

Introduction to Modern Cryptography

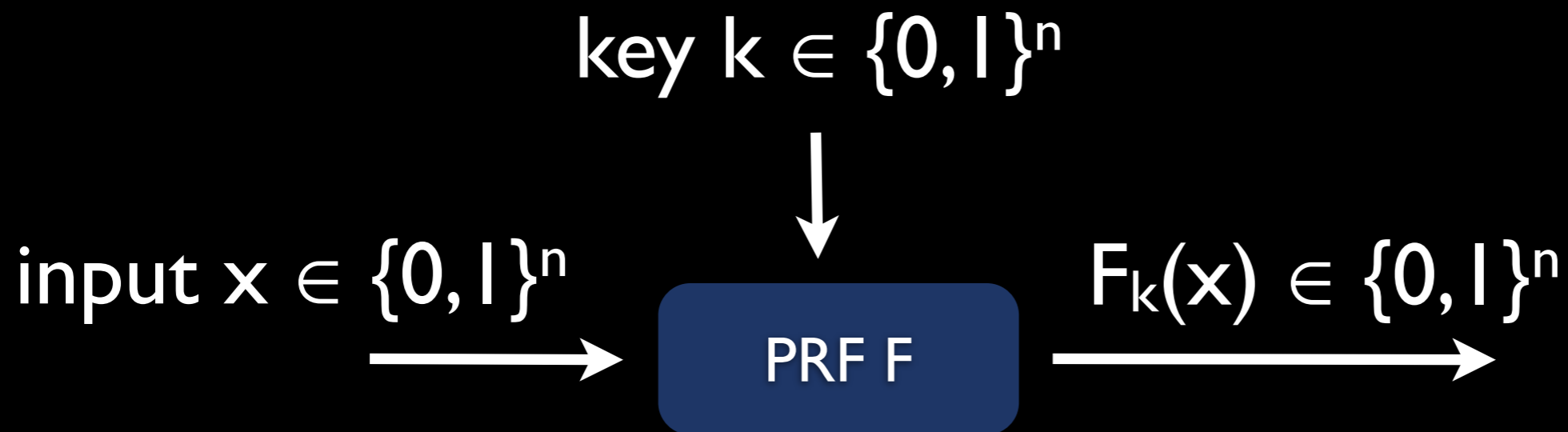


4th lecture:

Pseudorandom Functions and
Chosen-Plaintext Security

some of these slides are copied from or heavily inspired by the
University College London MSc InfoSec 2010 course given by Jens Groth
Thank you very much!

PRG vs PRF



- existence of PRF \Leftrightarrow existence of PRG
- both can be based on one-way functions

Battle of Midway (1942)

- Midway Atoll: [Wikipedia](#), [Google Maps](#)
- important naval battle between the USA and Japan in World War II ([Wikipedia](#))
- decided by **cryptographic skills**
- US tricked Japanese into **acting as encryption oracle**
- bottom line: the use of **CPA secure encryption** could have changed the course of world history

Breaking News (1 Nov 2012)

