

Introduction to Modern Cryptography

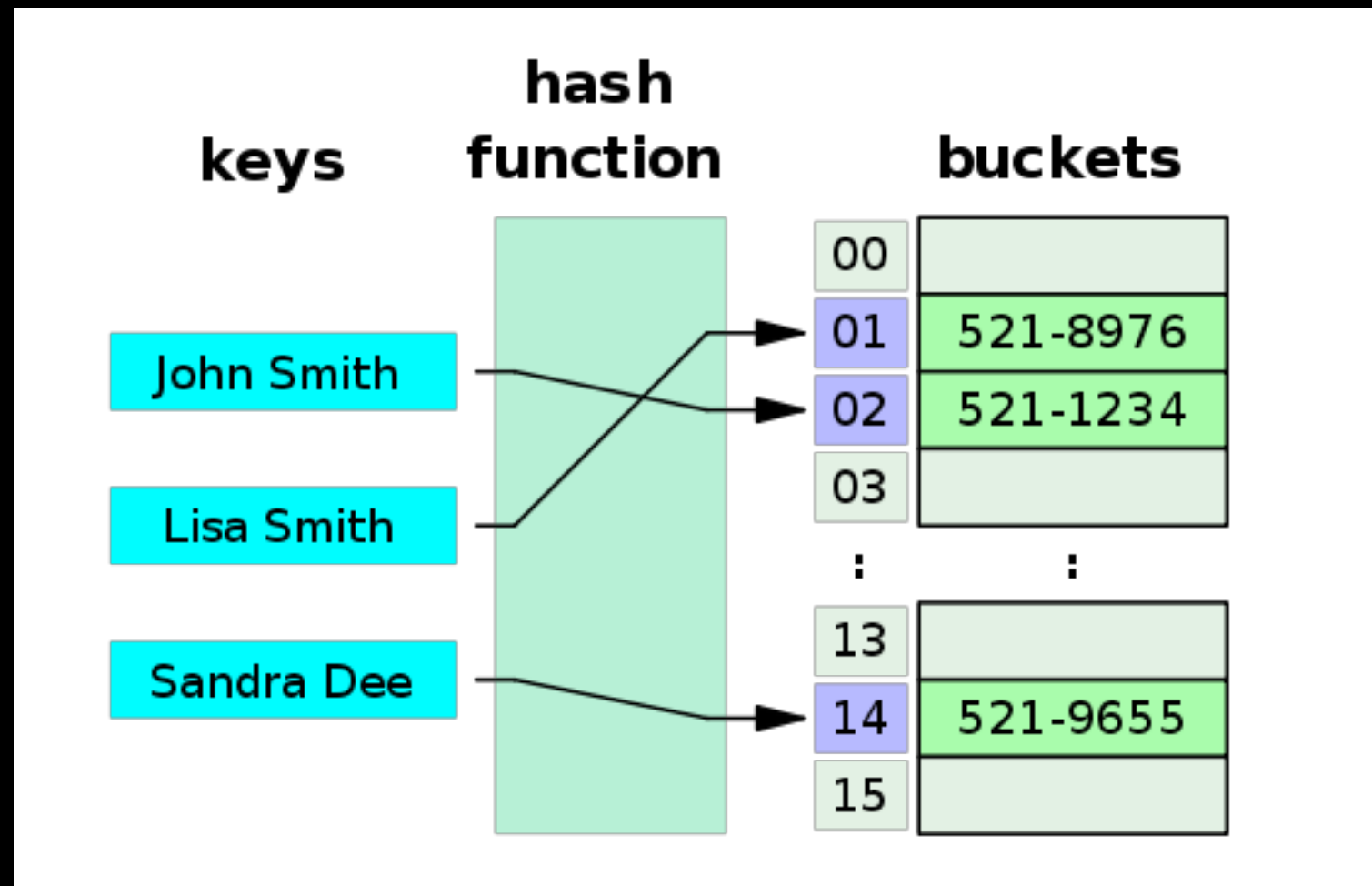


6th lecture:

Collision-Resistant Hash Functions

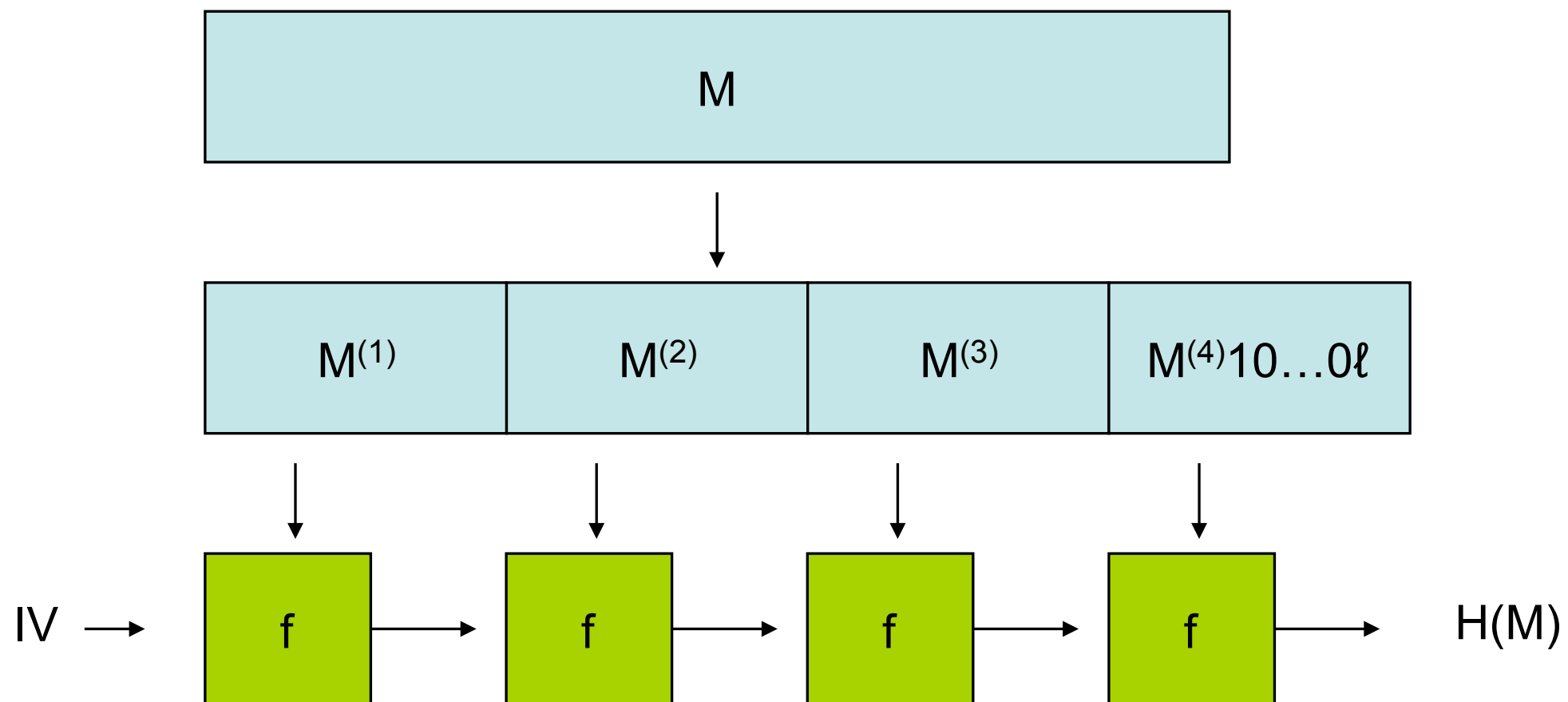
quite a few of these slides are copied from or heavily inspired by the
University College London MSc InfoSec 2010 course given by Jens Groth
Thank you very much!

