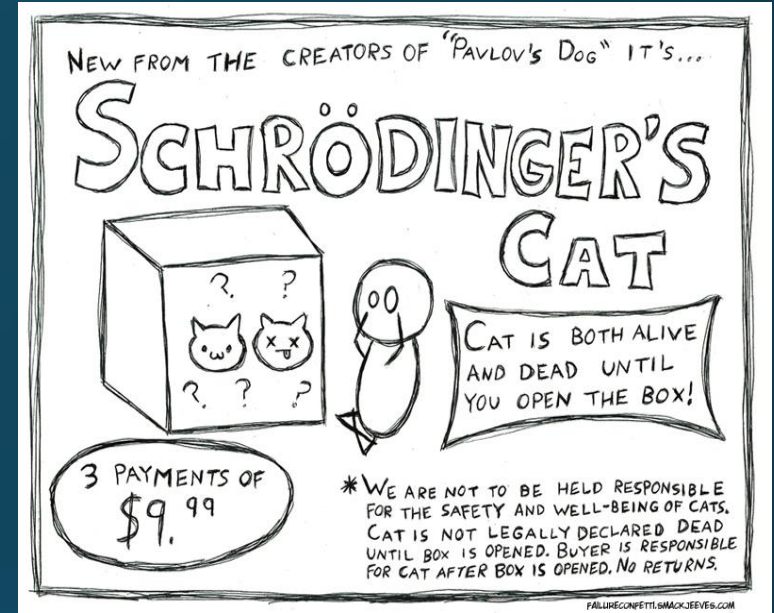
Measuring Qubits

- Measuring a qubit:
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 - filter the photon through one of the modes
- Collapses the wavefunction to one of the two states:
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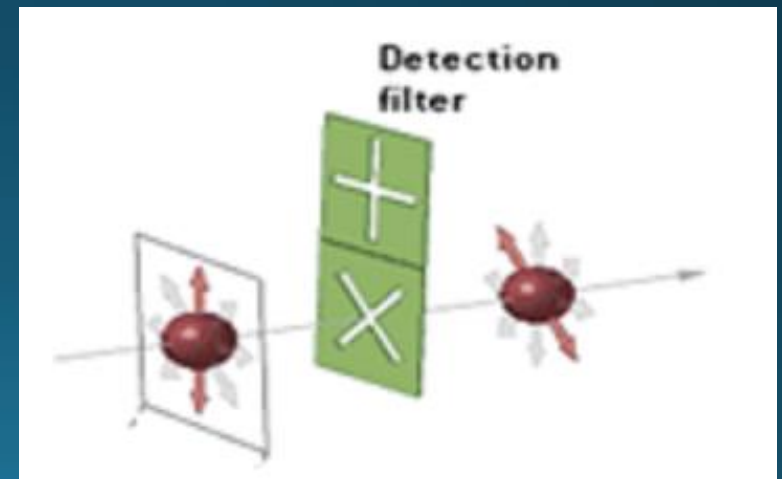
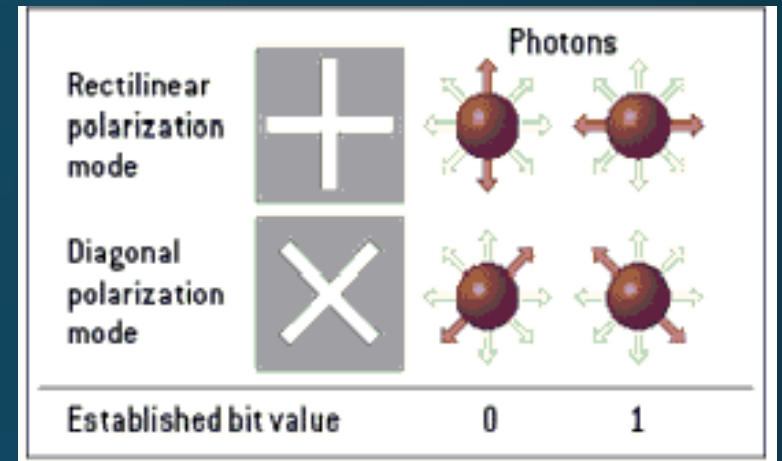


Measuring Qubits

- Measuring a qubit:
 - opening the box
 - filter the photon through one of the modes
- Collapses the wavefunction to one of the two states:
 - The cat is DEAD or ALIVE
 - The polarization is VERTICAL or HORIZONTAL
 - The bit is 0 or 1

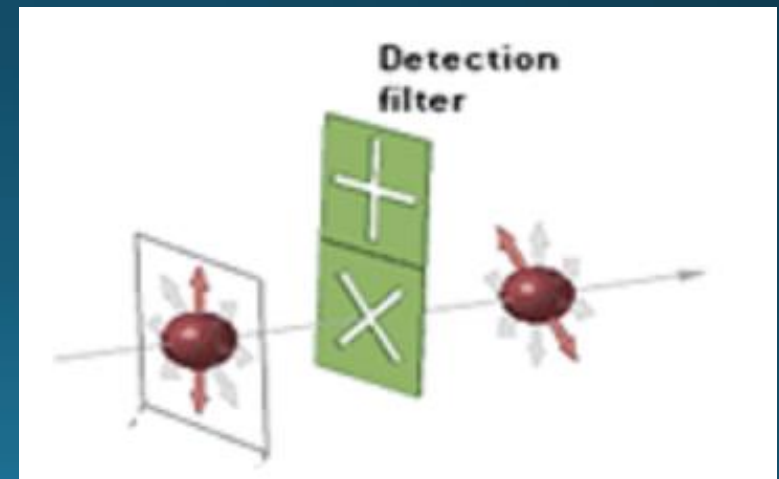
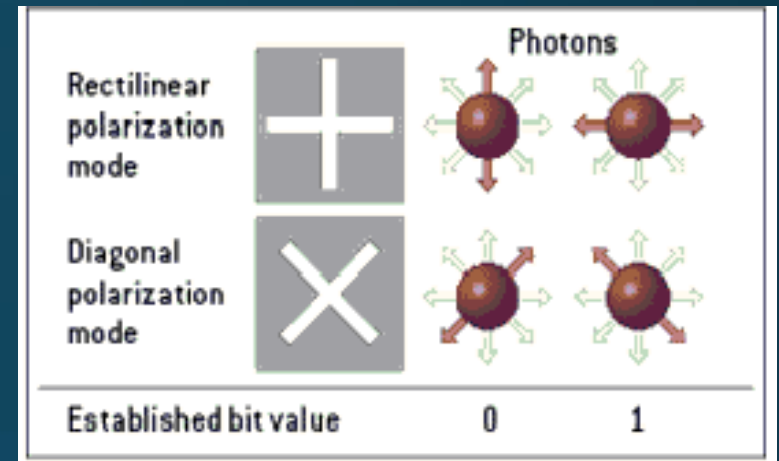