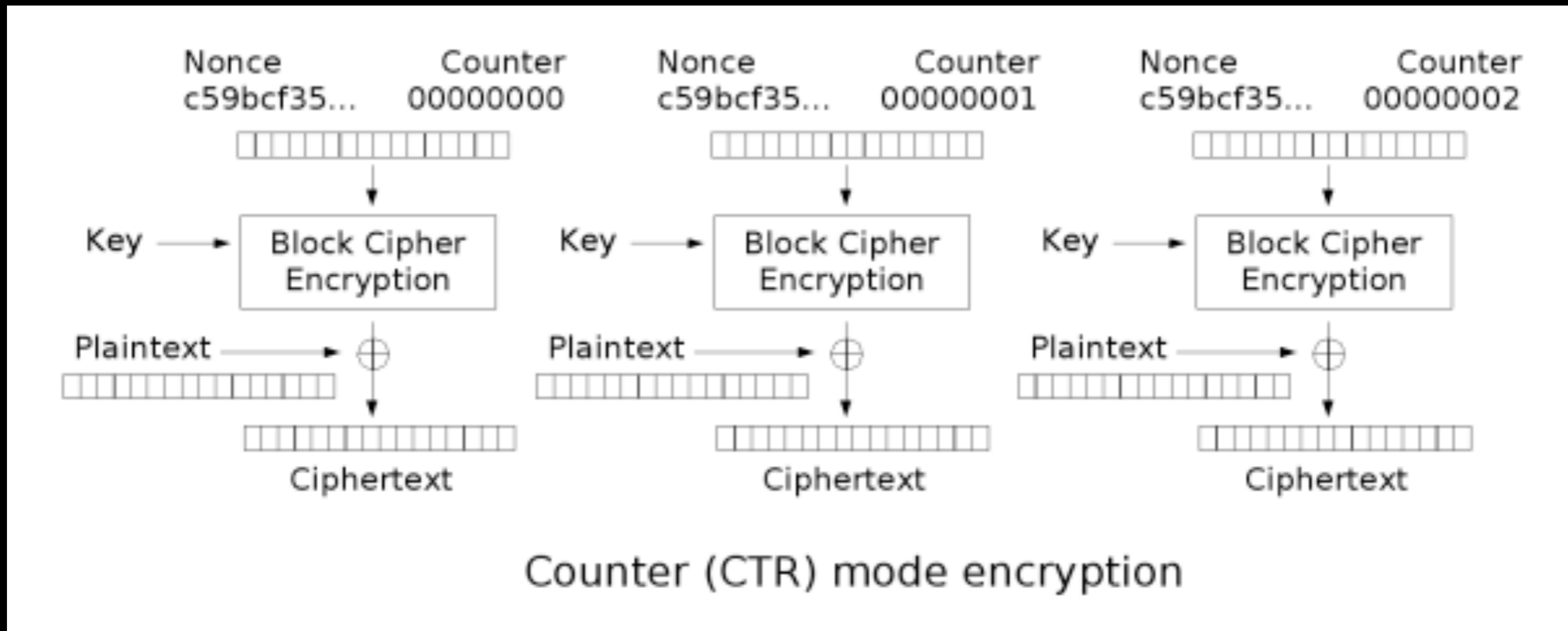


AOAKN HVPKD FNFJU YIDDC
RQXR DJHFP GOVEN MIAPX
PABUZ WYND CMPNW HJREH
NLXKE HENIK ONOIB AKEL4
HAOTA RBQRH DJOFM TPZEH
LKXEH RECHT JRZCQ FNKTO
KLDTS EQIRU AOAKN 27 1525/0.

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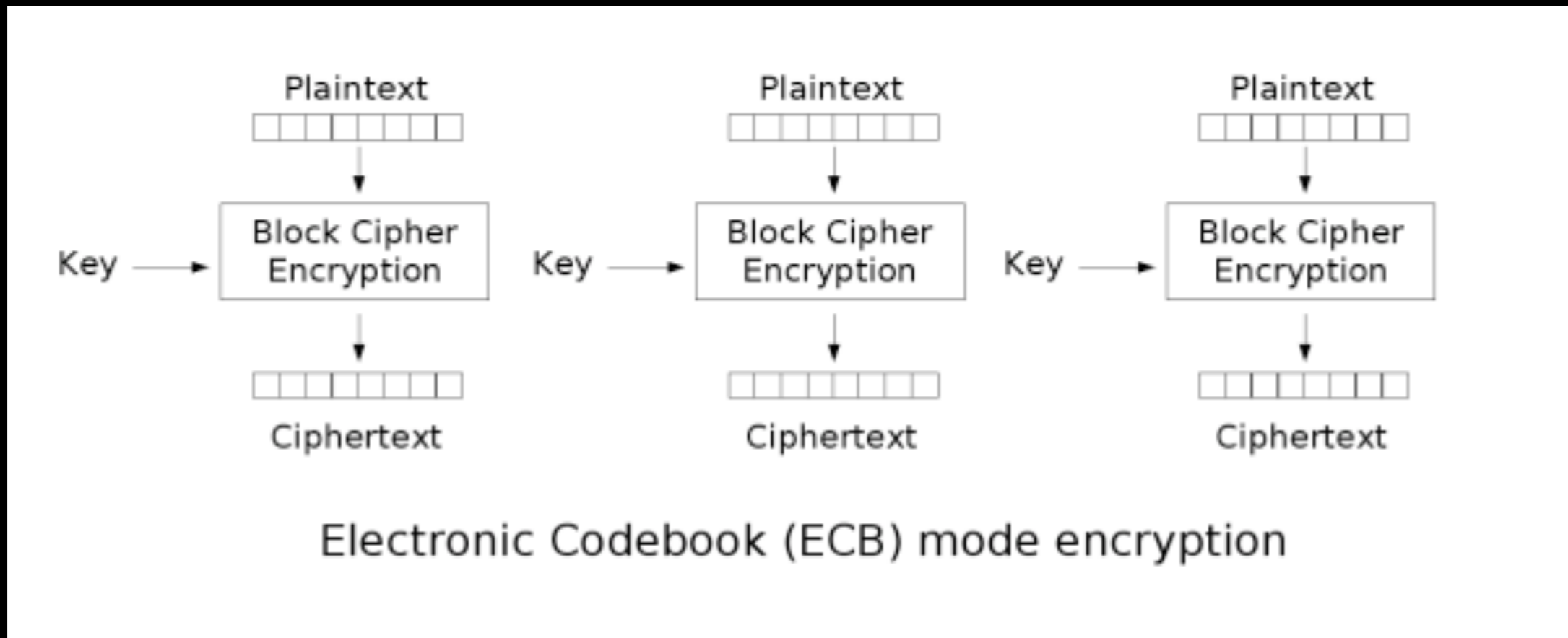
- another challenge for you to solve
- watch the BBC news story