Hash Tables



- used in data structures to store associative arrays
- try to minimize collisions!

Merkle-Damgård Construction

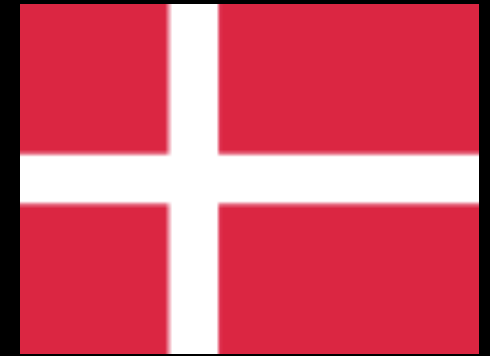


Ralph C. Merkle



- co-inventor of public-key crypto
- Merkle puzzles, Merkle trees
- new ideas are hard to publish
- now interested in nanotechnology and cryonics

Ivan Bjerre Damgård



- most publishing cryptographer in the world
- my PhD advisor
- amazing person
- plays the fiddle

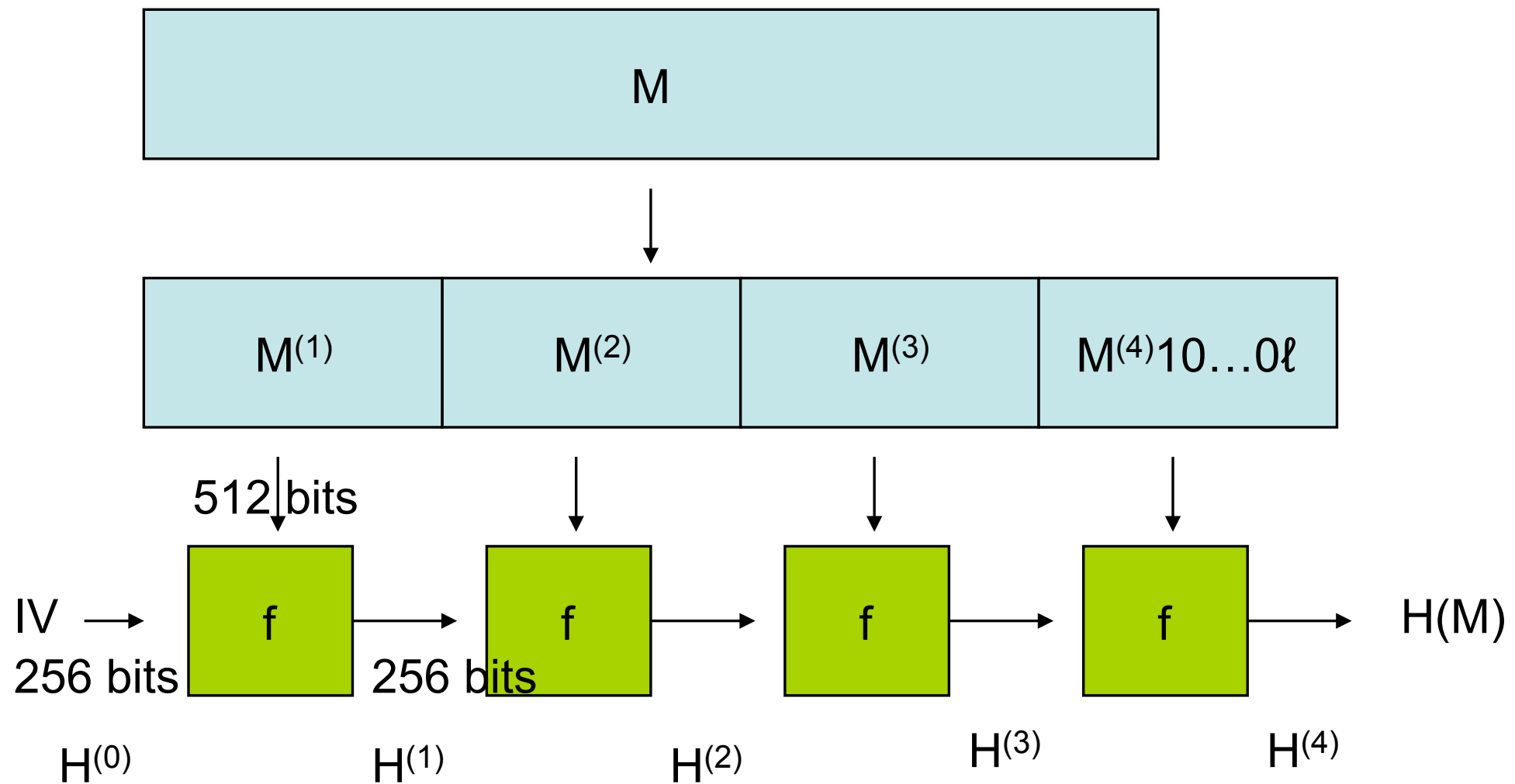
Birthday Attacks on Hash Functions

- In a class of N students with random birthdays $b_1, \dots, b_N \leftarrow \{1, 2, \dots, 365\}$. How large does N need to be such that $\Pr[\text{exists } i \neq j : b_i = b_j] > 1/2$?
- Answer: $N \geq \sqrt{365} \approx 23$