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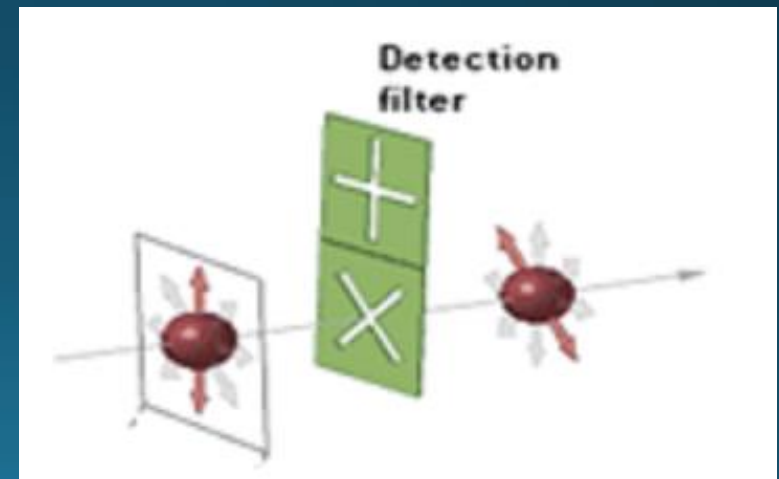
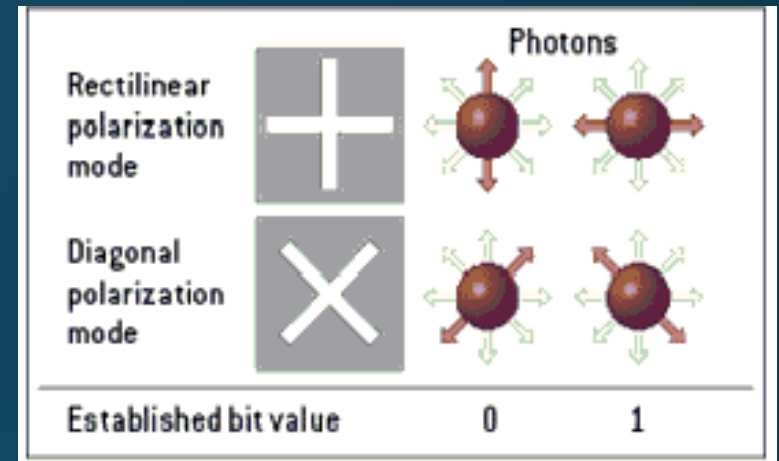
Measuring Qubits

1. Measuring the qubit in the “wrong basis”



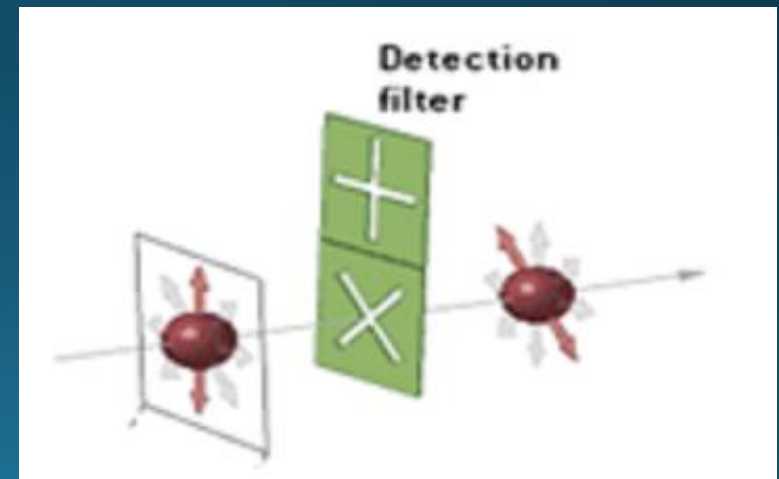
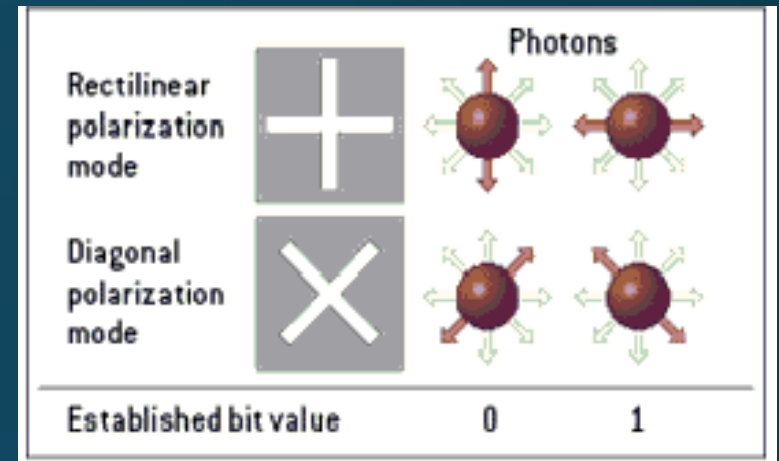
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Measuring Qubits

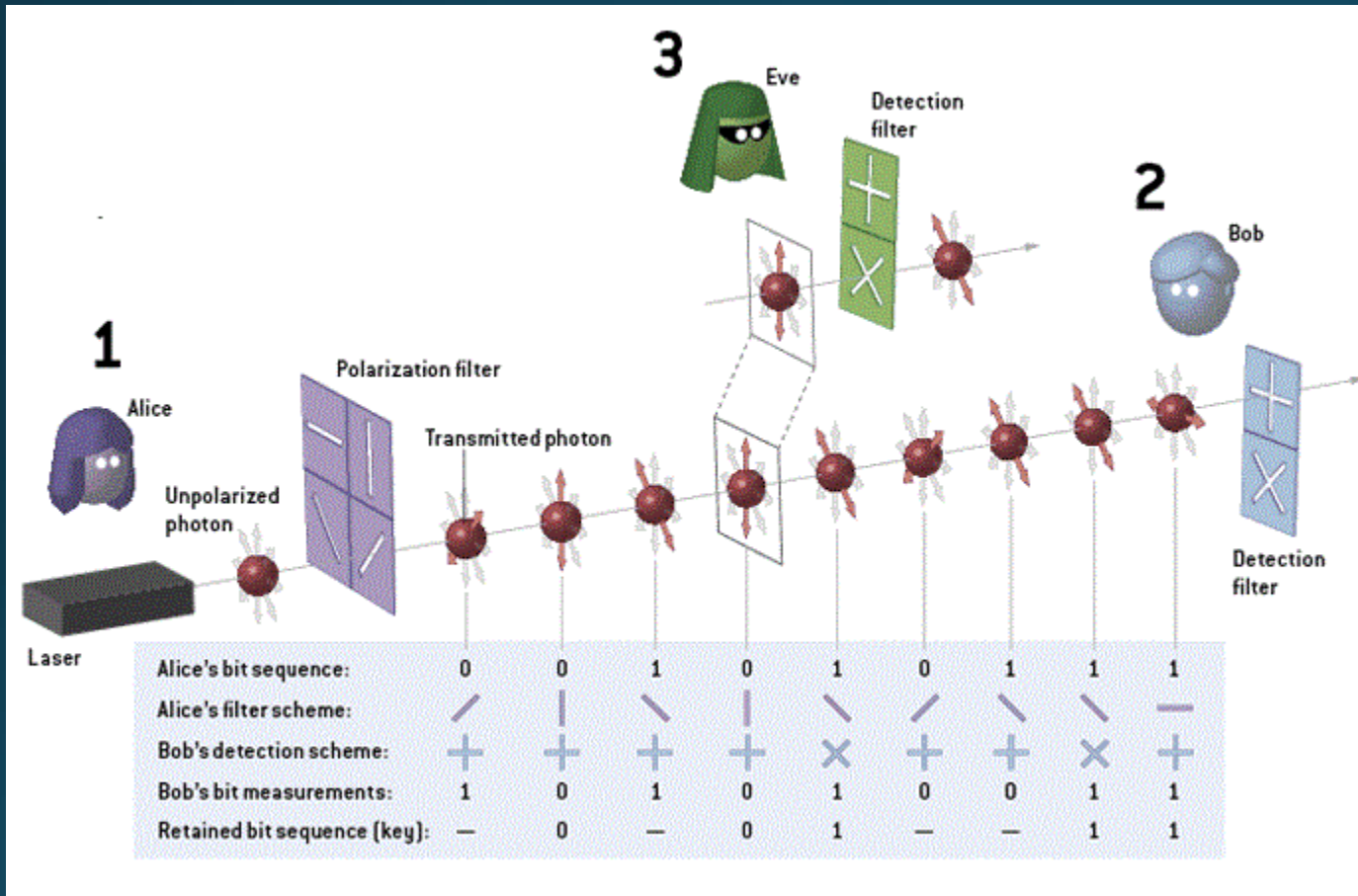
1. Measuring the qubit in the “wrong basis”
2. No information gain! (we get 0 or 1 with $P=0.5$)
3. It changes the state to one of the states corresponding to the new basis



