## Solutions for Exercises

## Chapter 7

140 Show that the following are  $\Pi_1^{\text{ZF}}$  or  $\Pi_1^{\text{ZFC}}$ . Solution.

- 1.  $y = \wp(x) \iff \forall z(z \subset x \leftrightarrow z \in y)$
- 2.  $x <_1 y \iff z_{\text{FC}} \forall z \neg (z \text{ is an injection } y \rightarrow x)$
- 3.  $\alpha$  is an initial number  $\Leftrightarrow \alpha \in OR \land \alpha \notin \omega \land \forall \beta \in \alpha (\beta <_1 \alpha)$
- 4.  $\gamma < cf(\alpha) \iff \forall x((x \subset \alpha \land \bigcup x = \alpha) \rightarrow \gamma <_1 x)$
- 5.  $\alpha$  is regular  $\Leftrightarrow \forall \gamma < \alpha(\gamma < cf(\alpha))$

What about the following?

- 1.  $x \leq_1 y$
- 2.  $x =_1 y$
- 4.  $\beta = cf(\alpha)$

These formulas are not  $\Pi_1^{\text{ZFC}}$ : if  $\phi(x, y)$  is a formula equivalent in ZFC to  $x \leq_1 y$  or  $x =_1 y$  or  $y = \operatorname{cf}(x)$ , then ZFC proves  $\exists x \in \operatorname{OR} \neg \phi(x, y)$ . If K is a countable transitive inner model of ZFC (or of a sufficient fragment of ZFC to prove the above sentence), then for some ordinal  $\alpha \in K, \ \neg \phi^K(\alpha, \omega)$ . But since K is countable,  $\alpha$  must be countable, so  $\phi(\alpha, \omega)$  holds, and hence  $\phi$  does not persist downward and is not  $\Pi_1$ .

- 3.  $\alpha = \omega_1$  is not  $\Pi_1^{\text{ZFC}}$ : if we have an inner transitive model K of (a sufficient fragment of) ZFC containing  $\omega_1$  such that for some countable  $\alpha \in K$ , K does not contain any bijection between  $\alpha$  and  $\omega$ , then  $\omega_1^K \leq \alpha < \omega_1$ . Constructing such a model is beyond the scope of this exercise, though.
- 5. " $\alpha$  is weakly inaccessible" is  $\Pi_1^{\text{ZF}}$ : it can be written as " $\alpha$  is an initial  $\land \alpha$  is regular  $\land \forall \beta < \alpha \exists \gamma < \alpha (\beta <_1 \gamma)$ ".

**142** Every  $L_{\alpha}$  is transitive. **L** is transitive. *Solution*.

Induction w.r.t.  $\alpha$ . The only non-trivial case is the successor step. Suppose that  $x \in L_{\alpha+1}$  and  $y \in x$ . Then  $x \subset L_{\alpha}$  and hence  $y \in L_{\alpha}$ . By IH,  $y \subset L_{\alpha}$ . Thus,  $y = \{u \in L_{\alpha} \mid (u \in y)^{L_{\alpha}}\} \in L_{\alpha+1}$ .  $\Box$ 

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- 1. For  $a \in \mathbf{L}$ , if  $\mathbf{L} \cap \wp(a) \subset \mathbf{L}_{\alpha}$ , then  $\mathbf{L} \cap \wp(a) \in \mathbf{L}_{\alpha+1}$ .
- 2. If  $a \in \mathbf{L}$ , then  $\mathbf{L} \cap \wp(a) \in \mathbf{L}$ .
- 3. *Thus*, the Powerset Axiom holds in L.

Solution.

- 1. If  $a \in \mathbf{L}$ , and  $\mathbf{L} \cap \wp(a) \subset \mathcal{L}_{\alpha}$ , then  $a \in \mathcal{L}_{\alpha}$ , and we can write  $\mathbf{L} \cap \wp(a) = \{ y \in \mathcal{L}_{\alpha} \mid (y \subset a)^{\mathcal{L}_{\alpha}} \} \in \mathrm{Def}(\mathcal{L}_{\alpha}) = \mathcal{L}_{\alpha+1}.$
- 2. Define the operation  $h : \mathbf{L} \to OR$  by h(x) = "the least  $\xi$  such that  $x \in L_{\xi}$ ". If  $a \in \mathbf{L}$ , then since  $\mathbf{L} \cap \wp(a) \subset \mathbf{L}$  we have  $\mathbf{L} \cap \wp(a) \subset L_{\alpha}$  and  $\mathbf{L} \cap \wp(a) \in L_{\alpha+1} \subset \mathbf{L}$  for  $\alpha = \bigcup_{x \in \mathbf{L} \cap \wp(a)} h(x)$ .
- 3. Now for any  $a, b \in \mathbf{L}$ ,  $b \in \mathbf{L} \cap \wp(a) \Leftrightarrow b \in \wp(a) \Leftrightarrow b \subset a \Leftrightarrow (b \subset a)^{\mathbf{L}}$ . So the Powerset Axiom holds in  $\mathbf{L}$ , and  $(\wp(a))^{\mathbf{L}} = \wp(a) \cap \mathbf{L}$ .