- Task: Find a generic collision-attack on the hash function $H: \{0, 1\}^* \rightarrow \{0, 1\}^n$

Finding Meaningful Collisions

- It is {hard, difficult, challenging, impossible} to {imagine, believe} that we will {find, locate, hire} another {employee, person} having similar {abilities, skills, character} as Alice. She has done a {great, super} job
- $4 \times 2 \times 3 \times 2 \times 3 \times 2 = 288$ possibilities
- prepare $2^{n/2}$ of those and $2^{n/2}$ of bad ones

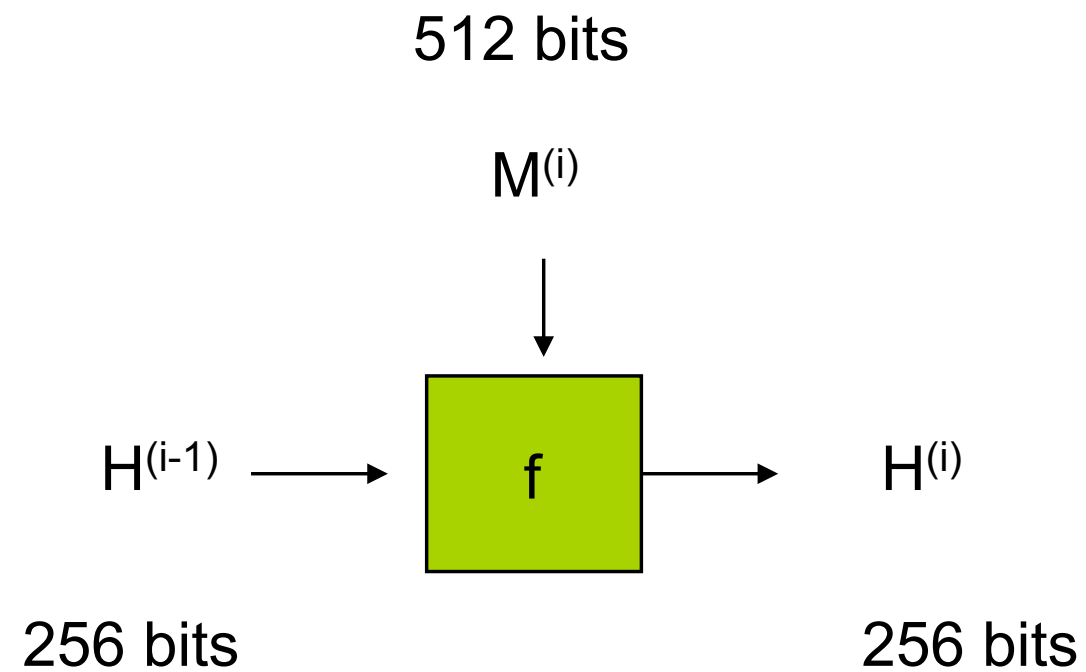
SHA-256



- $H(0) = 6a09e667\ bb67ae85\ 3c6ef372\ a54ff53a$
 $510e527f\ 9b05688c\ 1f83d9ab\ 5be0cd19$

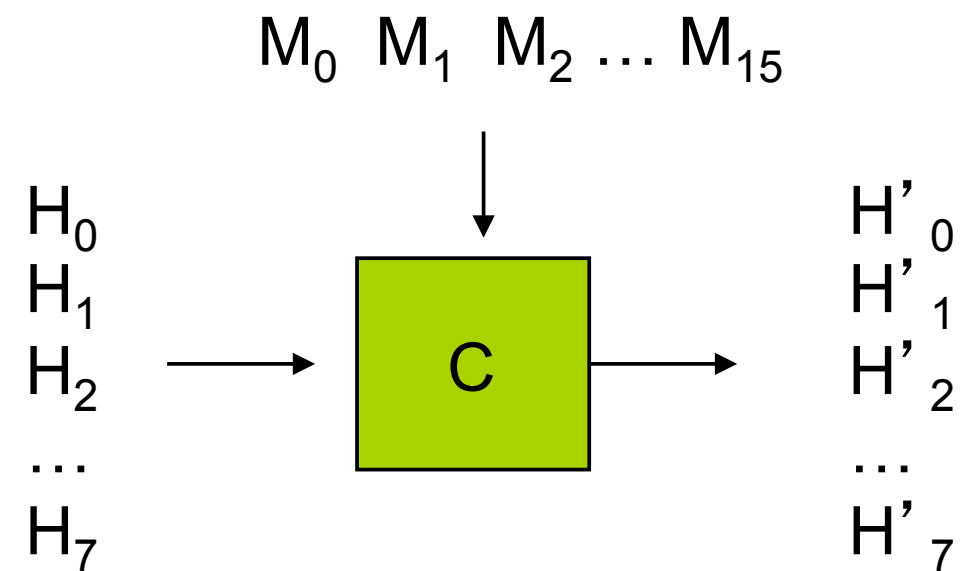
- first 32 bits of fractional parts of square roots of 2,3,5,7,11,13,17