BB84 QKD Scheme

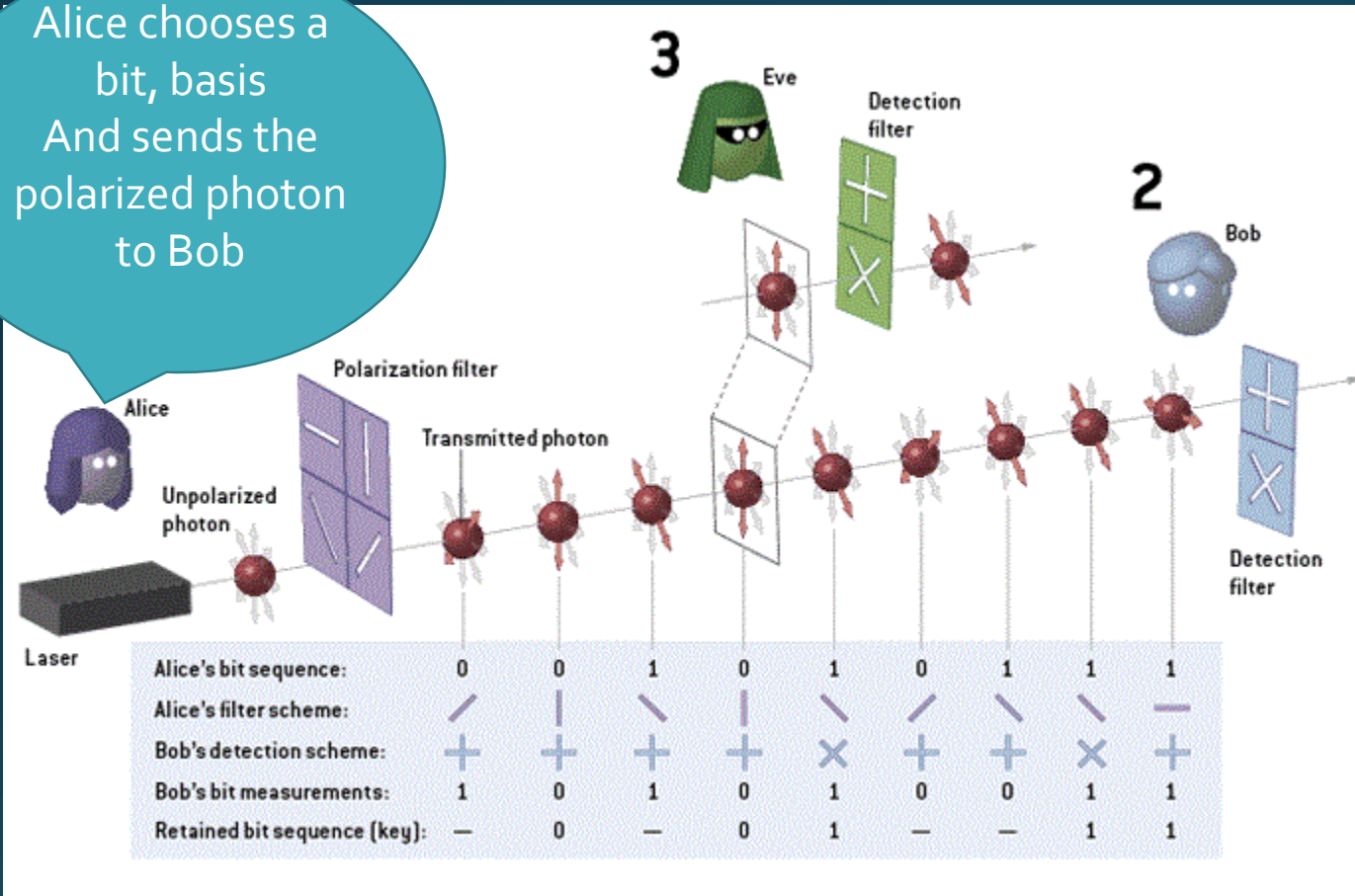
1. Alice chooses a random bit (0,1) and basis (+ or X)
2. She sends the qubit to Bob with the appropriate polarization
3. Bob measures the qubit with a random basis
4. Alice and Bob compare the string of bases they used and only keep those bits where they used the same basis
5. Error estimation and correction
6. Privacy amplification

BB84 QKD in Action



BB84 QKD in Action

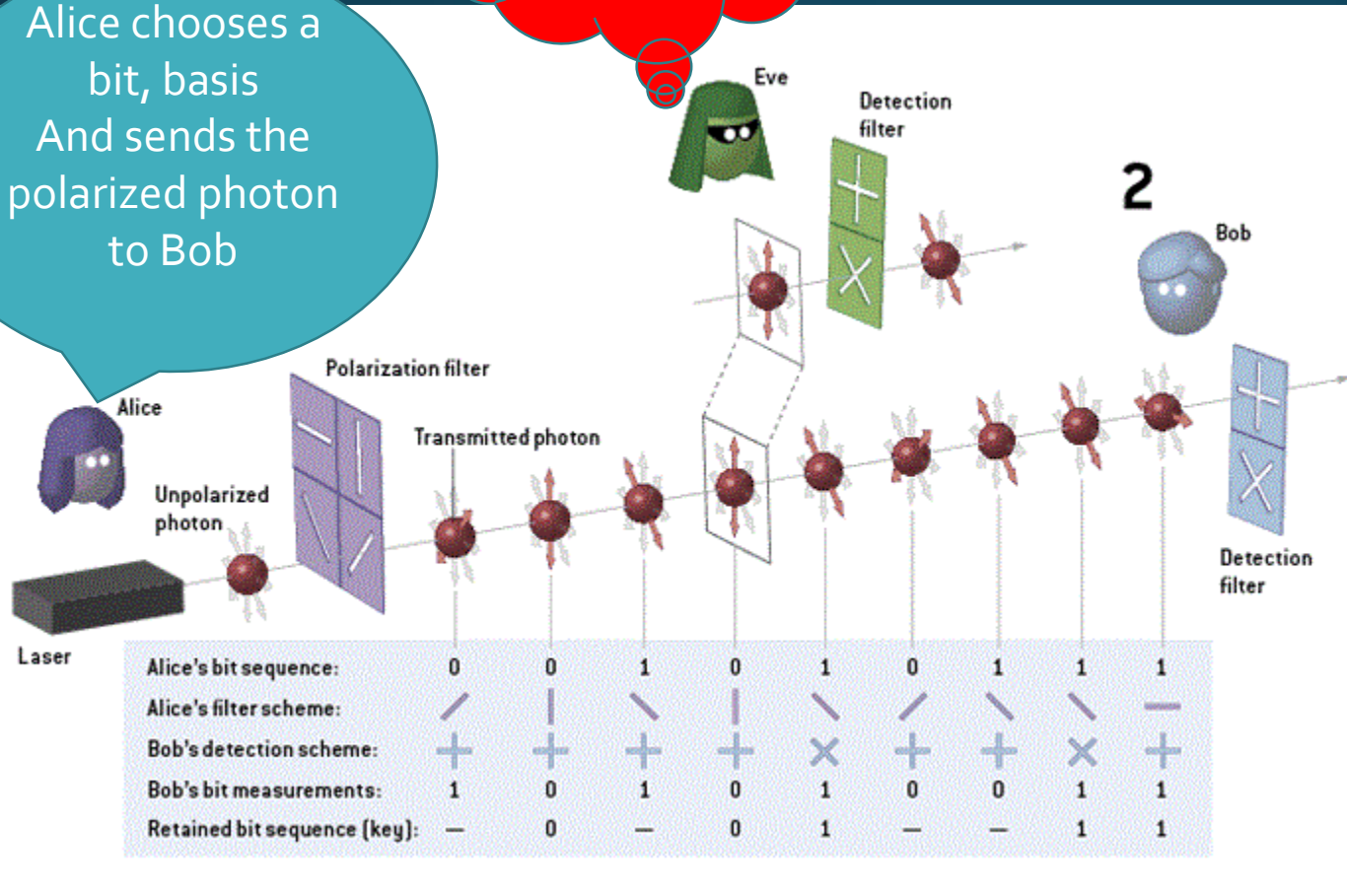
Alice chooses a bit, basis
And sends the polarized photon to Bob