**152** Show that Collection holds in **L**. *Solution.* 

Let  $a \in \mathbf{L}$ , and suppose that  $\forall x \in a \exists y \in \mathbf{L}\Psi^{\mathbf{L}}$ . Define the operation  $h : a \to OR$  by h(x) = the least  $\xi$  such that  $\exists y \in \mathcal{L}_{\xi}\Psi^{\mathbf{L}}$ . Now construct  $\alpha = \bigcup \{h(x) \mid x \in a\}$ . Then  $\forall x \in a \exists y \in \mathcal{L}_{\alpha}\Psi^{\mathbf{L}}$ . Since  $\mathcal{L}_{\alpha} \in \mathrm{Def}(\mathcal{L}_{\alpha}) = \mathcal{L}_{\alpha+1} \subset \mathbf{L}$ , Collection holds.

## **157** Prove Lemma 7.28:

- 1. The intersection of two clubs is a club,
- 2. if each  $C_x$  (for every element x of a set a) is club, then so is  $\bigcap_{x \in a} C_x$ ,

3. if each  $C_{\xi}$  ( $\xi \in OR$ ) is club, then so is  $\left\{ \alpha \in OR \mid \alpha \in \bigcap_{\xi < \alpha} C_{\xi} \right\}$ .

## Solution.

- 1. This is a special case of (ii).
- 2. Let  $C_x$  be club for all  $x \in a$ . Now, if  $\alpha$  is a limit ordinal in which  $\bigcap_{x \in a} C_x$  is unbounded, then all the  $C_x$  are unbounded in  $\alpha$ , so  $\alpha \in \bigcap_{x \in a} C_x$ . To show that  $\bigcap_{x \in a} C_x$  is unbounded, let  $\alpha \in OR$ , We can define a sequence  $\alpha_{i \in \omega}$  by setting  $\alpha_0 = \alpha$ ,  $\alpha_{n+1,x} =$  "the smallest ordinal  $\geq \alpha_n$  in  $C_x$ ", and  $\alpha_{n+1} = \bigcup_{x \in a} \alpha_{n+1,x}$ . Now either  $\alpha' = \bigcup_{n \in \omega} \alpha_n$  is equal to  $\alpha_n$ for some  $n \in \omega$ , or it is a limit ordinal in which all the  $C_x$  are unbounded. In both cases,  $\alpha' \in \bigcap_{x \in a} C_x$ .
- 3. Let  $C_{\xi}$  be club for all  $\xi \in OR$ , and let  $C = \left\{ \alpha \in OR \mid \alpha \in \bigcap_{\xi < \alpha} C_{\xi} \right\}$ . If  $\alpha$  is a limit ordinal in which C is unbounded, then for all  $\xi$ ,  $C_{\xi} \cup (\xi + 1) \supset C$  is unbounded in  $\alpha$ , and therefore for  $\xi < \alpha$  so is  $C_{\xi}$ , and  $\alpha \in C_{\xi}$ . It follows that  $\alpha \in C$ . To show that C is unbounded, let  $\alpha \in OR$ . We can define a sequence  $\alpha_{i \in \omega}$  by setting  $\alpha_0 = \alpha$  and  $\alpha_{n+1} =$  "the smallest ordinal  $\geq \alpha_n$  in  $\bigcap_{\xi < \alpha_n} C_{\xi}$ ". Now either  $\alpha' = \bigcup_{n \in \omega} \alpha_n$  is equal to  $\alpha_n$  for some  $n \in \omega$ , or it is a limit ordinal in which all the  $C_{\xi}$  are unbounded for  $\xi < \alpha'$ . In both cases,  $\alpha' \in C$ .

**159.** Show that, in the reflection principle,  $\{\alpha \mid A_{\alpha} \prec_{\Sigma} A\}$  is closed. *Solution*.

Assume that  $C_{\Sigma} = \{\xi \in \text{OR} \mid A_{\xi} \prec_{\Sigma} A\}$  is unbounded in the limit ordinal  $\alpha$ . We need to show that, for  $\Phi \in \Sigma$ , the equivalence  $\Phi^{A_{\alpha}} \leftrightarrow \Phi^{A}$  holds on parameters from  $A_{\alpha}$ .