SHA-256 compression function



SHA-256 compression function

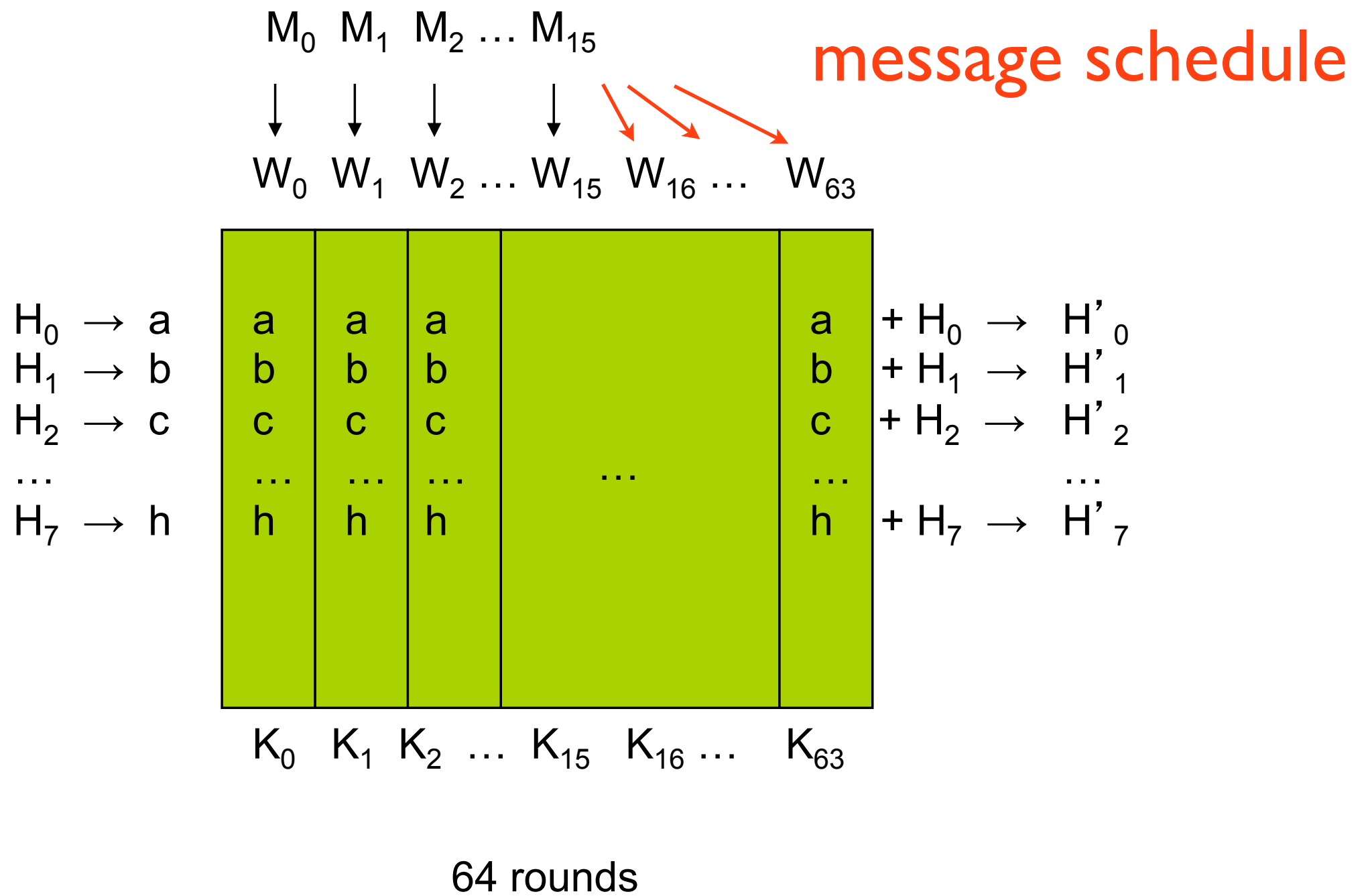
512 bits = 16 x 32 bit words



256 bits =
8 x 32 bit words

256 bits =
8 x 32 bit words

SHA-256 compression function

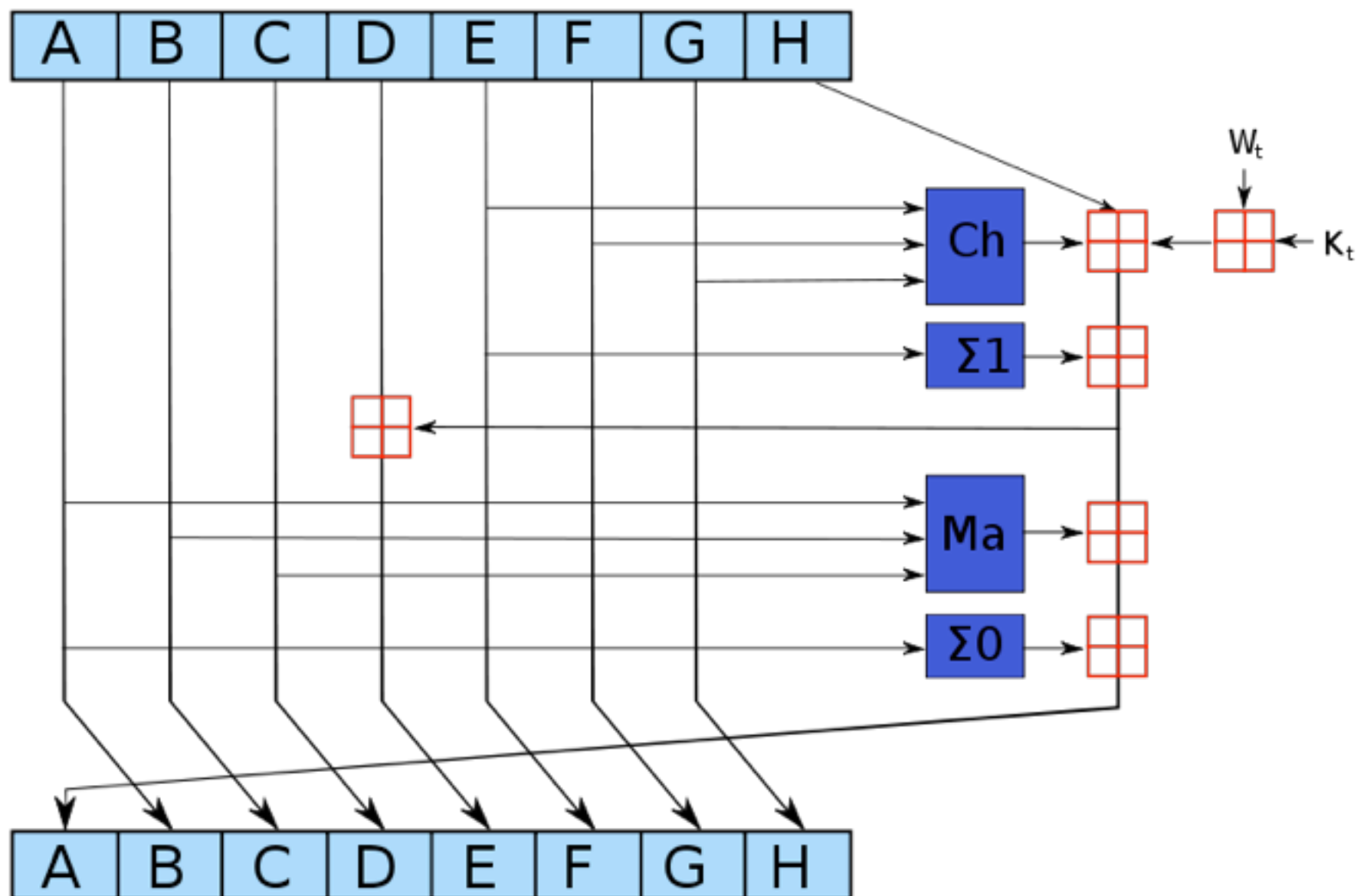


Message schedule

- $W_0 = M_0$
- ...
- $W_{15} = M_{15}$
- $W_{16} = W_0 + \sigma_0(W_1) + W_9 + \sigma_1(W_{14})$
- ...
- $W_{63} = W_{47} + \sigma_0(W_{48}) + W_{56} + \sigma_1(W_{61})$

- $\sigma_0(W) = \text{ROTR}^7(W) \oplus \text{ROTR}^{18}(W) \oplus \text{SHR}^3(W)$