BB84 Quantum Cryptography

Knows some bits!!

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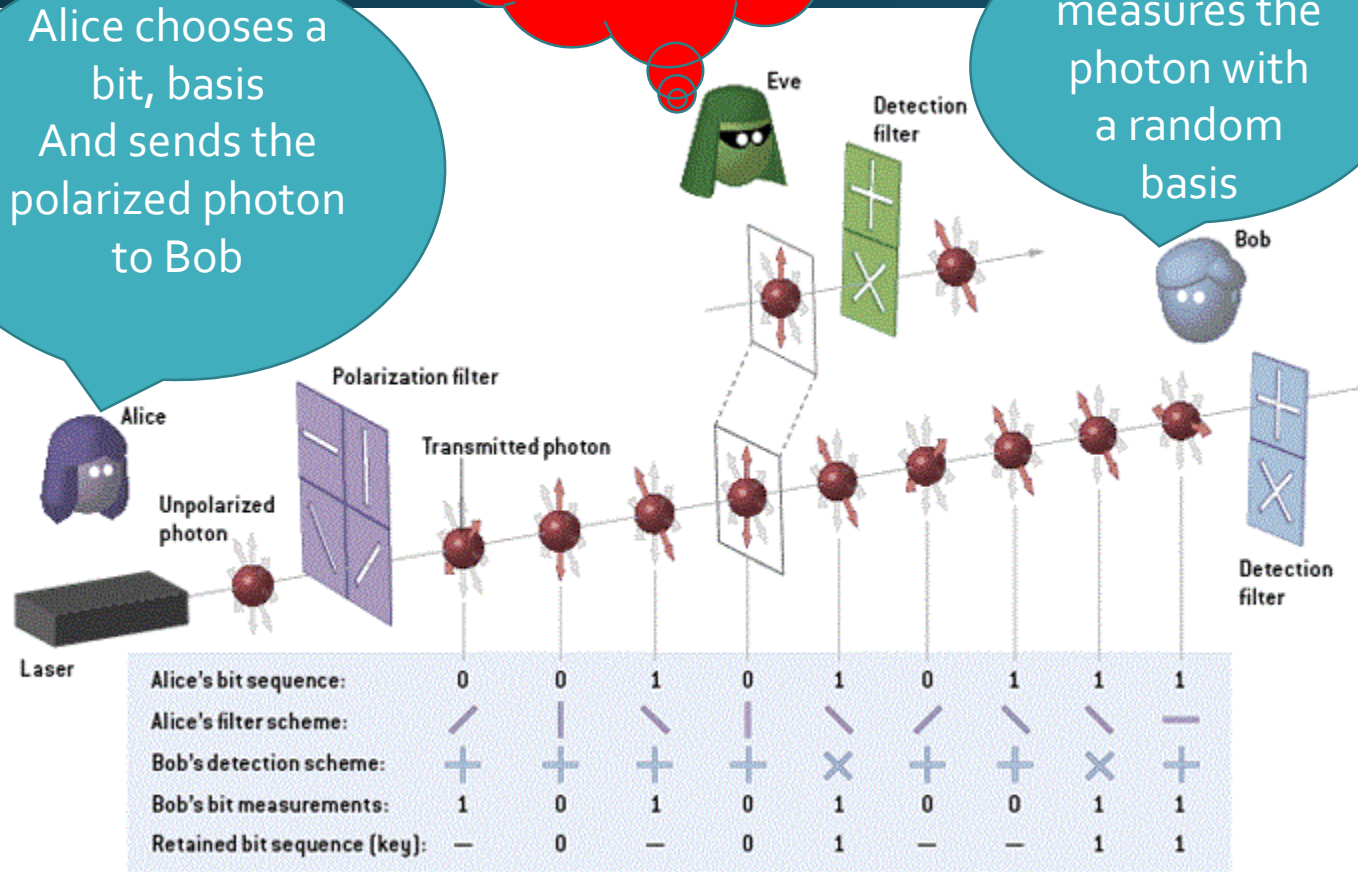


BB84 Quantum Cryptography

Knows some bits!!

