## Solutions to Exercises

## Chapter 4

54. Assume that the set $a$ is transitive. Show:
55. $a \in \mathbf{G}$ iff $\in$ is well-founded on $a$,
56. $a \subset \mathbf{T R}$ iff $\in$ is transitive on $a$.

Thus, an ordinal is the same as a transitive set on which $\in$ is a transitive and well-founded relation. (This is the standard definition of the notion.)
Solution.

1. If $a \notin \mathbf{G}$, then there would exist a $b \ni a$ such that $\forall y \in b(y \cap b \neq \emptyset)$, and then $a \cap b$ would be a nonempty subset of $a$ with no $\in$-minimal element (since for any $y \in a \cap b, y \subset a$ by transititive of $a$, and so $y \cap(a \cap b)=y \cap b \neq \emptyset)$. Conversely, if $c \subset a$ were a nonempty set without an $\in$-minimal element, then $b=c \cup\{a\}$ would satisfy $\forall y \in b: y \cap b \neq \emptyset$, and $a \notin \mathbf{G}$. It follows that $a \in \mathbf{G}$ iff $\in$ is well-founded on $a$.
2. $a \subset \mathbf{T R}$ iff for all $x, y, z$ with $z \in a,(x \in y \in z \Rightarrow x \in z)$.
$\in$ is transitive on $a$ iff for all $x, y, z$ with $x, y, z \in a,(x \in y \in z \Rightarrow x \in z)$.
By transitivity of $a$, the two are equivalent.
3. Show:
4. $\alpha \leq \beta \Leftrightarrow \alpha \subset \beta$,
5. if $K$ is a non-empty class of ordinals, then $\bigcap K$ is the least element of $K$,
6. if $A$ is a set of ordinals, then $\bigcup A$ is an ordinal that is the sup of $A$ (the least ordinal $\geq$ every $\alpha \in A$ ).

## Solution.

1. $(\alpha \in \beta \vee \alpha=\beta) \Rightarrow \alpha \subset \beta \Rightarrow \beta \notin \alpha \Rightarrow(\alpha \in \beta \vee \alpha=\beta)$
2. A reformulation of transfinite induction yields: if $K \subset O R$ is nonempty, then $\exists \alpha \in O R:(\forall \beta<$ $\alpha(\beta \notin K)) \wedge \alpha \in K$, i.e. $K$ has a least element. Let $\alpha \in K$ be the least element. Then $\bigcap K \leq \alpha$, and $\alpha \leq \bigcap K$ is trivial.
3. $\bigcup K$ is the union of transitive sets of ordinals, therefore itself a transitive set of ordinals, therefore an ordinal. It is also the smallest set containing as subsets all the sets in $K$, so by (1) it is the sup of $K$.
4. Assume that $(A, \prec)$ is a well-ordering and $B \subset A$.

Show that type $(B, \prec) \leqslant \operatorname{type}(A, \prec)$.
Solution. Suppose that $\alpha$ and $\beta$ are the order types of $A$ and $B$. If $\beta \nless \alpha$, then $\alpha<\beta$. Let $f: \beta \rightarrow B$ and $g: A \rightarrow \alpha$ be order-isomorphisms. Then $g f: \beta \rightarrow \alpha$ is an order-preserving injection such that $g f(\alpha)<\alpha$. This contradicts Lemma 4.8 (p. 29).
64. Prove Theorem 4.13: suppose that $\varepsilon$ is a well-founded relation on the class $\mathbf{U}$ such that for all $a \in \mathbf{U},\{b \in \mathbf{U} \mid b \varepsilon a\}$ is a set, then for every operation $H: \mathbf{V} \rightarrow \mathbf{V}$ there is a unique operation $F: \mathbf{U} \rightarrow \mathbf{V}$ such that for all $a \in \mathbf{U}$ :

$$
F(a)=H(F \mid\{b \in \mathbf{U} \mid b \varepsilon a\})
$$

Solution
Let us call a function or operation $f$ good if $\operatorname{Dom}(F) \subset \mathbf{U}$ and for all $a \in \operatorname{Dom}(F),\{b \in \mathbf{U} \mid$ $b \varepsilon a\} \subset \operatorname{Dom}(F)$ and $f(a)=H(f \mid\{b \in \mathbf{U} \mid b \varepsilon a\})$.
If $f$ and $g$ are both good, then they agree on their common domain: otherwise, if $a$ is the $\varepsilon$-minimal element on which they disagree, then $f(a)=H(f \mid\{b \in \mathbf{U} \mid b \varepsilon a\})=H(g \mid\{b \in \mathbf{U} \mid b \varepsilon a\})=g(a)$, a contradiction. It follows that the union $F$ of all good functions is an operation that satisfies the recursion equation on its domain.
Furthermore, $\operatorname{Dom}(F)=\mathbf{U}$. For otherwise, let $a$ be the $\varepsilon$-minimal element of $\mathbf{U}-\operatorname{Dom}(F)$. Then $\left\{b \in \mathbf{U} \mid b \varepsilon^{*} a\right\}$ would be a set, and setting $\left.f=F \mid\left\{b \in \mathbf{U} \mid b \varepsilon^{*} a\right\}\right)$ and $f^{\prime}=f \cup\{a, H(f)\}$ we would obtain a good function $f^{\prime}$ with $f^{\prime} \not \subset F$, a contradiction.
Since good operations on $\mathbf{U}$ have to agree on their domain, $F$ is unique.
65. Let $a_{0} \in \mathbf{V}$ be a set and $G: \mathbf{V} \rightarrow \mathbf{V}$ an operation. Show: there exists a unique operation $F:$ OR $\rightarrow \mathbf{V}$ on OR such that

- $F(0)=a_{0}$,
- $F(\alpha+1)=G(F(\alpha))$,
- for limits $\gamma: F(\gamma)=\bigcup_{\xi<\gamma} F(\xi)$.