- $\sigma_1(W) = \text{ROTR}^{17}(W) \oplus \text{ROTR}^{19}(W) \oplus \text{SHR}^{10}(W)$

Single SHA-256 round



- $K_0 = 428a2f98$
- $K_1 = 71374491$
- ...
- $K_{63} = c67178f2$
- cube roots of first 64 primes

One iteration in a SHA-2 family compression function. The blue components perform the following operations:

$$\text{Ch}(E, F, G) = (E \wedge F) \oplus (\neg E \wedge G)$$

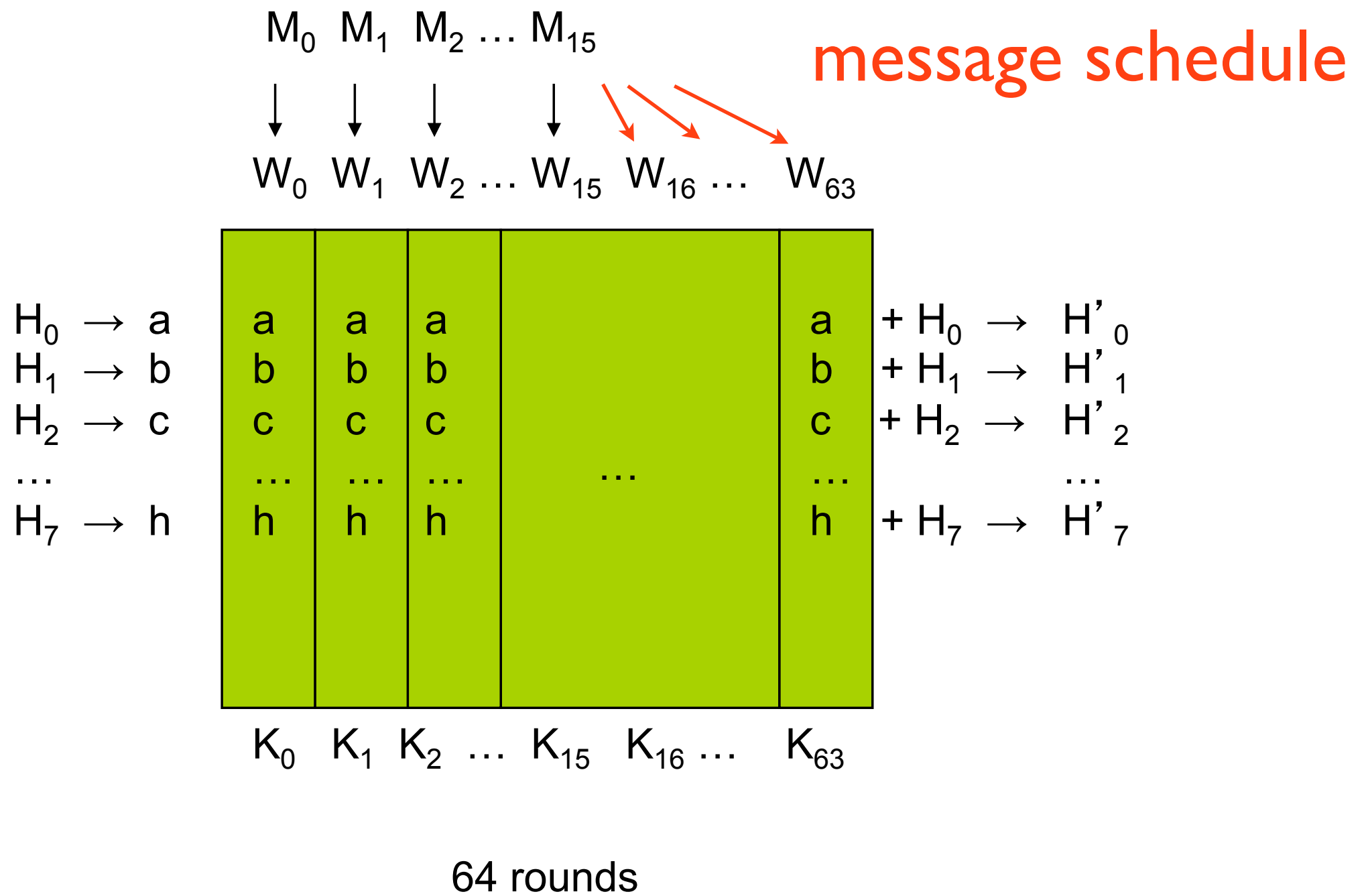
$$\text{Ma}(A, B, C) = (A \wedge B) \oplus (A \wedge C) \oplus (B \wedge C)$$