Bob measures the photon with a random basis

Alice chooses a bit, basis
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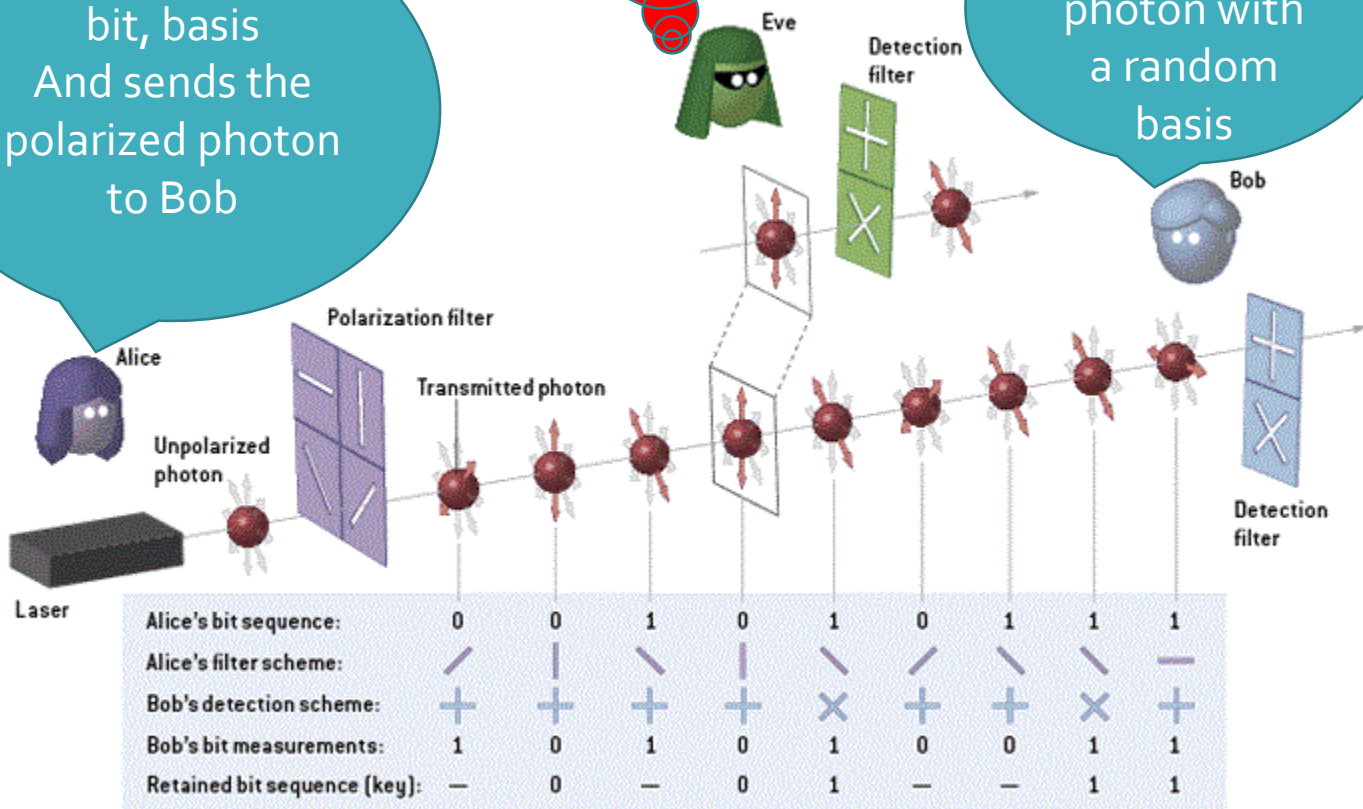


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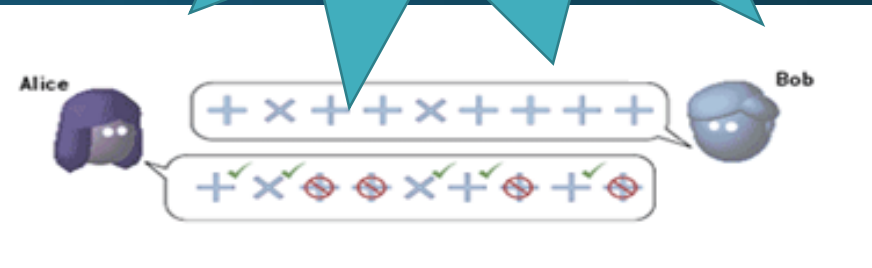
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- Confirm bases used
- Error estimation
- Privacy Amplification

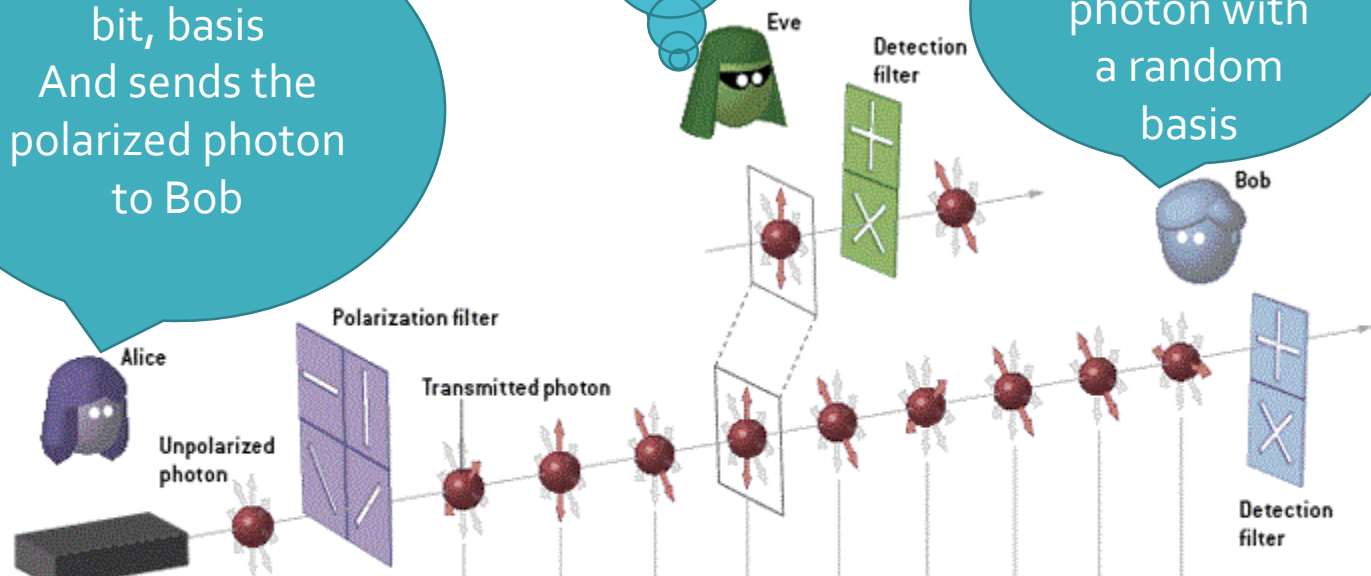


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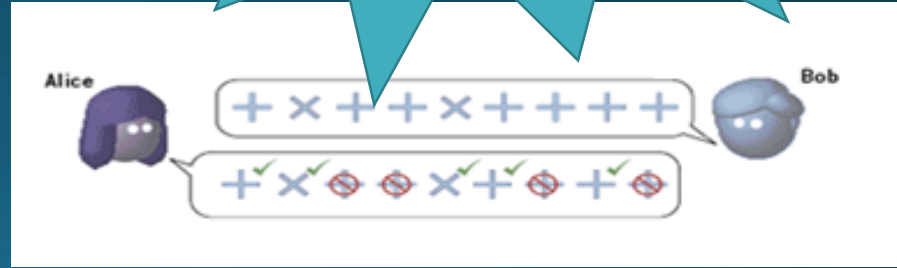
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Bob measures the photon with a random basis



Laser									
Alice's bit sequence:	0	0	1	0	1	0	1	1	1
Alice's filter scheme:	↘	↑	↘	↑	↘	↘	↘	↘	↘
Bob's detection scheme:	+	+	+	+	×	+	+	×	+
Bob's bit measurements:	1	0	1	0	1	0	0	1	1
Retained bit sequence (key):	-	0	-	0	1	-	-	1	1

