

UNIVERSITEIT VAN AMSTERDAM INSTITUTE FOR LOGIC, LANGUAGE AND COMPUTATION

> Core Logic 2005/2006; 1st Semester dr Benedikt Löwe

Homework Set #9

Deadline: November 15th, 2005

Exercise 29 (total of six points).

Let  $\mathcal{L} = \{R\}$  be a language with one binary relation symbol. Consider the following seven  $\mathcal{L}$ -sentences:

$$\begin{split} \varphi_{(i)} &:= \quad \forall x \neg Rxx \\ \varphi_{(ii)} &:= \quad \forall x \forall y (x \neq y \rightarrow (Rxy \lor Ryx)) \\ \varphi_{(iii)} &:= \quad \forall x \forall y \forall z ((Rxy \land Ryz) \rightarrow Rxz) \\ \varphi_{(iv)} &:= \quad \forall x \exists y \exists z (Ryx \land Rxz) \\ \varphi_{ME} &:= \quad \exists x \forall y (Ryx \lor x = y) \\ \varphi_{LEP} &:= \quad \forall x \exists y \forall z (Rxz \rightarrow (Rzy \lor y = z)) \end{split}$$

Check whether the following sets of sentences are consistent. If they are, give a model. If they aren't, derive a contradiction (2 points each).

(1) { $\varphi_{(i)}, \varphi_{(iii)}, \varphi_{(iv)}, \varphi_{ME}$ }, (2) { $\varphi_{(i)}, \varphi_{(iii)}, \varphi_{LEP}, \neg \varphi_{ME}$ }, (3) { $\varphi_{(i)}, \varphi_{(ii)}, \varphi_{(iii)}, \varphi_{LEP}, \neg \varphi_{ME}$ },

## Exercise 30 (total of nine points).

We are modelling Achilles and the turtle as a transfinite process on the real line  $\mathbb{R}$ . Please give arguments for all answers.

(1) Achilles' position at time t is given by  $A_t$ , the turtle's position is given by  $T_t$ . We start with  $A_0 := 0$  and  $T_0 := 1$ . For every index i, we define  $A_{i+1} := A_i + |T_i - A_i|$ ,  $T_{i+1} := T_i + \frac{1}{2} \cdot |T_i - A_i|$ , and

$$T_{\infty} := \lim_{i \in \mathbb{N}} T_i,$$
$$A_{\infty} := \lim_{i \in \mathbb{N}} A_i,$$
$$T_{\infty + \infty} := \lim_{i \in \mathbb{N}} T_{\infty + i}, \text{ and}$$
$$A_{\infty + \infty} := \lim_{i \in \mathbb{N}} A_{\infty + i}.$$

Determine the least index i such that  $A_i = T_i$  (1 point). Where is Achilles at time  $\infty + \infty$  (2 points)?

(2) Now the positions are given by  $A_t^*$  and  $T_t^*$  defined as follows. For each index  $i \in \{0, 1, 2, ..., \infty, \infty + 1, \infty + 2, \infty + 3, ...\}$ , we define the *value* v(i) as follows:

$$\mathbf{v}(i) := n \text{ if } i = n \text{ or } i = \infty + n.$$

We start with  $A_0^* := 0$  and  $T_0^* := 1$ . For every index *i*, we define  $A_{i+1}^* := A_i^* + \frac{1}{2^{v(i)}}$ ,  $T_{i+1}^* := T_i^* + \frac{1}{2^{v(i)+1}}$ , and

$$T_{\infty}^* := \lim_{i \in \mathbb{N}} T_i^*,$$
  

$$A_{\infty}^* := \lim_{i \in \mathbb{N}} A_i^*,$$
  

$$T_{\infty+\infty}^* := \lim_{i \in \mathbb{N}} T_{\infty+i}^*, \text{ and }$$
  

$$A_{\infty+\infty}^* := \lim_{i \in \mathbb{N}} A_{\infty+i}^*.$$

Show that for every natural number n, we have  $T_n = T_n^* = A_{n+1} = A_{n+1}^*$  (2 points). Compute  $A_{\infty+5}^*$ ,  $T_{\infty+12}^*$ ,  $A_{\infty+\infty}^*$  and  $T_{\infty+\infty}^*$  (1 point each).

## Exercise 31 (total of seven points).

Consider the language of arithmetic  $\mathcal{L} = \{ \dot{+}, \dot{\times}, \dot{S}, \dot{0}, \dot{1}, \dot{<}, \dot{=} \}$  and its standard model  $\mathbf{N} := \langle \mathbb{N}, +, \cdot, \text{succ}, 0, 1, <, = \rangle$ . (Here succ is the successor function  $n \mapsto n + 1$ .) The language of arithmetic allows to define formulas that describe the natural numbers:

$$\chi_n(x) := x = \underbrace{\dot{S} \dots \dot{S}}_{n \text{ times}} \dot{0}$$

We say that a set of  $\mathcal{L}$ -sentences T is an *arithmetic* if  $\mathbf{N} \models T$ . Prove that every arithmetic has a model which is not isomorphic to N (7 points).

**Hint.** Define an extension  $\mathcal{L}^* := \mathcal{L} \cup \{\dot{c}\}$  of  $\mathcal{L}$  where  $\dot{c}$  is a constant symbol and look at the theory  $T^* := T \cup \{\neg \chi_n(\dot{c}); n \in \mathbb{N}\}$ . Prove that there is no value c of  $\dot{c}$  such that  $\langle \mathbf{N}, c \rangle$  is a model of  $T^*$ . Prove that  $T^*$  is consistent by using the compactness theorem. Use these two facts to prove the claim. (You may use that isomorphic models satisfy the same sentences.)

For students with a mathematical logic background: If T is sufficiently strong, then you can say that the interpretation of  $\dot{c}$  must be an "infinite element". Make this statement mathematically precise and prove it (3 extra points).