$$\Sigma_0(A) = (A \ggg 2) \oplus (A \ggg 13) \oplus (A \ggg 22)$$

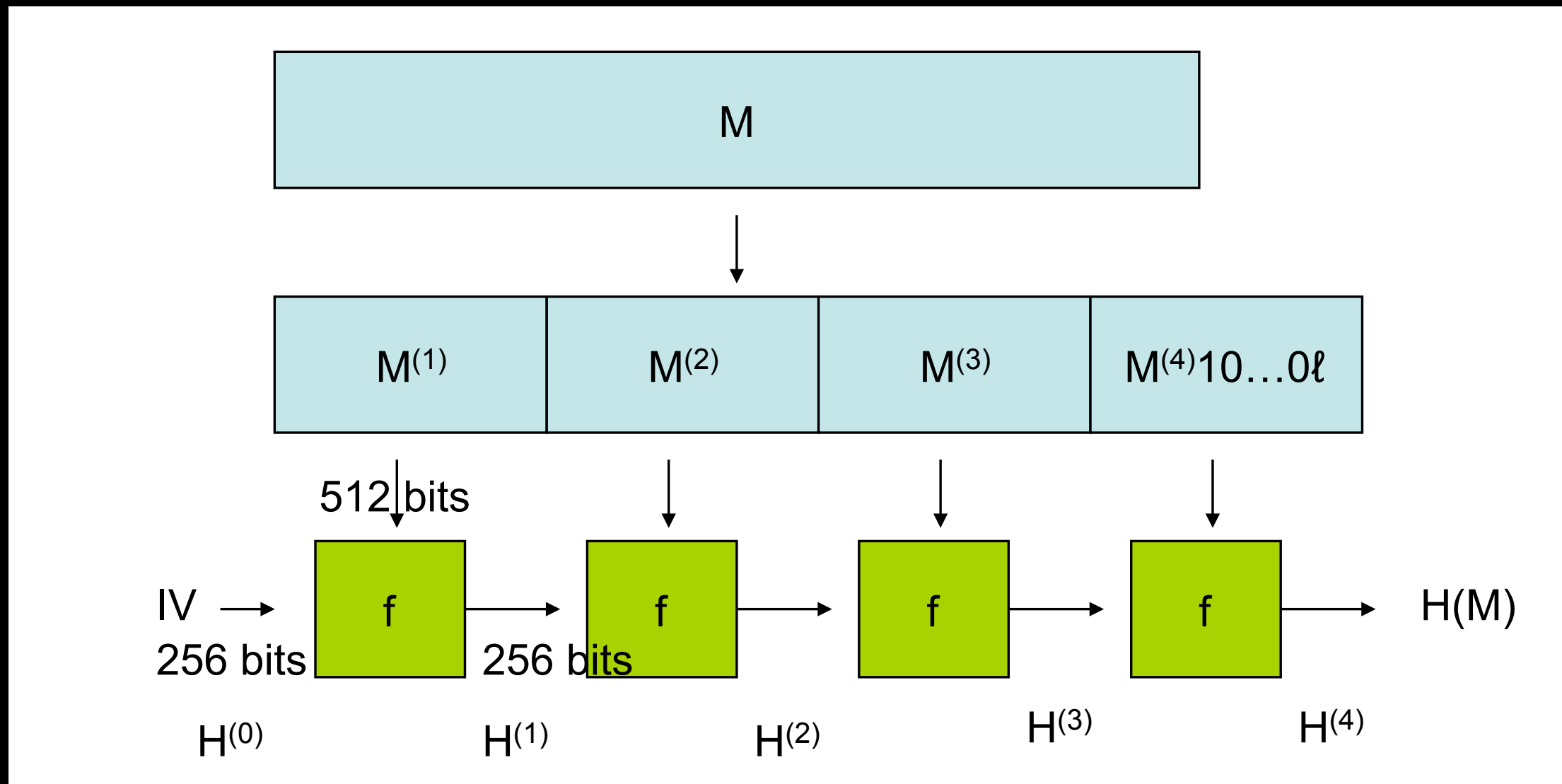
$$\Sigma_1(E) = (E \ggg 6) \oplus (E \ggg 11) \oplus (E \ggg 25)$$

The bitwise rotation uses different constants for SHA-512. The given numbers are for SHA-256. The red \boxplus is an addition modulo 2^{32} .