-Confirm bases used
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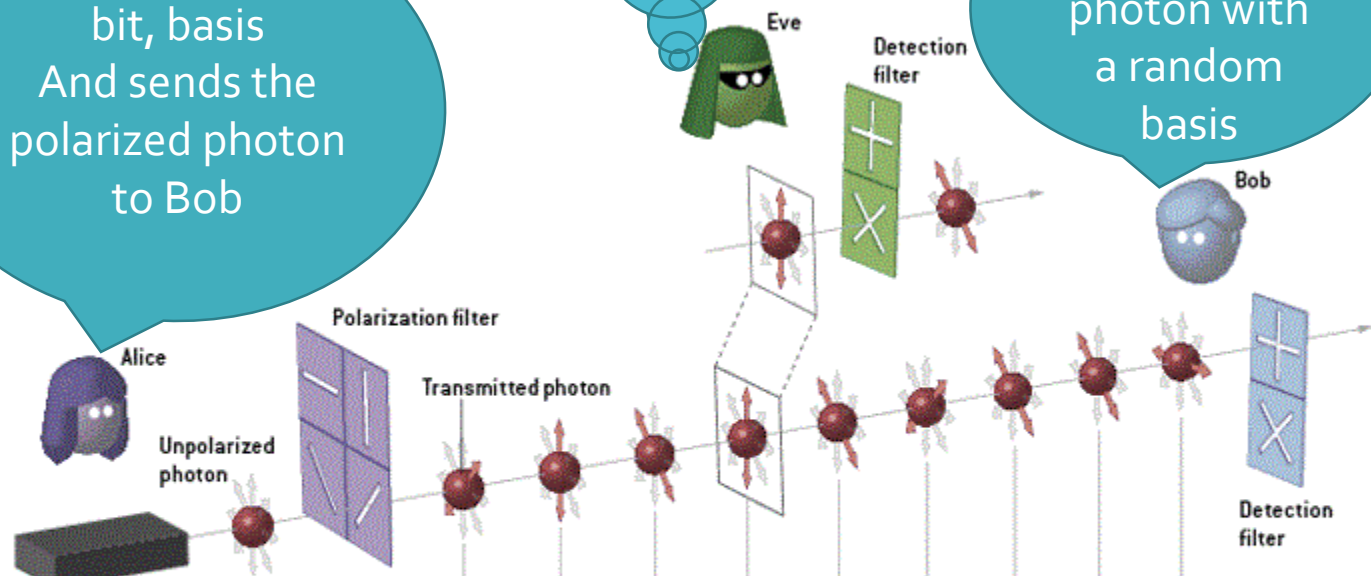


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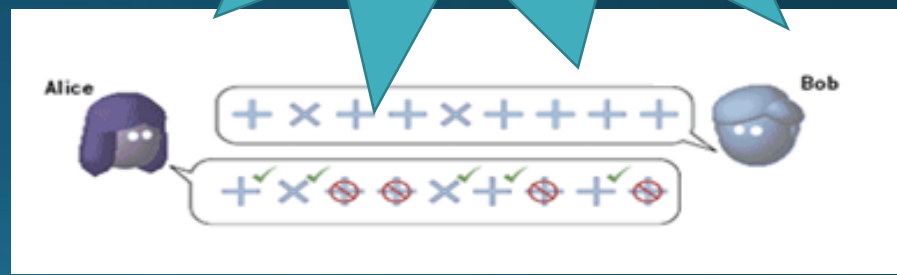
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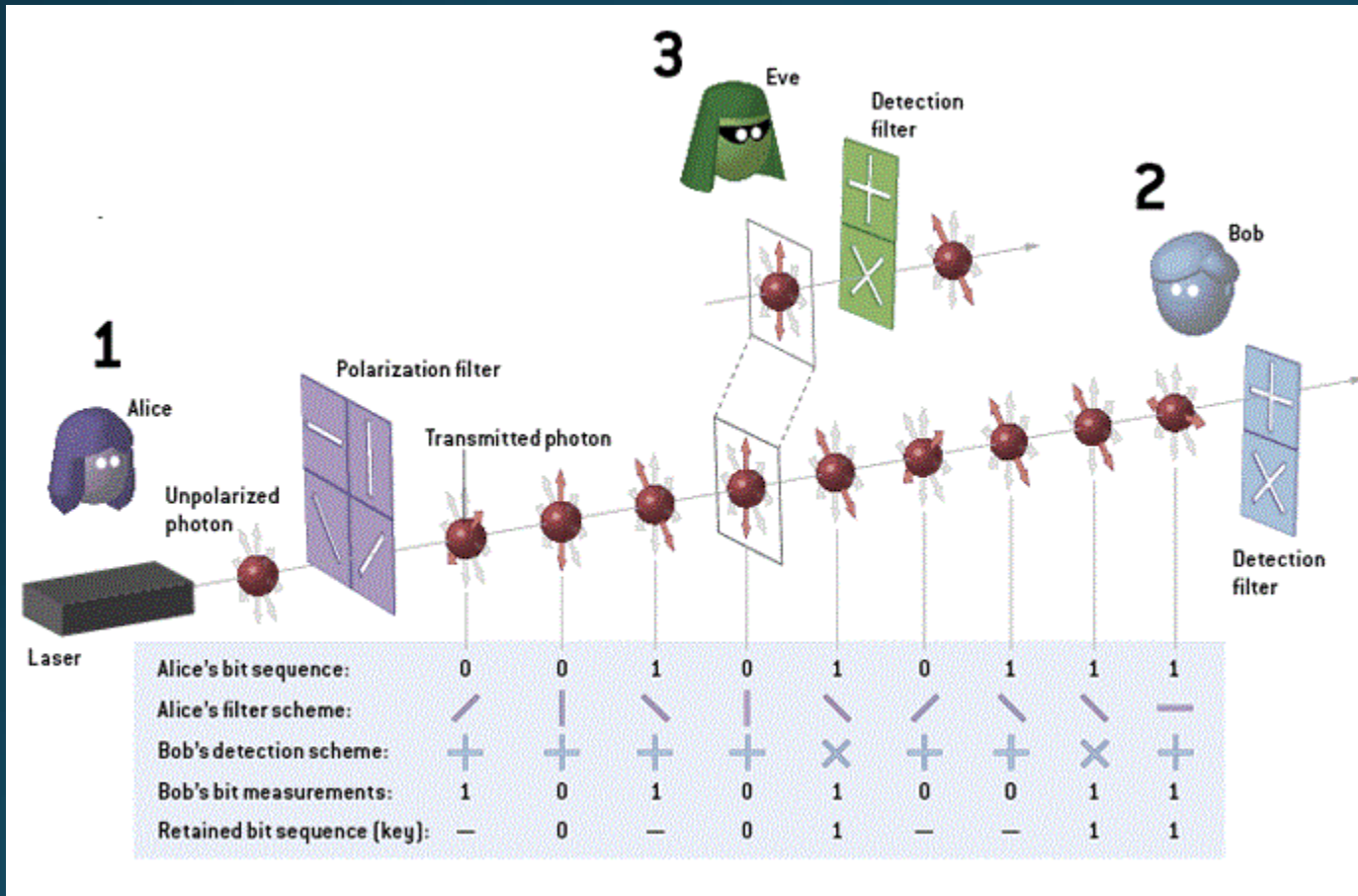
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This scheme can be proven to be perfectly secure!