Solution
Applying the Recursion Theorem on OR to the operation $H$ defined by

- $H(f)=G(f(\alpha))$ if $\operatorname{Dom}(f)=\alpha+1, \alpha \in \mathrm{OR}$
- $H(f)=\bigcup_{\xi<\gamma} f(\xi)$ if $\operatorname{Dom}(f)=\gamma$ and $\gamma$ is a limit ordinal $\neq 0$.
- $H(x)=a_{0}$ otherwise
yields an operation satisfying the requirements.

70. Show that the single recursion equation $H \uparrow \alpha=\bigcup_{\xi<\alpha} H(H \uparrow \xi)$ defines the same operation as the one defined in Definition 4.14 by three equations. (And, of course, $H \downarrow \alpha=\bigcap_{\xi<\alpha} H(H \downarrow \xi)$ is a single equation defining the greatest fixed point hierarchy - cf. Exercise 72.)
Solution.
Claim: $H \uparrow \alpha \subset H \uparrow(\alpha+1)$.
Induction. Obviously, $H \uparrow 0 \subset H \uparrow 1$. And if $H \uparrow \alpha \subset H \uparrow(\alpha+1)$, then, by monotonicity, $H \uparrow(\alpha+1) \subset$ $H \uparrow(\alpha+2)$. Finally, if $\gamma$ is a limit, then, if $\xi<\gamma, H \uparrow \xi \subset \bigcup_{\xi<\gamma} H \uparrow \xi$, hence, $H(H \uparrow \xi) \subset H\left(\bigcup_{\xi<\gamma} H \uparrow \xi\right)$; and so $H \uparrow \gamma=\bigcup_{\xi<\gamma} H \uparrow \xi \subset \bigcup_{\xi<\gamma} H(H \uparrow \xi) \subset H\left(\bigcup_{\xi<\gamma} H \uparrow \xi\right)=H \uparrow(\gamma+1)$.
Now:
$H \uparrow 0=\emptyset=\bigcup_{\xi<0} H(H \uparrow \xi) ;$
$H \uparrow(\alpha+1)=H(H \uparrow \alpha)=H(H \uparrow \alpha) \cup H \uparrow \alpha($ since $H \uparrow \alpha \subset H(H \uparrow \alpha))=H(H \uparrow \alpha) \cup \bigcup_{\xi<\alpha} H(H \uparrow \xi)($ by IH $)$ $=\bigcup_{\xi<\alpha+1} H(H \uparrow \xi)$;
$H \uparrow \gamma=\bigcup_{\xi<\gamma} H \uparrow \xi=\bigcup_{\xi<\gamma} H(H \uparrow \xi)$ (since $H \uparrow \xi \subset H \uparrow(\xi+1) \subset H \uparrow \gamma$ ).
72 Let $H$ be a monotone operator over a set $\mathbf{U}$. The greatest fixed point hierarchy is the sequence $\{H \downarrow \alpha\}_{\alpha}$ recursively defined by

- $H \downarrow 0=\mathbf{U}$,
- $H \downarrow(\alpha+1)=H(H \downarrow \alpha)$,
- $H \downarrow \gamma=\bigcap_{\xi<\gamma} H \downarrow \xi($ for limits $\gamma)$.

Show that:

1. the hierarchy is descending, i.e., that $\alpha<\beta \Rightarrow H \Downarrow \beta \subset H \downarrow \alpha$.
2. some stage $H \rrbracket \alpha_{0}$ is a fixed point of $H$.
3. $H \rrbracket \alpha_{0}=\bigcap_{\alpha} H \downarrow \alpha$ is the greatest fixed point of $H$.

Try to generalize for the case where $\mathbf{U}$ may be a proper class.
Solution.
Consider the dual operator of $H^{d}: \wp\left(\mathbf{U}_{\rightarrow} \wp(\mathbf{U})\right.$ of $H$ defined by $H^{d}(X)={ }_{\text {def }} \mathbf{U}-H(\mathbf{U}-X)$. By exercise 52 , the least fixed point of $H^{d}$ corresponds to the greatest fixed point of $H$. It is easy to see that the least fixed point hierarchy of $H^{d}$ also corresponds exactly to the greatest fixed point hierarchy of $H$, i.e. $H \downarrow \alpha=\mathbf{U}-H^{d} \uparrow \alpha$ for $\alpha \in O R$. The desired properties follow.
If $\mathbf{U}$ is a class, then the first property will still hold, but the second and third properties may fail. For instance, consider the operator $H: \mathbf{V} \rightarrow \mathbf{V}$ given by $H(X)=X-\{\alpha\}$ if $\alpha$ is the least ordinal in $X$, and $H(X)=\emptyset$ if $X$ does not contain any ordinals. Then for all ordinals $\alpha, H \downarrow \alpha=\mathbf{V}-\alpha$, and yet the only fixed point of $H$ is $H \downarrow=H \uparrow=\emptyset$.