SHA-256 compression function



SHA-256



- SHA-256: $\{0,1\}^* \rightarrow \{0,1\}^{256}$ for input messages of length $|M| < 2^{64}$, i.e. M of size at most 2 billion GigaBytes
- **believed to be hard** to find a collision

Number Of Hashes to Find Collision

	Output bits	Birthday	Shortcut
MD4	128	2^{64}	2^2
RIPEMD	128	2^{64}	2^{18}
MD5	160	2^{80}	2^{21}
RIPEMD-160	160	2^{80}	
SHA-0	160	2^{80}	2^{34}
SHA-1	160	2^{80}	2^{51}
SHA-224	224	2^{112}	
SHA-256	256	2^{128}	

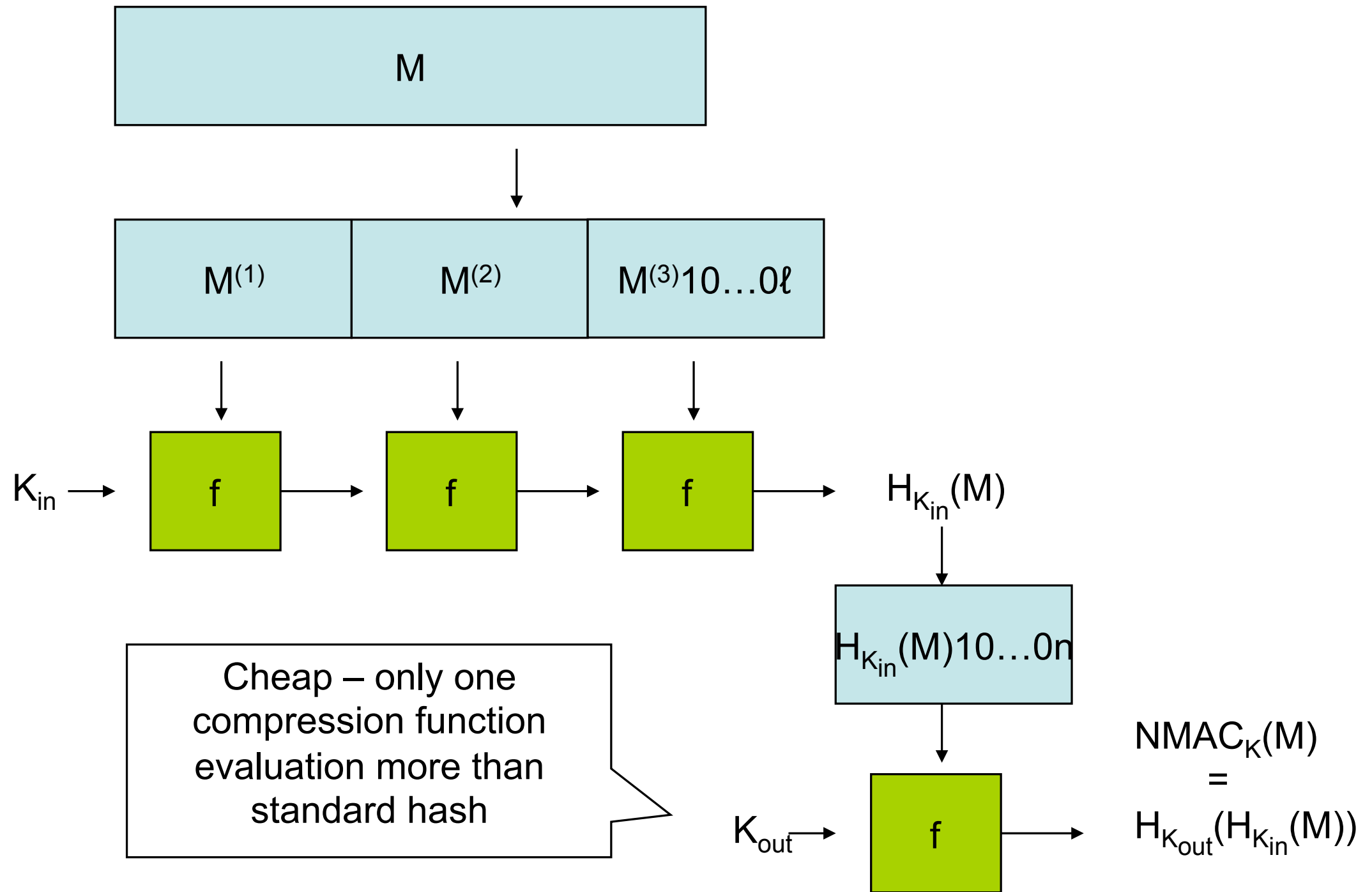
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SHA-224	224	2^{112}	
SHA-256	256	2^{128}	

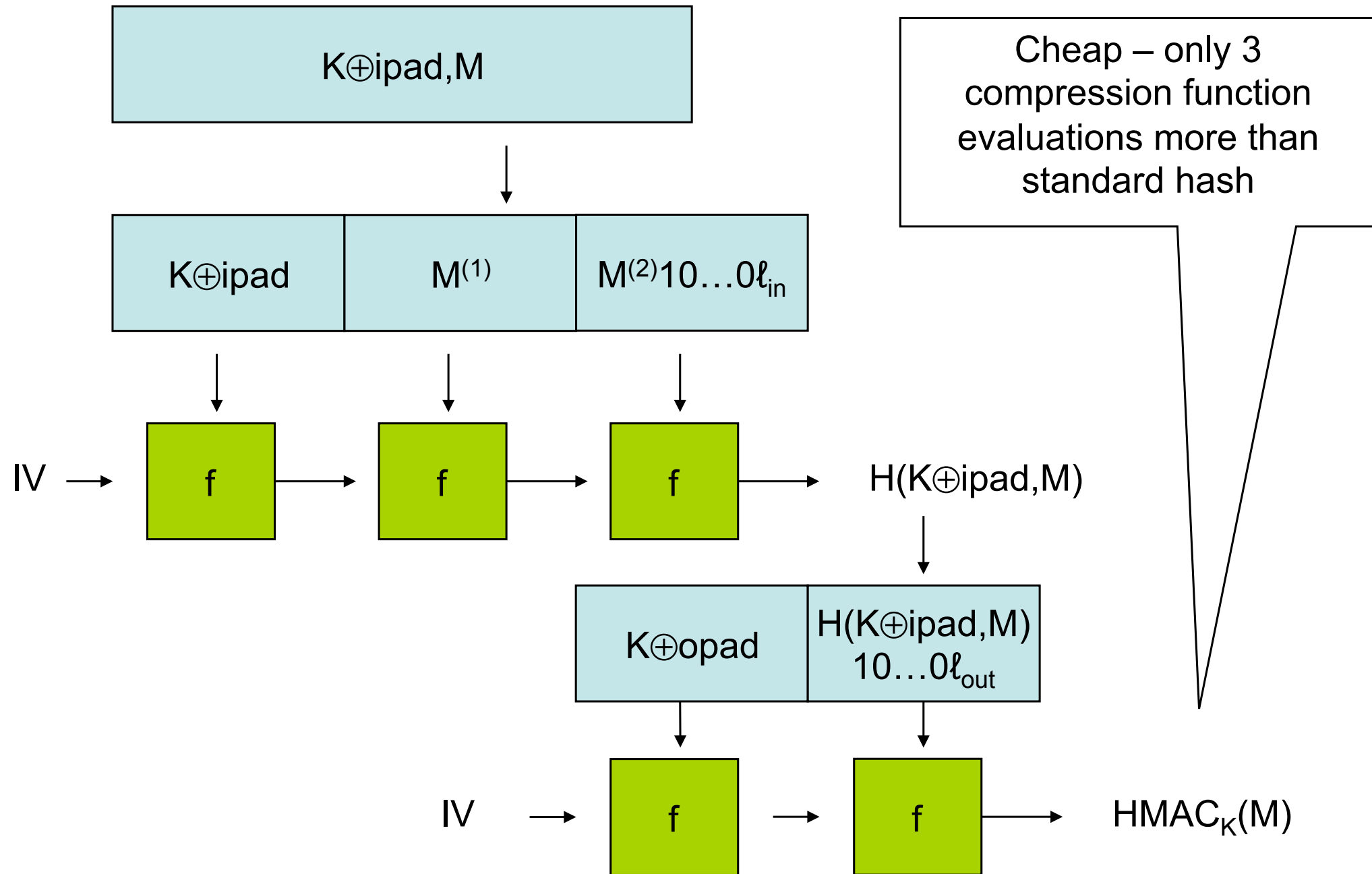
creating a rogue
Certification Authority
certificate