Induction w.r.t. the nr of logical symbols in  $\Phi$ . The only problem arises when  $\Phi$  is a quantification; say,  $\Phi = \exists z \Psi(x, y, z)$ . Assume that  $a, b \in A_{\alpha}$ .

 $(\Rightarrow)$  Suppose that  $[\exists z \Psi(a, b, z)]^{A_{\alpha}}$ . Say,  $c \in A_{\alpha}$  is such that  $[\Psi(a, b, c)]^{A_{\alpha}}$ . By IH on  $\Psi$ , it follows that  $[\Psi(a, b, c)]^{A}$ . Hence,  $[\exists z \Psi(a, b, z)]^{A}$ .

 $(\Leftarrow)$  Conversely, suppose that  $[\exists z \Psi(a, b, z)]^A$ . Choose  $\xi < \alpha$  in  $C_{\Sigma}$  such that  $a, b \in A_{\xi}$ . Then since  $\xi \in C_{\Sigma}$ , we also have that  $[\exists z \Psi(a, b, z)]^{A_{\xi}}$ . Thus,  $c \in A_{\xi}$  exists such that  $[\Psi(a, b, c)]^{A_{\xi}}$ , and it follows that  $[\Psi(a, b, c)]^A$ . By IH on  $\Psi$ ,  $[\Psi(a, b, c)]^{A_{\alpha}}$ . Therefore,  $[\exists z \Psi(a, b, z)]^{A_{\alpha}}$ .

Note that, by the Reflection Principle,  $\{\alpha \mid A_{\alpha} \prec_{\Sigma} A\}$  is unbounded as well, and therefore club.

**160.** Suppose that the initial  $\lambda$  is strongly inaccessible (Definition 6.24 p. 50). Show that  $\alpha < \lambda$  exists such that  $V_{\alpha} \prec V_{\lambda}$ . Show that the smallest such  $\alpha$  has  $cf(\alpha) = \omega$ . Solution.

Since  $V_{\lambda}$  models ZF, for any finite subformula-closed set of formulas  $\Sigma$ ,  $C_{\Sigma} = \{ \alpha \mid A_{\alpha} \prec_{\Sigma} A \}$ 

is closed and unbounded in  $\lambda$ . Let  $\alpha_{\Sigma}$  be the smallest element of  $C_{\Sigma}$ . Technically we need the apparatus from Section 7.6 to be able to use the notion of satisfaction and range over sets of formulas. Assuming this apparatus, set  $A = \{\alpha_{\Sigma} \mid \Sigma \text{ a finite subformula-closed set of formulas}\}$  and  $\alpha = \bigcup A$ .

Let  $\Sigma$  be an arbitrary finite subformula-closed set of formulas. Now for any  $\beta < \alpha$  we can find  $\Sigma'$  such that  $\beta < \alpha_{\Sigma'}$ , and hence  $\beta < \alpha_{\Sigma \cup \Sigma'} \in C_{\Sigma}$ . It follows that either  $\alpha = \alpha_{\Sigma \cup \Sigma'}$  for some  $\Sigma'$ , or  $A \cap C_{\Sigma}$  is unbounded in  $\alpha$ , and in both cases  $\alpha \in C_{\Sigma}$  and  $V_{\alpha} \prec_{\Sigma} V_{\lambda}$ . Since we chose  $\Sigma$  arbitrary,  $V_{\alpha} \prec V_{\lambda}$ . For any  $\beta$  with this property, we have  $\alpha_{\Sigma} \leq \beta$  for all  $\Sigma$ , and therefore  $\alpha \leq \beta$ .

Now for any finite  $\Sigma$ ,  $\exists \beta(V_{\beta} \prec_{\Sigma} V)$  can be expressed as a first-order formula, Obviously this formula holds in  $V_{\lambda}$ . If it were to hold in  $V_{\alpha_{\Sigma}}$ , then for some  $\beta < \alpha_{\Sigma}$  we would have  $V_{\beta} \prec_{\Sigma} V_{\alpha_{\Sigma}} \prec_{\Sigma} V_{\lambda}$ , contradicting the minimality of  $\alpha_{\Sigma}$ . Thus  $V_{\alpha_{\Sigma}} \not\prec V_{\lambda}$ , and  $\alpha_{\Sigma} < \alpha$ . It follows that  $\{\alpha_{\Sigma}\}$  is a countable cofinal subset of  $\alpha$ .