Counter (CTR) mode



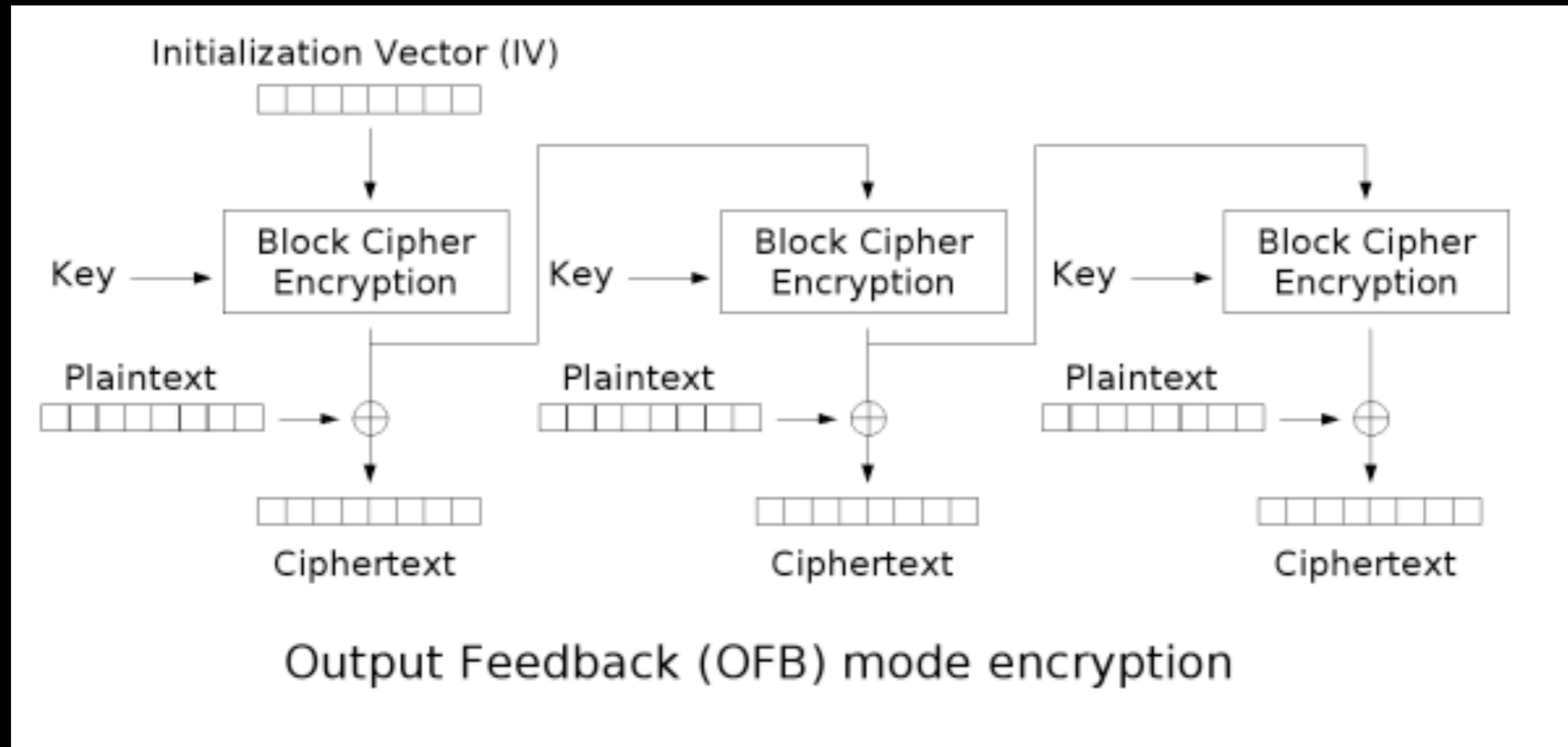
- CTR mode is CPA-secure if F (the Block Cipher) is a pseudorandom function
- can be precomputed and fully parallelized
- allows random access