MACs from Hash Functions

NMAC-SHA-256

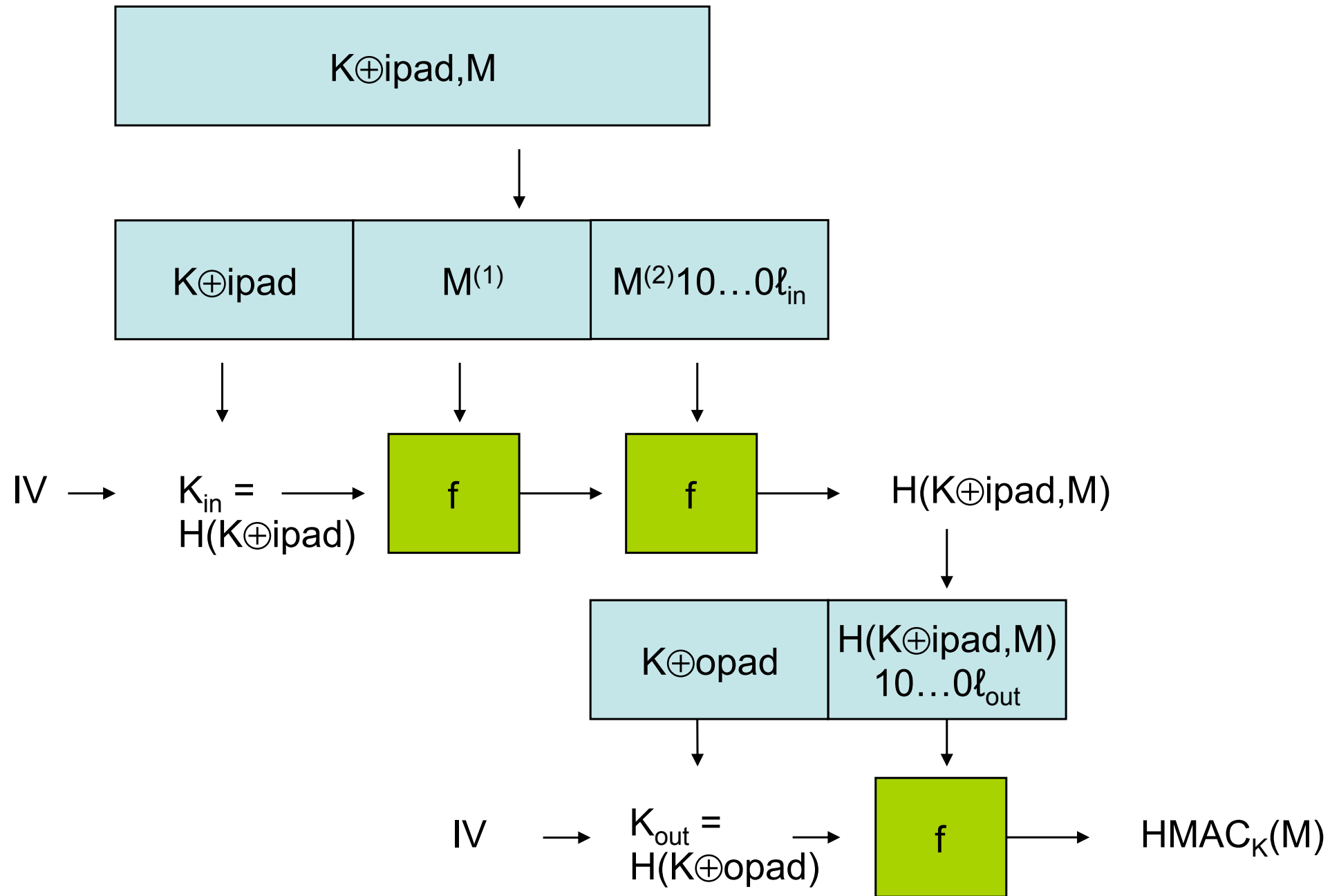


HMAC-SHA-256



- $\text{HMAC}_K(M) = H(K \oplus \text{opad}, H(K \oplus \text{ipad}, M))$
 $\text{ipad} = 3636\dots36$, $\text{opad} = 5c5c\dots5c$

HMAC vs NMAC



- HMAC can be seen as variant of NMAC
- as secure if $H(x \oplus \text{ipad}), H(x \oplus \text{opad})$ is pseudorandom