-Confirm bases used
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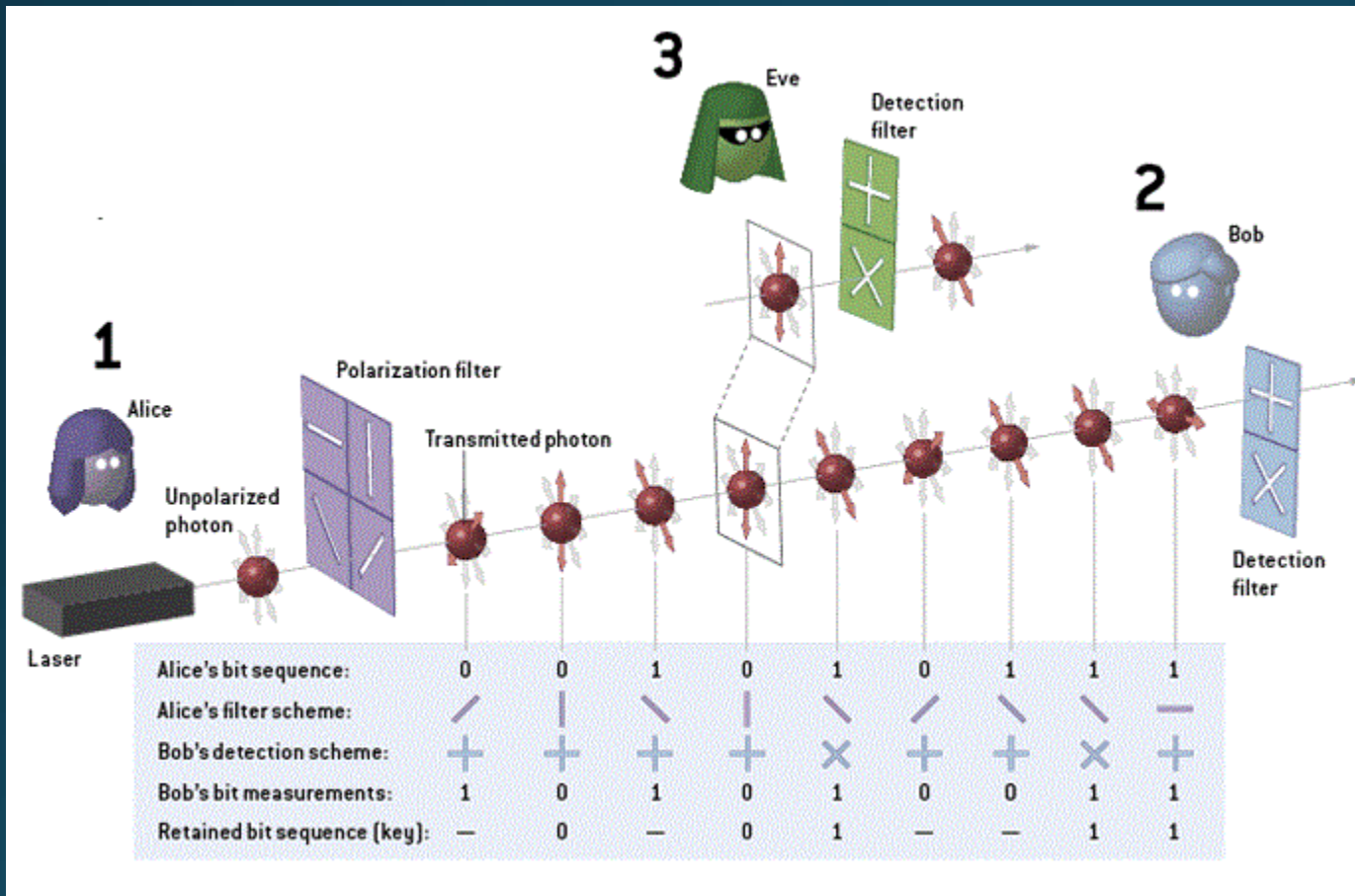