Electronic Code Book (ECB)



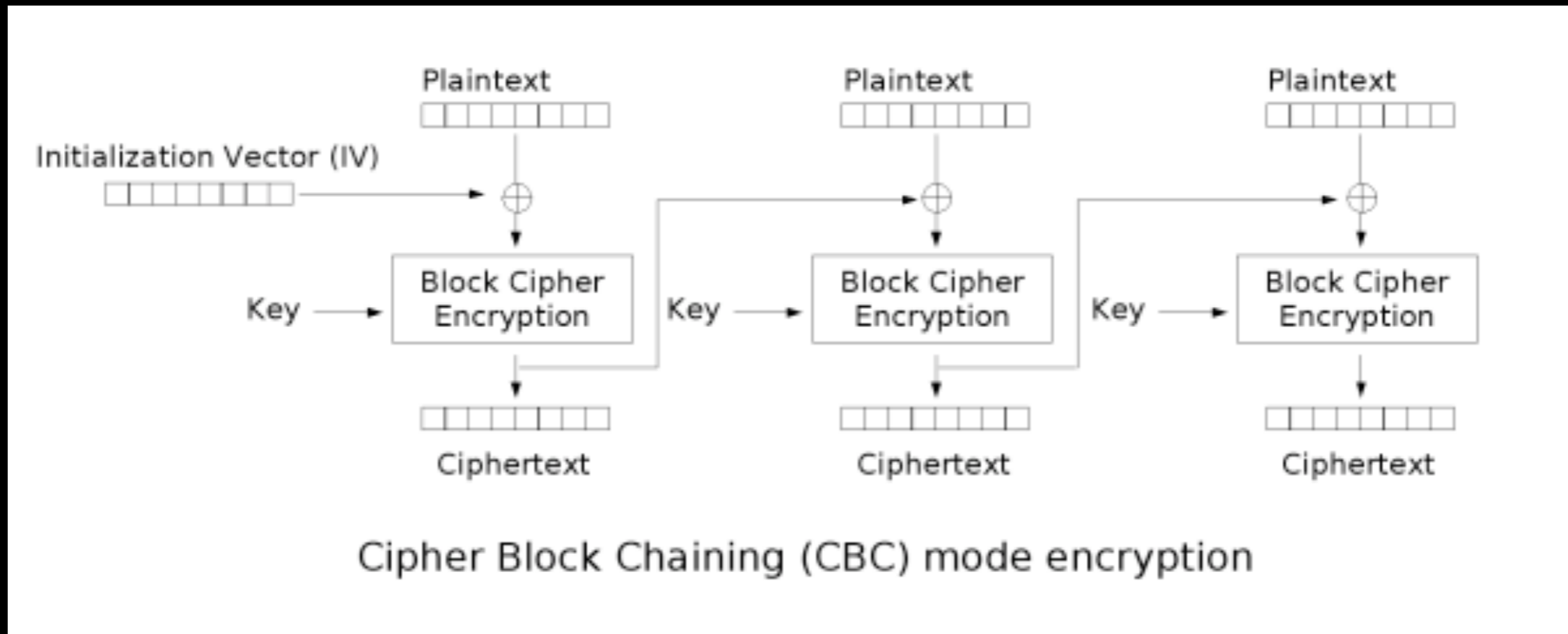
- highly insecure, should **never** be used
- see example on [wikipedia](#)

Output Feedback (OFB)



- if F is pseudorandom function, then OFB is CPA-secure
- advantage: pseudorandom stream can be precomputed

Cipher Block Chaining (CBC)



- if F is pseudorandom permutation, then CBC is CPA-secure
- drawback: encryption is sequential