Detecting an eavesdropper



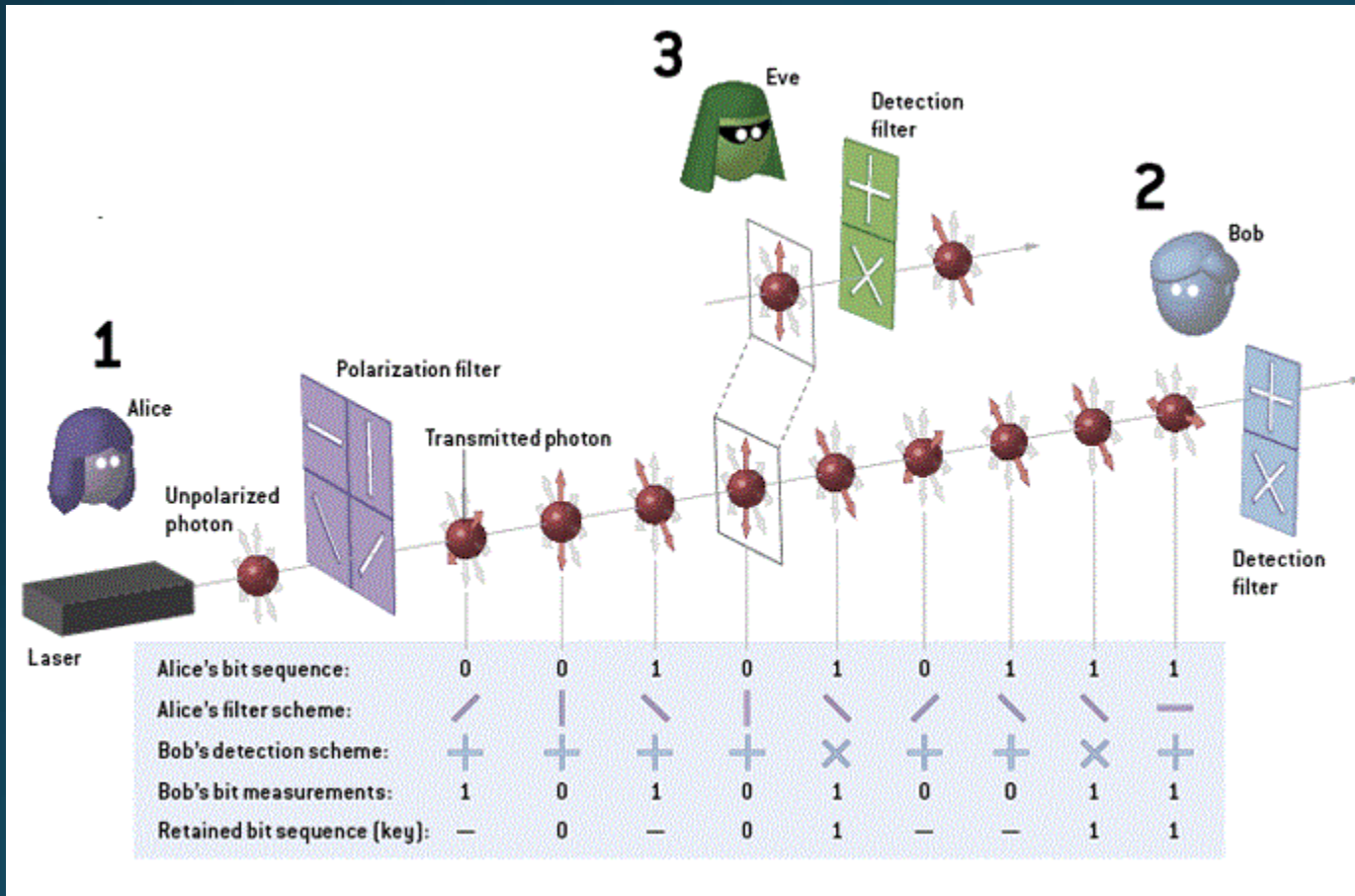
Detecting an eavesdropper

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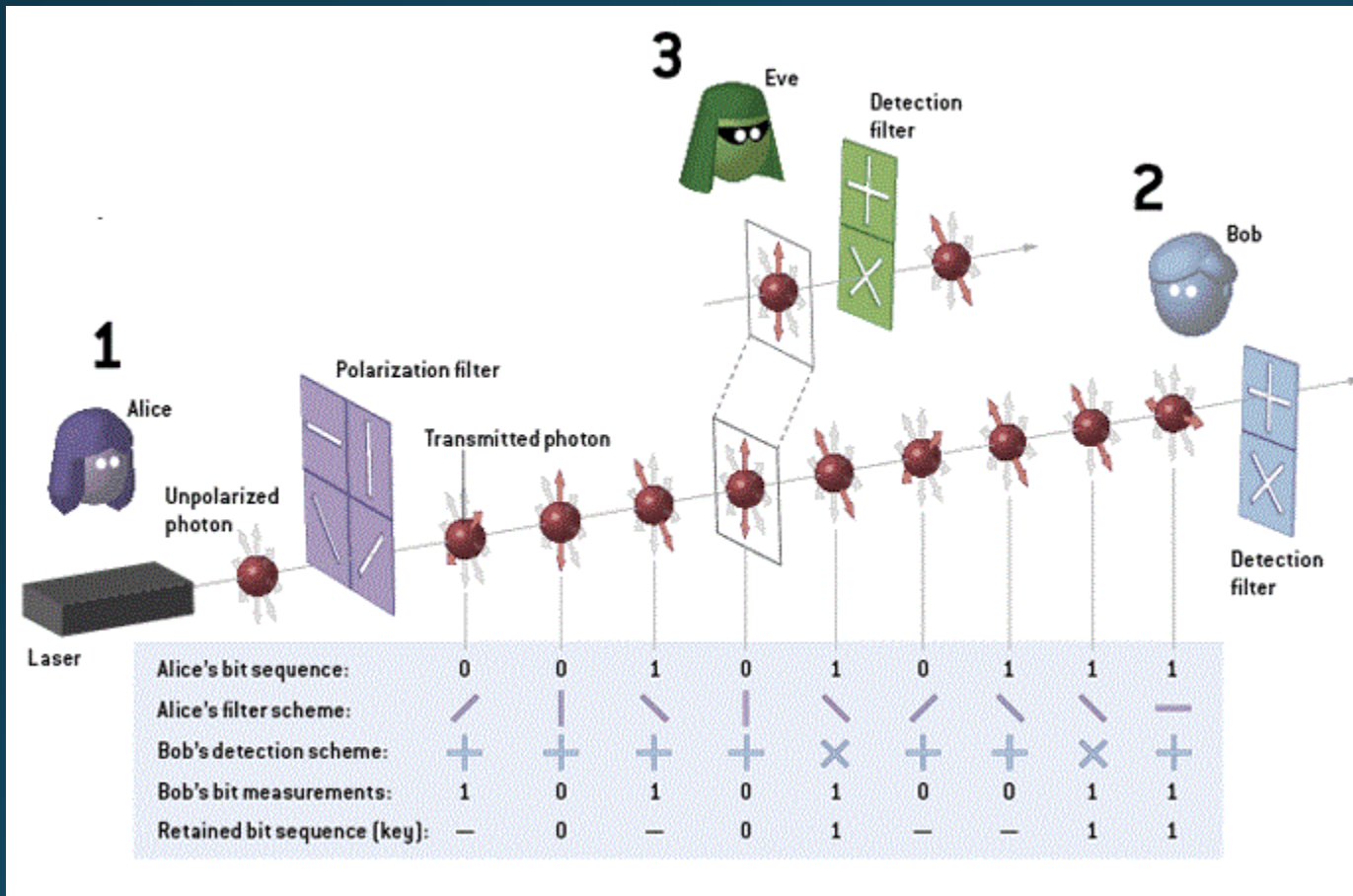


Detecting an eavesdropper

- If Eve measures a state in the wrong basis she will change the state of the photon
- This might introduce errors that can be detected by Alice and Bob



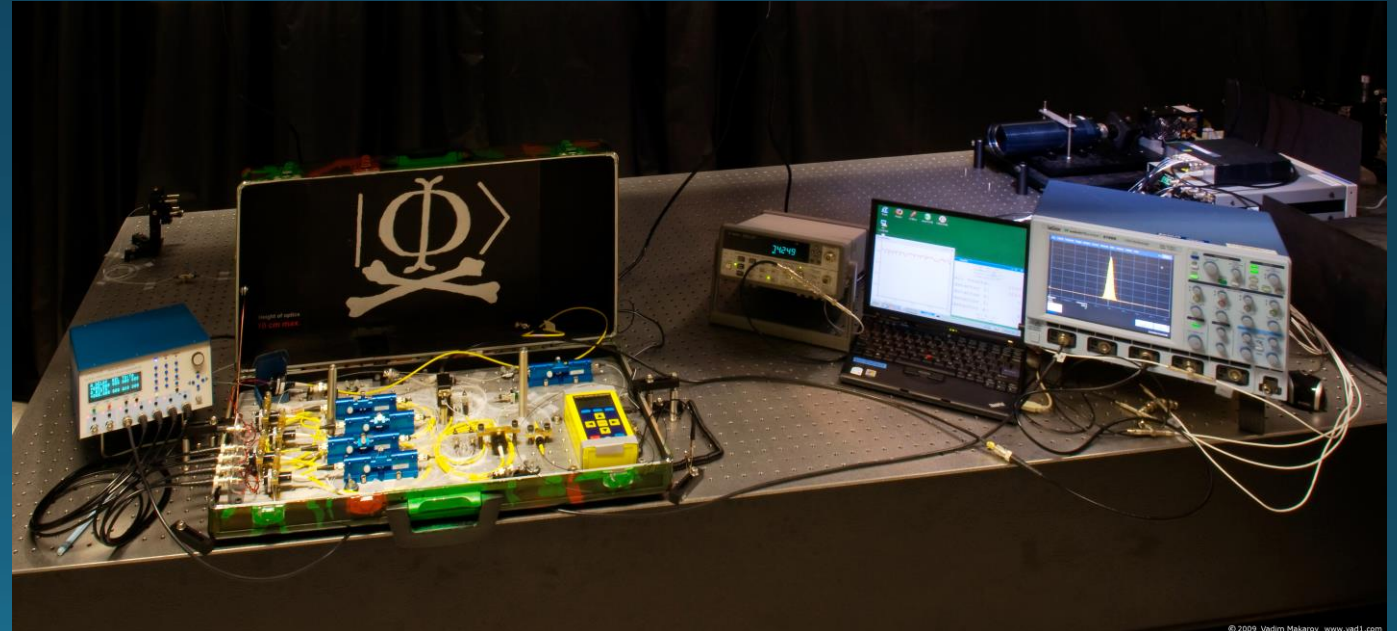
Detecting an eavesdropper



- If Eve measures a state in the wrong basis she will change the state of the photon
- This might introduce errors that can be detected by Alice and Bob
- If too many errors are detected they know that there was an eavesdropper and abort

Applications

- Commercial QKD systems already exist
- 2007 Voting in Geneva [5]
- Approximately 4 commercial companies
- and 5 Quantum Key Distribution Networks



Too good to be true?

- Distances: ~200km using optic fiber
and much less through free space (air)
- Expensive equipment
- Imperfect implementations, at least two successful attacks

Conclusions

- Quantum Cryptography only relies on laws of nature
- Post-quantum cryptography relies on primitive that are difficult for quantum and classical computers
- Quantum Key Distribution allows two parties to share a key using a public quantum channel
- QKD schemes are perfectly secure, possible and work in practice although the implementation of them so far is not perfect

References

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I SURVIVED

SUCK IT, SCHRODINGER